

# Quantum Cryptography

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*BSc IW visit to ILLC*

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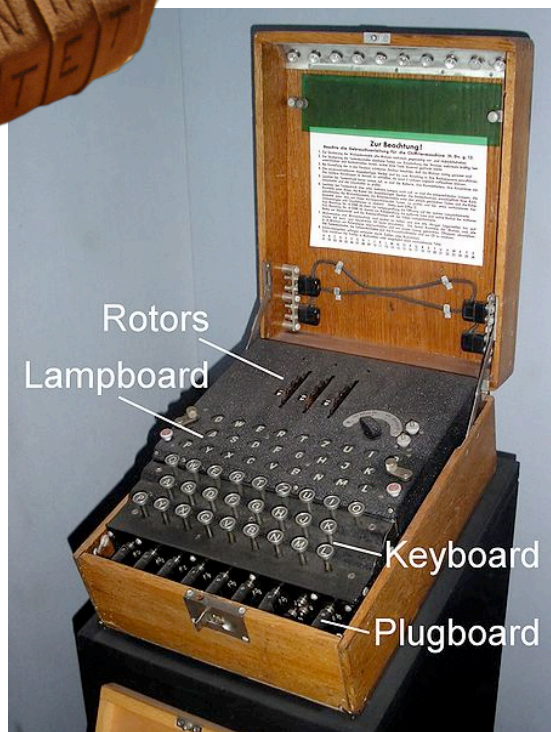


# Classical Cryptography

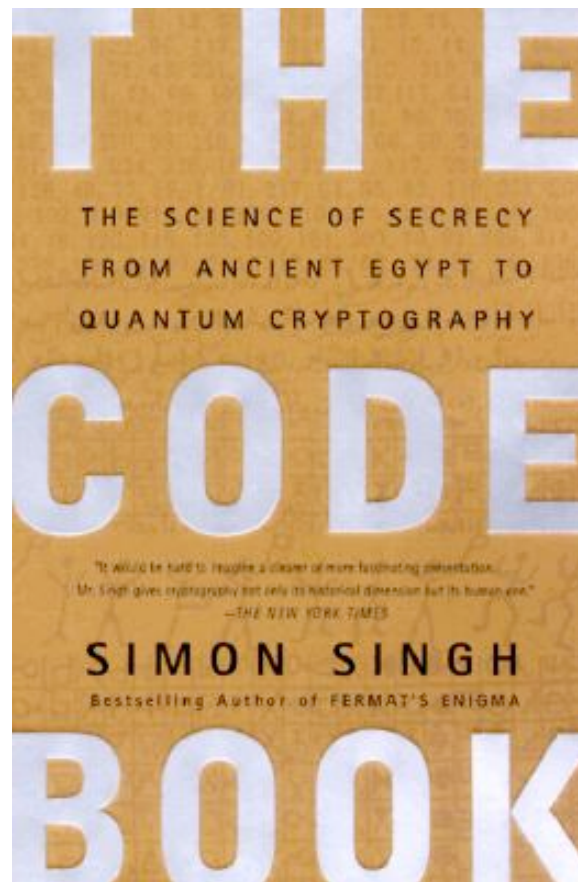
- 3000 years of fascinating history
- Until 1970: **private communication** was the only goal



Scytale



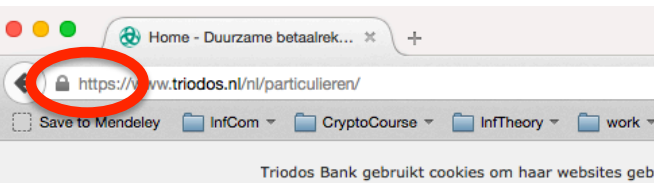
Enigma



# Modern Cryptography

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- is **everywhere!**
- is concerned with all settings where people **do not trust** each other



beleggers | Over Triodos Bank

theken

bankzake  
duurzaam

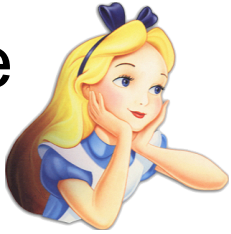


# Secure Encryption

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$m = \text{'doe you'}$

Alice



$k = 0101\ 1011$



Eve



Bob



$k = 0101\ 1011$

- Goal: Eve **does not learn** the message
- Setting: Alice and Bob share a secret key  $k$



# Perfectly Secure Encryption: One-Time Pad

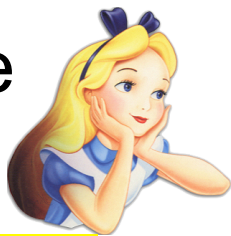
5

$m = 0000\ 1111$

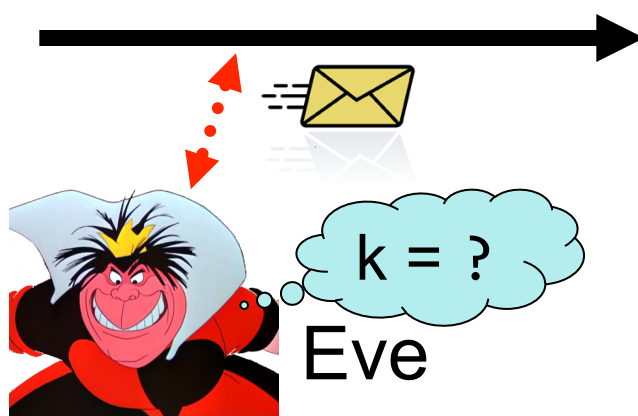
$c = m \oplus k = 0101\ 0100$

$m = c \oplus k = 0000\ 1111$

Alice



$k = 0101\ 1011$



Bob



$k = 0101\ 1011$

- Goal: Eve **does not learn** the message
- Setting: Alice and Bob share a key  $k$
- Recipe:

$m = 0000\ 1111$

$c = 0101\ 0100$

$k = 0101\ 1011$

$k = 0101\ 1011$

$c = m \oplus k = 0101\ 0100$

$c \oplus k = 0000\ 1111$

$c \oplus k = m \oplus k \oplus k = m \oplus 0 = m$

x	y	$x \oplus y$
0	0	0
0	1	1
1	0	1
1	1	0

- It is perfectly secure!

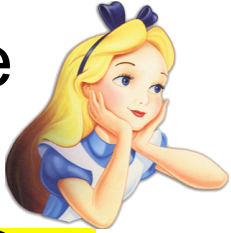
# Problems With One-Time Pad

$m = 0000\ 1111$

$c = m \oplus k = 0101\ 0100$

$m = c \oplus k = 0000\ 1111$

Alice



$k = 0101\ 1011$



Eve



Bob



$k = 0101\ 1011$

- The key has to be **as long as** the message.
- The key can only be **used once**.
- In practice, other encryption schemes (such as [AES](#)) are used which allow to encrypt long messages with short keys.

BSc Informatica course:  
Information & Communication

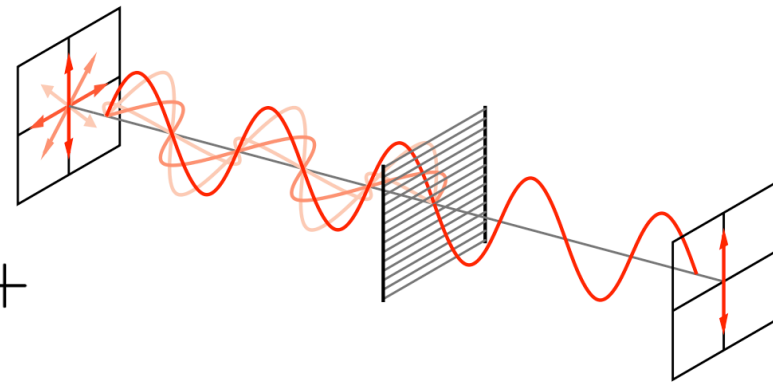
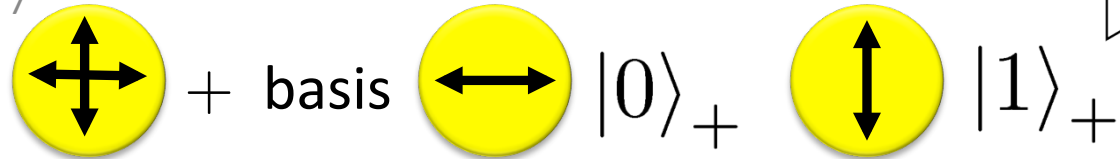
Master of Logic course:  
Information Theory

Master of Logic course:  
Modern Cryptography

Master of Logic course:  
Computational Complexity

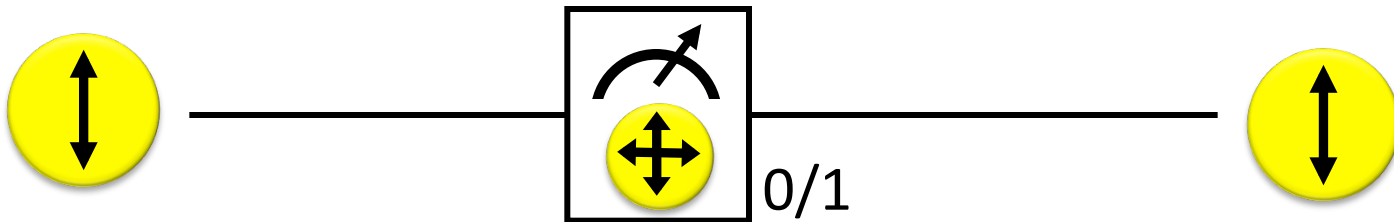
# Quantum Mechanics

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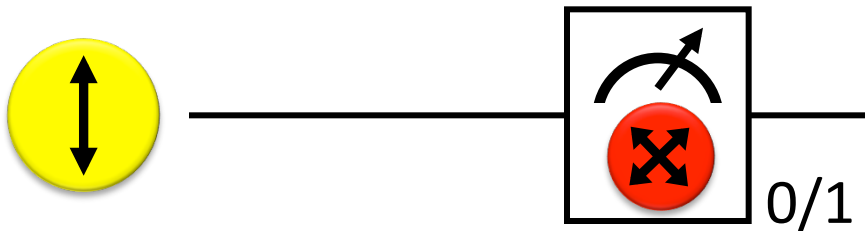


Measurements:

with prob. 1 yields 1

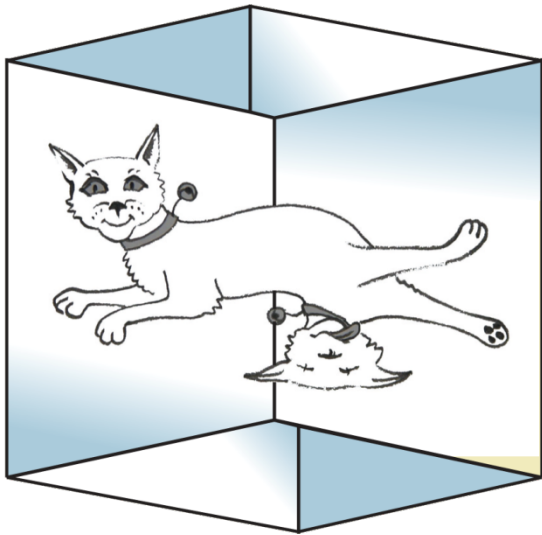


with prob.  $\frac{1}{2}$  yields 0

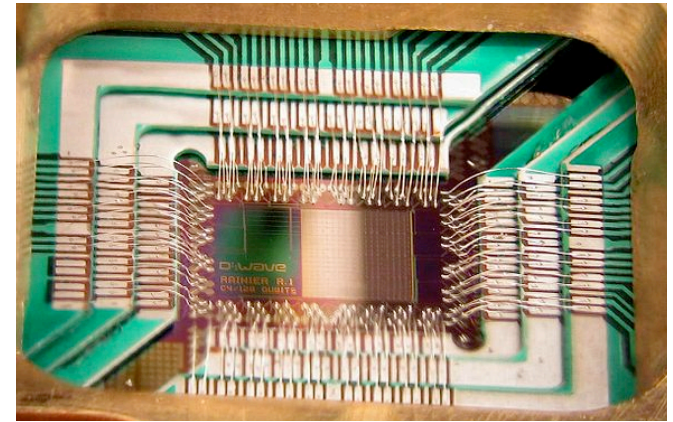


with prob.  $\frac{1}{2}$  yields 1





0



# Wonderland of Quantum Mechanics

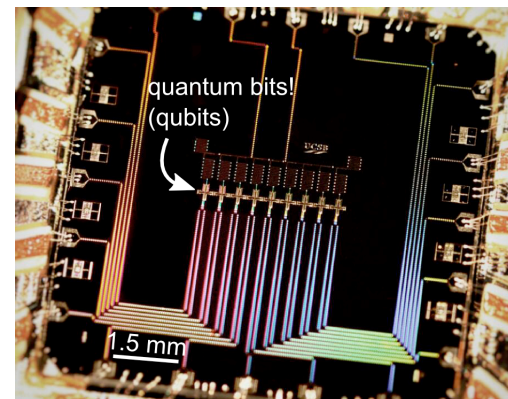
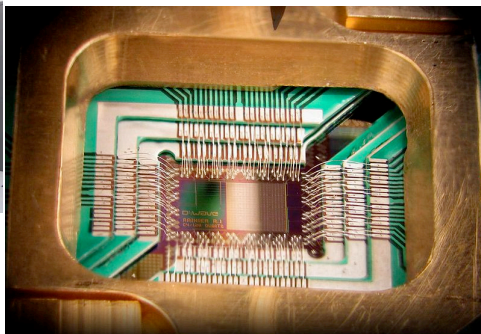




# Can We Build Quantum Computers?

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- Possible to build in theory, no fundamental theoretical obstacles have been found yet.



Martinis group (UCSB)  
9 qubits

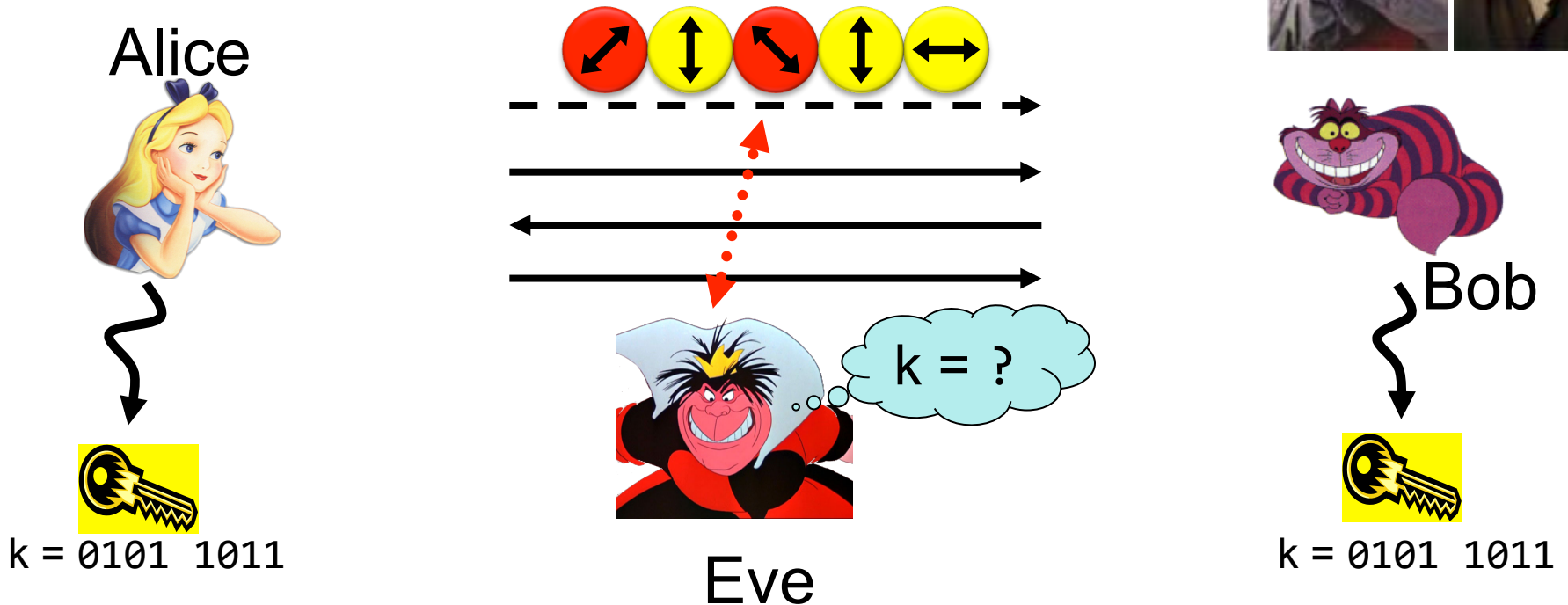
- Canadian company “D-Wave” claims to have build a quantum computer with 1024 qubits. Did they?
- 2014: Martinis group “[acquired](#)” by Google
- 2014/15: 135+50 Mio € investment Delft
- 2015: QuSoft center in Amsterdam

Master of Logic course:  
Quantum Computing

# Quantum Key Distribution (QKD)

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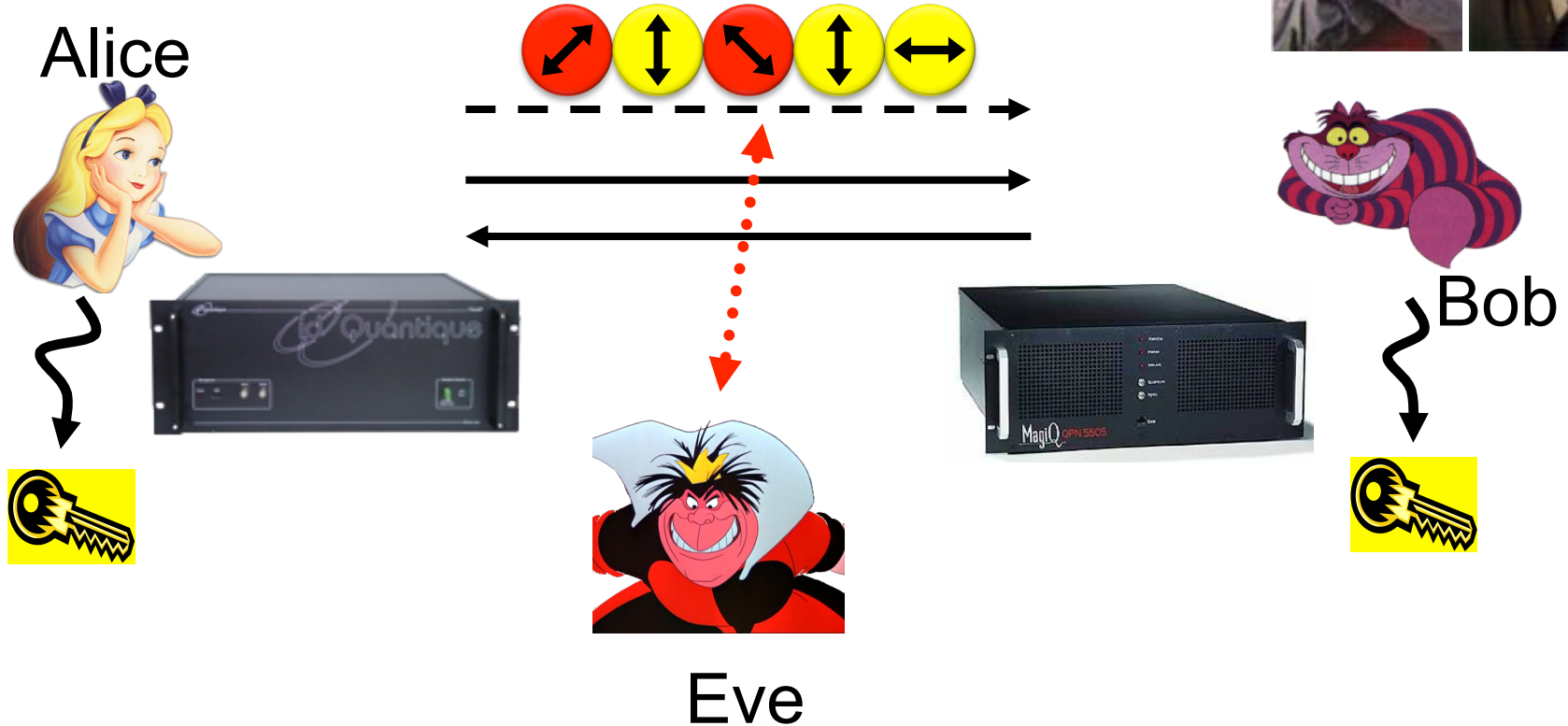
[Bennett Brassard 84]



- Offers a **quantum solution** to the key-exchange problem
- Puts the players into the starting position to use symmetric-key cryptography (such as the one-time pad)

# Quantum Key Distribution (QKD)

11 [Bennett Brassard 84]



- **technically feasible**: no quantum computer required, only quantum communication



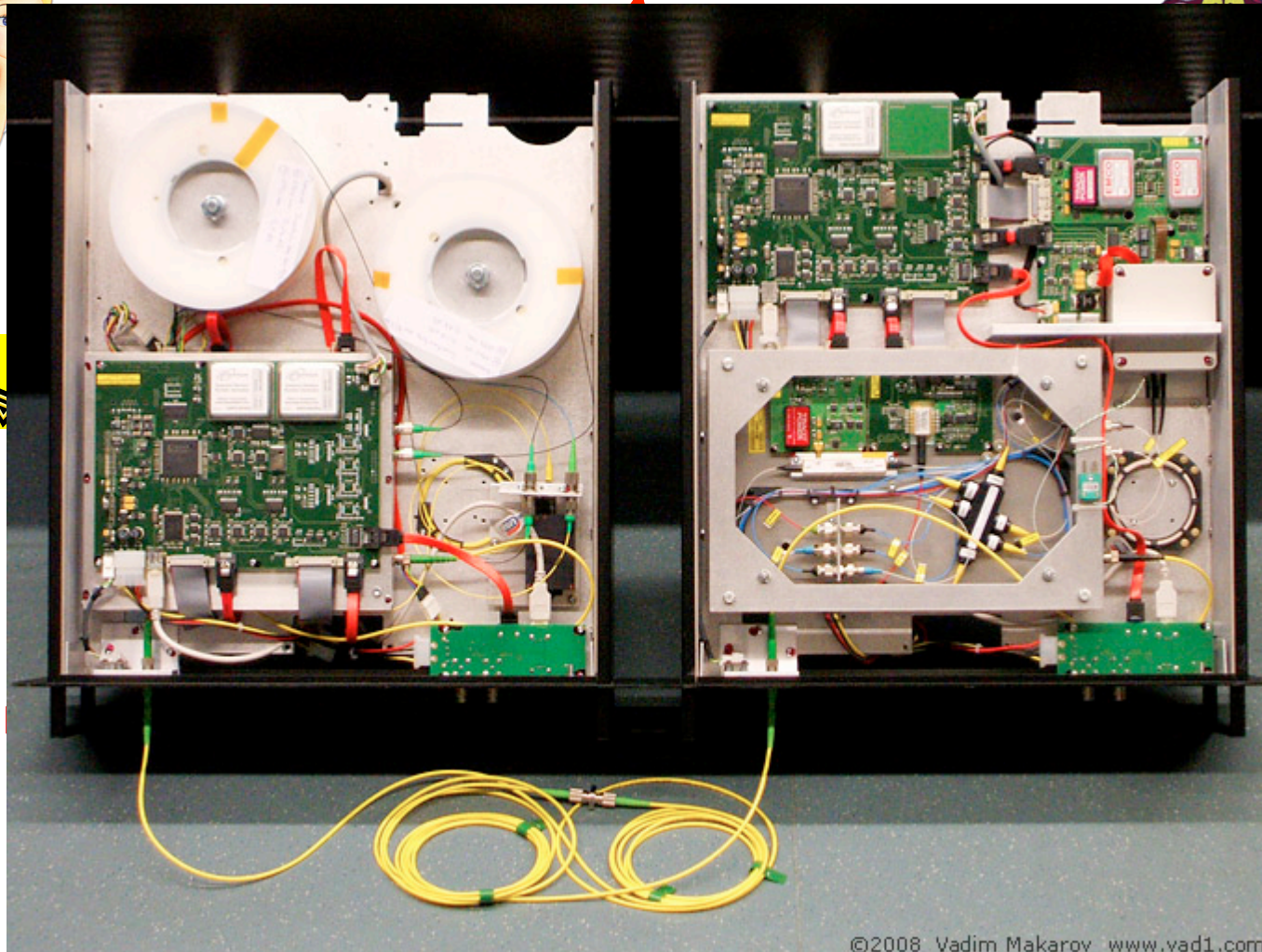
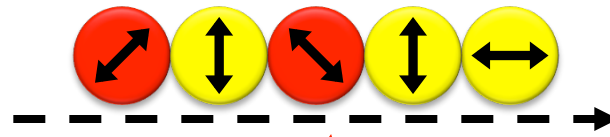
# Quantum Key Distribution (QKD)

12

[Bennett Brassard 84]



Alice



Bob



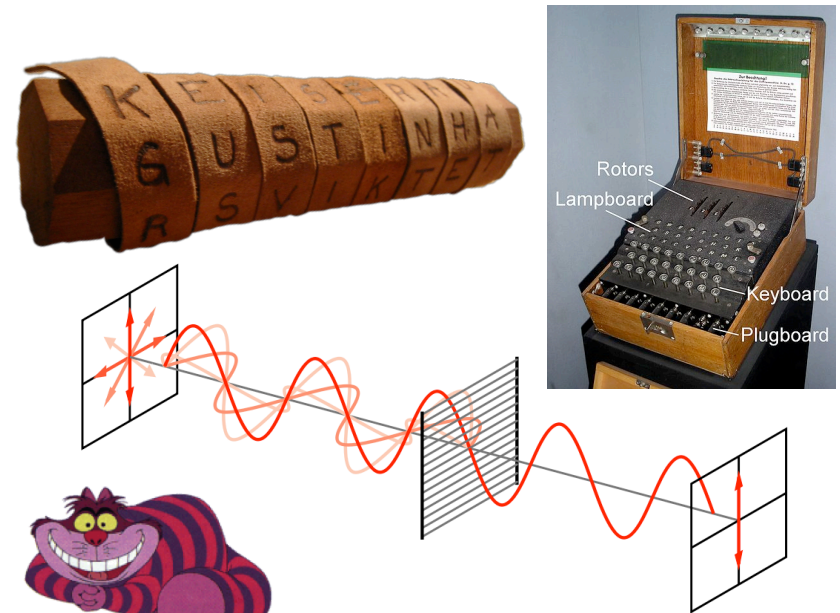
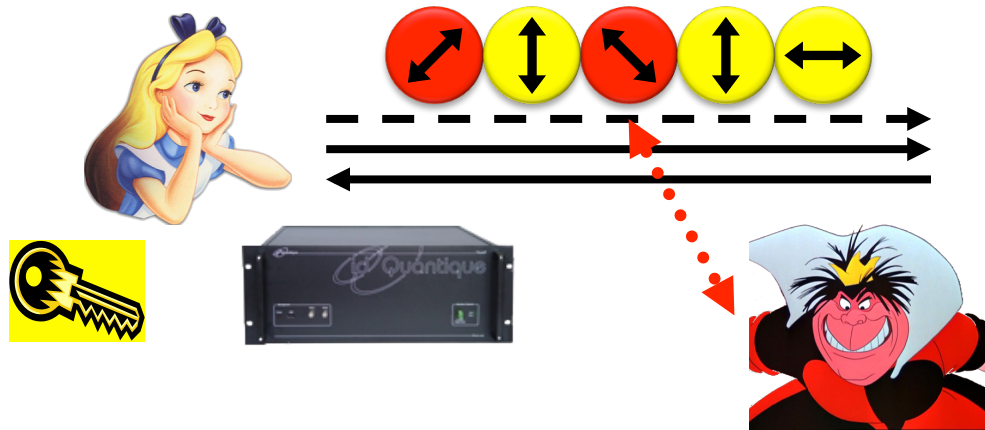
■ tech  
only



# Summary

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- One-Time Pad
- Quantum Key Distribution



Master of Logic course:  
Modern Cryptography

BSc Informatica course:  
Information & Communication

Master of Logic course:  
Quantum Computing

Master of Logic course:  
Information Theory

Master of Logic course:  
Computational Complexity

# Perfect Security

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$$m = ?$$



$$c = m \oplus k = 0101 \ 0100$$

$$m = c \oplus k = ?$$

Alice





$k = ?$





$k = ?$

Eve



Bob



$k = ?$

- Given that
  - is it possible that
    - Yes, if
  - is it possible that
    - Yes, if
  - it is possible that
    - Yes, if
- In fact, every  $m$  is possible.
- Hence, the one-time pad is **perfectly secure!**

$$c = 0101 \ 0100,$$

$$m = 0000 \ 0000 \ ?$$

$$k = 0101 \ 0100.$$

$$m = 1111 \ 1111 \ ?$$

$$k = 1010 \ 1011.$$

$$m = 0101 \ 0101 \ ?$$

$$k = 0000 \ 0001$$

x	y	$x \oplus y$
0	0	0
0	1	1
1	0	1
1	1	0