

# 5th Exercise sheet Model Theory

## 21 Feb 2017

**Exercise 1** Show that DLO is not  $\lambda$ -categorical for any  $\lambda > \omega$ .

**Exercise 2** Show that the embedding  $(\mathbb{Q}, <) \subseteq (\mathbb{R}, <)$  is elementary.

**Exercise 3** By a *graph* we will mean a pair  $(V, E)$  where  $V$  is a non-empty set and  $E$  is a binary relation on  $V$  which is both symmetric and irreflexive. We will refer to the elements of  $V$  as the *vertices* and the elements of  $E$  as the *edges*. If  $xEy$  holds for two  $x, y \in V$ , we say that  $x$  and  $y$  are *adjacent*.

A graph  $(V, E)$  will be called *random* if for any two finite sets of vertices  $X$  and  $Y$  which are disjoint there is a vertex  $v \notin X \cup Y$  which adjacent to all of the vertices in  $X$  and to none of the vertices in  $Y$ . We will write  $RG$  for the theory of random graphs.

Show that the theory  $RG$  is  $\omega$ -categorical, and hence complete.

**Exercise 4** The aim of this exercise is to prove Beth's Definability Theorem.

Let  $L$  be a language a  $P$  be a predicate symbol not in  $L$ , and let  $T$  be an  $L \cup \{P\}$ -theory.  $T$  *defines  $P$  implicitly* if any  $L$ -structure  $M$  has at most one expansion to an  $L \cup \{P\}$ -structure which models  $T$ . There is another way of saying this: let  $T'$  be the theory  $T$  with all occurrences of  $P$  replaced by  $P'$ , another predicate symbol not in  $L$ . Then  $T$  *defines  $P$  implicitly* iff

$$T \cup T' \models \forall x_1, \dots, x_n (P(x_1, \dots, x_n) \leftrightarrow P'(x_1, \dots, x_n)).$$

$T$  *defines  $P$  explicitly*, if there is an  $L$ -formula  $\varphi(x_1, \dots, x_n)$  such that

$$T \models \forall x_1, \dots, x_n (P(x_1, \dots, x_n) \leftrightarrow \varphi(x_1, \dots, x_n)).$$

Show that  $T$  defines  $P$  implicitly if and only if  $T$  defines  $P$  explicitly.

*Hint:* Use the Craig Interpolation Theorem.