## Information & Communication Exercise Sheet #3

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Out: Tuesday, 11 January 2016

(due: Friday, 15 January 2016, 14:00, by email or in my ILLC post box)

## To be solved in Class

- 1. **Huffman** [CT, 5.38] Find the Huffman *D*-ary code for  $(p_1, p_2, p_3, p_4, p_5, p_6) = (\frac{6}{25}, \frac{6}{25}, \frac{4}{25}, \frac{4}{25}, \frac{3}{25}, \frac{2}{25})$  and the expected codeword length
  - (a) for D = 2,
  - (b) for D = 4.
- 2. Error Penalty Suppose that an engineer believes that a source X can be described by the distribution  $Q_X$  given by the following table:

$$\begin{array}{c|ccccc} x & \texttt{a} & \texttt{b} & \texttt{c} \\ \hline Q_X(x) & ^1\!/_2 & ^1\!/_4 & ^1\!/_4 \end{array}$$

In fact, however, the source follows the distribution  $P_X$ :

$$\begin{array}{c|ccccc} x & \mathsf{a} & \mathsf{b} & \mathsf{c} \\ \hline P_X(x) & ^1\!/_4 & ^1\!/_2 & ^1\!/_4 \end{array}$$

- (a) Design a code for X based on the wrong distribution  $Q_X$ .
- (b) Design a code for X based on the correct distribution  $P_X$ .
- (c) Compute the expected number of bits per symbol used by each of these codes when X is sampled from  $P_X$ . How big is the difference?
- (d) Explain how this number relates to the Kullback-Leibler divergence

$$D(P_X || Q_X) = \sum_x P_X(x) \log \frac{P_X(x)}{Q_X(x)}.$$

- 3. Let X, Y, Z be binary random variables such that I(X; Y) = 0 and I(X; Z) = 0.
  - (a) Does it follow that I(X;Y,Z) = 0? If yes, prove it. If no, give a counterexample. **Hint:** Consider the case where X and Y are two independent uniform bits and  $Z = X \oplus Y$ .
  - (b) Does it follow that I(Y; Z) = 0? If yes, prove it. If no, give a counterexample.
- 4. For the Markov chain  $X \leftrightarrow Y \leftrightarrow \hat{X}$ , show that  $H(X|\hat{X}) \geq H(X|Y)$ .
- 5. [Cover-Thomas 2.32]. We are given the following joint distribution of  $X \in \{1, 2, 3\}$  and  $Y \in \{a, b, c\}$ :

$$P_{XY}(1,a) = P_{XY}(2,b) = P_{XY}(3,c) = 1/6$$
  
 $P_{XY}(1,b) = P_{XY}(1,c) = P_{XY}(2,a) = P_{XY}(2,c) = P_{XY}(3,a) = P_{XY}(3,b) = 1/12.$ 

Let  $\hat{X}(Y)$  be an estimator for X (based on Y) and let  $p_e = P(\hat{X} \neq X)$ .

- (a) Find an estimator  $\hat{X}(Y)$  for which the probability of error  $p_e$  is as small as possible.
- (b) Evaluate Fano's inequality for this problem and compare.

## Homework

- 1. Let X, Y, Z be arbitrary random variables, and let f be any deterministic function acting on  $\mathcal{Y}$ . In the following, replace "?" by " $\geq$ " or " $\leq$ " to obtain the correct inequalities, and reason each time with the help of an entropy diagram. **Hint:** H(f(Y)|Y) = 0.
  - (a) H(f(Y))?H(Y) 2 p.
  - (b) H(X|f(Y))?H(X|Y) 2 p.
  - (c) I(X;Z|Y) = 0 implies I(X;Z) ? I(X;Y) and I(X;Z) ? I(Y;Z).
- 2. For each statement below, specify a (different) joint distribution  $P_{XYZ}$  of random variables X, Y and Z such that the inequalities hold.
  - (a) There exists a y, such that H(X|Y=y)>H(X)
  - (b) I(X;Y) > I(X;Y|Z)
  - (c) I(X;Y) < I(X;Y|Z)

Note that the distributions have to be different from the ones seen as examples during the lecture.

3. Bottleneck. Suppose a Markov chain starts in one of n states, necks down to k < n states, and then fans back to m > k states. Thus  $X_1 \to X_2 \to X_3$ , i.e.,

$$P_{X_1X_2X_3}(x_1,x_2,x_3) = P_{X_1}(x_1) \cdot P_{X_2|X_1}(x_2|x_1) \cdot P_{X_3|X_2}(x_3|x_2)$$

for all  $x_1 \in \{1, 2, ..., n\}, x_2 \in \{1, 2, ..., k\}, x_3 \in \{1, 2, ..., m\}.$ 

- (a) Show that the (unconditional) dependence of  $X_1$  and  $X_3$  is limited by the bottleneck by proving that  $I(X_1; X_3) \leq \log k$ .
- (b) Evaluate  $I(X_1; X_3)$  for k = 1, and explain why no dependence can survive such a bottleneck.
- 4. Let A, B, C be random variables such that

$$I(A;B) = 0, (1)$$

$$I(A;C|B) = I(A;B|C), \tag{2}$$

$$H(A|BC) = 0. (3)$$

Which of the three relations  $\leq$ ,  $\geq$ , = holds between the quantities H(A) and H(C)? Prove your 3 p. answer.