

# Intuitionistic Logic

## Spring 2012

### Homework 5

(due Monday, 2nd April)

1. (a) Show that the following is valid: If  $\vdash_{\mathbf{IPC}} (\varphi \rightarrow \psi) \rightarrow (\chi \vee \theta)$ , then  $\vdash_{\mathbf{IPC}} (\varphi \rightarrow \psi) \rightarrow \chi$  or  $\vdash_{\mathbf{IPC}} (\varphi \rightarrow \psi) \rightarrow \theta$  or  $\vdash_{\mathbf{IPC}} (\varphi \rightarrow \psi) \rightarrow \varphi$ . [3 pts]  
(b)\* Give an example such that the first two alternatives of (a) do not apply, but the last one does. [2 pts]  
(c) Show, using the above, that, if  $\vdash (\neg p \rightarrow q \vee r) \rightarrow \varphi \vee \psi$ , then  $\vdash (\neg p \rightarrow q \vee r) \rightarrow \varphi$  or  $\vdash (\neg p \rightarrow q \vee r) \rightarrow \psi$ . [2 pts]
2. (a) Let  $\varphi$  be a propositional formula containing only  $\wedge, \rightarrow$  and  $\perp$ .  
Let  $\mathfrak{M}$  be a model and  $w$  a node in  $\mathfrak{M}$  such that  $w$  has proper successors (i.e., there is at least one  $v$  in  $\mathfrak{M}$  with  $wRv$  and  $w \neq v$ ). Suppose
  - $\varphi$  holds in all proper successors of  $w$  (i.e., for all  $v$  with  $wRv, w \neq v$ , we have  $\mathfrak{M}, v \models \varphi$ ), and
  - for all propositional variables  $p$ , we have that  $p$  is true in  $w$  iff  $p$  is true in all proper successors of  $w$  (i.e.,  $w \in V(p)$  iff  $\forall v (wRv, w \neq v \Rightarrow v \in V(p))$ ).(In other words, the valuation in  $w$  is maximal for propositional variables considering persistency).  
Show that  $\varphi$  is true in  $w$ . [4 pts]  
(b) Show on the basis of the above that, if  $\varphi$  is a propositional formula not containing  $\vee$  and  $\vdash_{\mathbf{IPC}} \varphi \rightarrow \psi \vee \chi$ , then  $\vdash_{\mathbf{IPC}} \varphi \rightarrow \psi$  or  $\vdash_{\mathbf{IPC}} \varphi \rightarrow \chi$ . [2 pts]
3. (a) Show that, if  $\Gamma \mid \varphi$ , then  $\Gamma \vdash \varphi$ <sup>1</sup>. [1 pt]  
(b) Show that, if  $\Gamma \vdash \delta$  for each  $\delta \in \Delta$ , and  $\Delta \vdash \gamma$  for each  $\gamma \in \Gamma$  ( $\Gamma$  and  $\Delta$  are equivalent), then  $\Gamma \mid \varphi$  iff  $\Delta \mid \varphi$ . [1 pt]  
(c) Show that, if  $\vdash \chi \rightarrow \varphi \vee \psi$  implies  $\vdash \chi \rightarrow \varphi$  or  $\vdash \chi \rightarrow \psi$ , for all subformulas  $\varphi \vee \psi$  of  $\chi$ , then  $\chi \mid \chi$ . (Hint: Show first that, under this condition,  $\chi \mid \theta$  iff  $\chi \vdash \theta$ , for all subformulas  $\theta$  of  $\chi$ ). [3 pts]

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<sup>1</sup>erratum: Add to the definition 21 of  $\mid$  (p. 25) the clause,  $\Gamma \mid \perp$  iff  $\Gamma \vdash \perp$ .