Introduction to Modern Cryptography



4th lecture:

Pseudorandom Functions and Chosen-Plaintext Security

some of these slides are copied from or heavily inspired by the University College London MSc InfoSec 2010 course given by Jens Groth Thank you very much!

PRG vs PRF

seed
$$s \in \{0,1\}^n$$

$$\longrightarrow PRGG$$

$$G(s) \in \{0,1\}^{\ell (n)}$$

$$key \ k \in \{0, I\}^n$$
 input $x \in \{0, I\}^n$
$$\xrightarrow{PRF \ F}$$

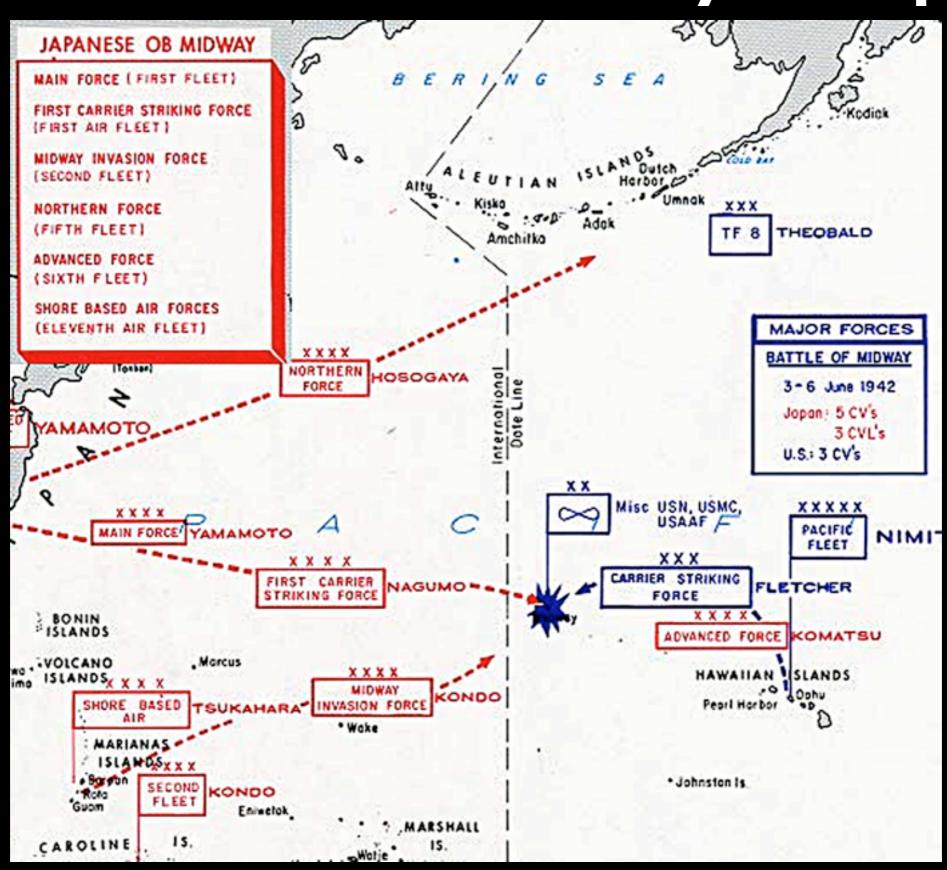
$$F_k(x) \in \{0, I\}^n$$

- existence of PRF ⇔ existence of PRG
- both can be based on one-way functions

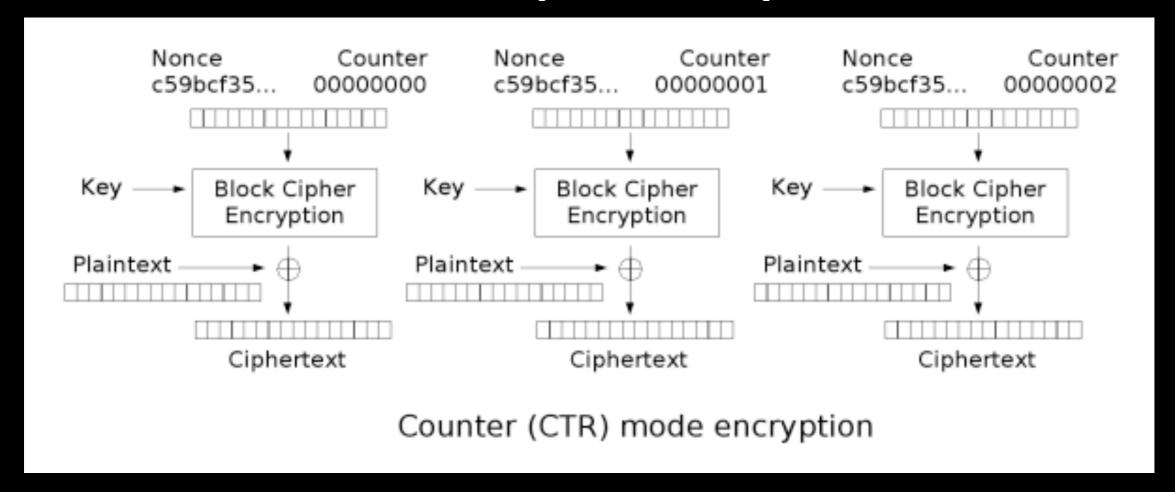
Battle of Midway (1942)

- Midway Atoll: Wikipedia, Google Maps
- important naval battle between the USA and Japan in World War II (<u>Wikipedia</u>)
- decided by cryptographic skills
- US tricked Japanese into acting as encryption oracle
- bottom line: the use of CPA secure encryption could have the course of world history

Battle of Midway Map

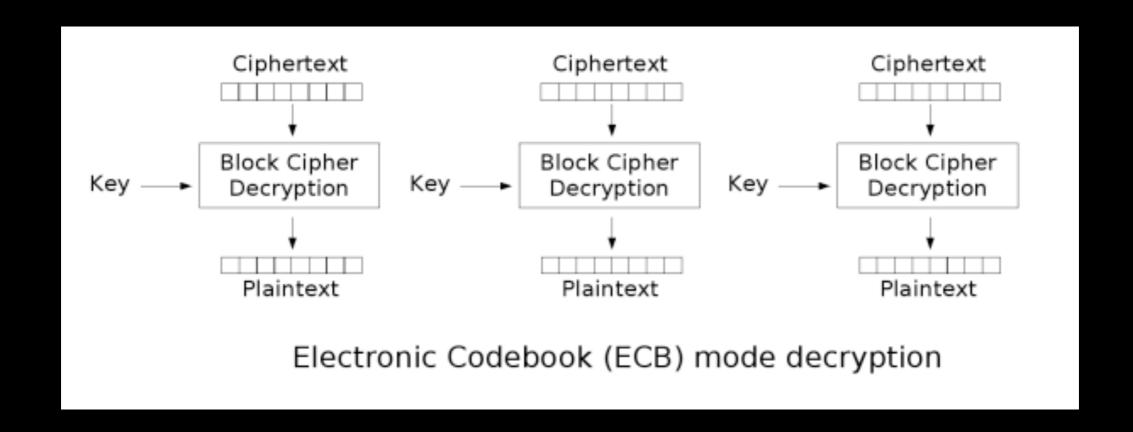


Counter (CTR) mode



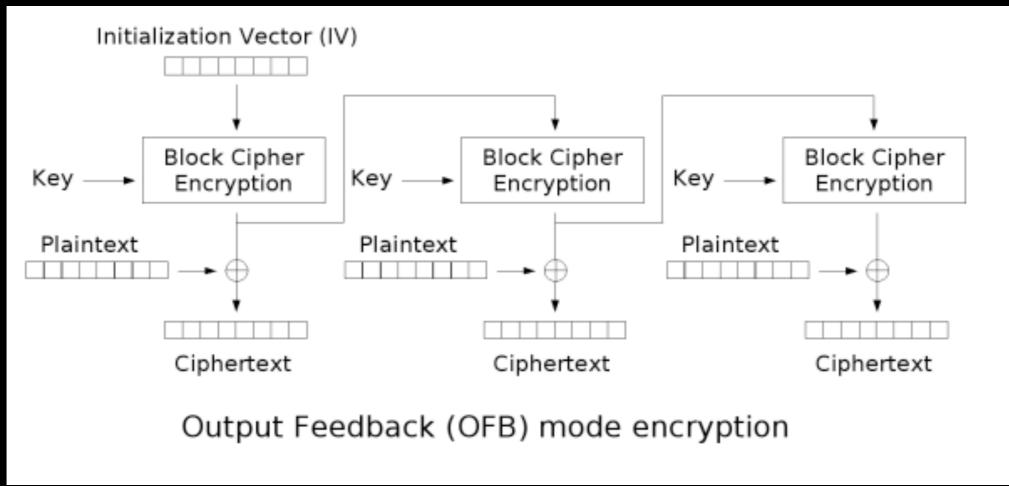
- CTR mode is CPA-secure if F (the Block Cipher) is a pseudorandom function
- can be precomputed and fully parallelized
- allows random access

Electronic Code Book (ECB)



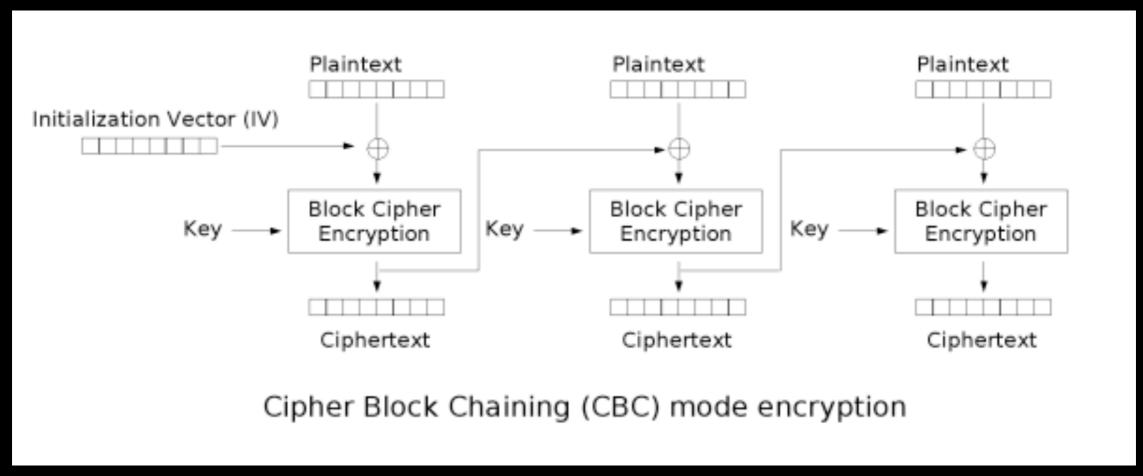
- highly insecure, should never be used
- see example on wikipedia

Output Feedback (OFB)



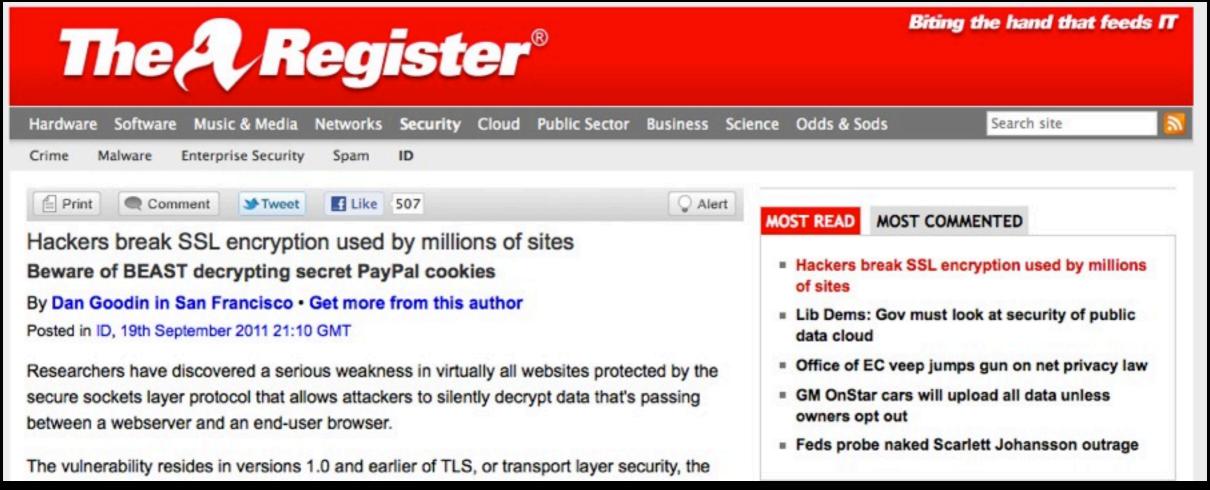
- if F is pseudorandom function, then OFB is CPA-secure
- advantage: pseudorandom stream can be precomputed

Cipher Block Chaining (CBC)



- if F is pseudorandom permutation, then CBC is CPA-secure
- drawback: encryption is sequential

Breaking News (23 Sep 2011)



- BEAST: Browser Exploit Against SSL/TLS
- error: reusing the IV of the last-sent CBC block for new connections, see <u>here</u>
- treated in this week's Exercise 3