2nd Homework sheet Proof Theory

- Deadline: 14 November.
- Submit your solutions by handing them to the lecturer at the *beginning* of the lecture.
- The homework sheet consists of two pages and there are exercises on the other side.
- Good luck!

In this exercise we work in intuitionistic propositional logic and the following Hilbert-style system for intuitionistic propositional logic:

- (i) Its axiom schemes are: $\varphi \lor \varphi \to \varphi, \varphi \to \varphi \land \varphi, \varphi \to \varphi \lor \psi, \varphi \land \psi \to \varphi, \varphi \lor \psi \to \psi \lor \varphi, \varphi \land \psi \to \psi \land \varphi, \bot \to \varphi.$
- (ii) Its inference rules are: from φ and $\varphi \to \psi$ infer ψ ; from $\varphi \to \psi$ and $\psi \to \chi$ infer $\varphi \to \chi$; from $\varphi \land \psi \to \chi$, infer $\varphi \to (\psi \to \chi)$; from $\varphi \to (\psi \to \chi)$, infer $\varphi \land \psi \to \chi$; from $\varphi \to \psi$ infer $\varphi \lor \chi \to \psi \lor \chi$.

In particular, $\Gamma \vdash_{\mathrm{IL}} \varphi$ will mean that there is a proof in this Hilbert-style proof calculus for intuitionistic propositional logic of φ from Γ .

The aim of the exercise is to give a syntactic proof of the disjunction property for intuitionistic logic. If Γ is a set of propositional formulas and φ is a formula, then we define a new relation $\Gamma | \varphi$ by induction of φ , as follows:

$$\begin{split} \Gamma|p &:= \quad \Gamma \vdash_{\mathrm{IL}} p \text{ for any propositional variable } p \\ \Gamma|\perp &:= \quad \Gamma \vdash_{\mathrm{IL}} \perp \\ \Gamma|\psi \wedge \chi &:= \quad \Gamma|\psi \text{ and } \Gamma|\chi \\ \Gamma|\psi \lor \chi &:= \quad \Gamma|\psi \text{ or } \Gamma|\chi \\ \Gamma|\psi \to \chi &:= \quad (\Gamma|\psi \text{ implies } \Gamma|\chi) \text{ and } \Gamma \vdash_{\mathrm{IL}} \psi \to \chi \end{split}$$

(a) (30 points) Show by induction on the structure of φ that $\Gamma | \varphi$ implies $\Gamma \vdash_{\mathrm{IL}} \varphi$.

- (b) (20 points) Show that $\Gamma | \perp$ implies $\Gamma | \varphi$ for any formula φ .
- (c) (40 points) Show that if $\Gamma | \gamma$ for all $\gamma \in \Gamma$ and $\Gamma \vdash_{\mathrm{IL}} \varphi$, then $\Gamma | \varphi$. *Hint:* Use induction the derivation of $\Gamma \vdash_{\mathrm{IL}} \varphi$. You do not have to treat all cases: it suffices to treat two axioms schemes and two inference rules of your choice.
- (d) (10 points) Deduce from (a) and (c) that $\vdash_{\mathrm{IL}} \varphi \lor \psi$ implies $\vdash_{\mathrm{IL}} \varphi$ or $\vdash_{\mathrm{IL}} \psi$.