

Homework 4

1. In the following, assume that x is not a free variable of ψ . Which of the following statements are intuitionistically valid? (If yes, give a proof, if not, give a countermodel).

- (a) $(\exists x\varphi(x) \rightarrow \psi) \rightarrow \forall x(\varphi(x) \rightarrow \psi)$
- (b) $\forall x(\varphi(x) \rightarrow \psi) \rightarrow (\exists x\varphi(x) \rightarrow \psi)$
- (c) $(\forall x\varphi(x) \rightarrow \psi) \rightarrow \exists x(\varphi(x) \rightarrow \psi)$
- (d) $\exists x(\varphi(x) \rightarrow \psi) \rightarrow (\forall x\varphi(x) \rightarrow \psi)$ [4 pts]

2. (a) Show that the following is valid: If $\vdash_{\mathbf{IPC}} (\varphi \rightarrow \psi) \rightarrow (\chi \vee \theta)$, then $\vdash_{\mathbf{IPC}} (\varphi \rightarrow \psi) \rightarrow \chi$ or $\vdash_{\mathbf{IPC}} (\varphi \rightarrow \psi) \rightarrow \theta$ or $\vdash_{\mathbf{IPC}} (\varphi \rightarrow \psi) \rightarrow \varphi$. [4 pts]

- (b) Give an example such that the first two alternatives of (a) do not apply, but the last one does. [2 pts]

- (c)* Let α be $(\neg\phi \rightarrow \psi \vee \chi) \rightarrow (\neg\phi \rightarrow \psi) \vee (\neg\phi \rightarrow \chi)$.

Show that, if $\vdash_{\mathbf{IPC}} \alpha \rightarrow \beta \vee \gamma$, then

$\vdash_{\mathbf{IPC}} \alpha \rightarrow \beta$ or $\vdash_{\mathbf{IPC}} \alpha \rightarrow \gamma$. [2 pts]

3. (a) Let φ contain only \wedge, \vee and \rightarrow but no \neg and no \perp . Let \mathfrak{M} be any Kripke-model (for the language of φ). Extend the model \mathfrak{M} to \mathfrak{M}^+ by adding one more node x at the top above all the nodes of \mathfrak{M} , and making all the propositional variables of φ true in x .

Show that, for all the nodes w in \mathfrak{M} we have:

$$\mathfrak{M}, w \models \varphi \text{ iff } \mathfrak{M}^+, w \models \varphi$$

(satisfaction in the old and new model is the same for φ). [4 pts]

- (b) Let φ contain only \wedge, \vee and \rightarrow but no \neg and no \perp . Show that $\vdash_{\mathbf{IPC}} \varphi$ iff $\vdash_{\mathbf{KC}} \varphi$. (You may use what is claimed about completeness of \mathbf{KC} in class.) [2 pts]