

Stochastic Simulation

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Chapter VI

Rare-Event Simulation

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A: Let's see an example to see whether this is a valid answer.

Rare-Event Simulation

Let's see whether this works... Let $Z = \mathbb{1}_{\{A\}}$, so that indeed $z = \mathbb{E}[Z] = \mathbb{P}(A)$. Z is then Bernoulli distributed with parameter z .

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We know that $\text{Var}[Z] = z(1 - z)$. This means that

$$\frac{\sigma_Z}{z} = \sqrt{\frac{1-z}{z}}$$

which behaves like $z^{-\frac{1}{2}}$ for small z .

Rare-Event Simulation

Recall the 95%-confidence interval for \hat{z}_R :

$$\left(\hat{z}_R - 1.96 \frac{\sigma_Z}{\sqrt{R}}, \hat{z}_R + 1.96 \frac{\sigma_Z}{\sqrt{R}} \right)$$

Suppose that we want to acquire a precision such that the width of our confidence interval is about 20% of the value of z . In other words:

$$0.1 = 1.96 \frac{\sigma_Z}{z\sqrt{R}}$$

or

$$R = \frac{100 \cdot 1.96^2 z(1-z)}{z^2}.$$

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$$R = \frac{100 \cdot 1.96^2 z(1-z)}{z^2}. \quad (1)$$

This number increases like z^{-1} towards ∞ as $z \downarrow 0$. Thus..

A: No, when z becomes small enough, there is always a point at which crude Monte Carlo will fail as R gets too large.

We will need other machinery in this setting. Importance sampling will turn out to be a useful tool.

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Formal setup:

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Quest: find well-performing algorithms such that the required R does not explode. What does this mean?

Recall the expression we just had:

$$R = \frac{100 \cdot 1.96^2 \text{Var}[Z(x)]}{z(x)^2}$$

We wish R to stay finite as $z(x) \downarrow 0$. This happens, when $Z(x)$ has a *bounded relative error*:

$$\limsup_{x \rightarrow \infty} \frac{\text{Var}[Z(x)]}{z(x)^2} < \infty.$$

Rare-Event Simulation

Bounded relative error:

$$\limsup_{x \rightarrow \infty} \frac{\text{Var}[Z(x)]}{z(x)^2} < \infty.$$

In practice, we often check whether a variant of this condition holds, *logarithmic efficiency*:

$$\limsup_{x \rightarrow \infty} \frac{\text{Var}[Z(x)]}{z(x)^{2-\epsilon}} = 0.$$

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Why logarithmic efficiency?

- ▶ The difference is minor from a practical point of view.
- ▶ Logarithmic efficiency is often easier to verify.

Hence, our quest: find logarithmically efficient estimators.

Rare-Event Simulation

Example. Let N be a geometric r.v. with success parameter π , i.e.

$\mathbb{P}(N = n) = \pi(1 - \pi)^{n-1}$, and consider

$z := \mathbb{P}(N \leq m) = 1 - (1 - \pi)^m$.

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For simulating z , we consider importance sampling such that N is instead simulated from a geometric distribution with a success parameter $\tilde{\pi}$, which does not depend on x .

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For simulating z , we consider importance sampling such that N is instead simulated from a geometric distribution with a success parameter $\tilde{\pi}$, which does not depend on x .

A family $\{Z(x)\}$ of estimators is now given by

$$Z(x) = \frac{1}{R} \sum_{i=1}^R \mathbb{1}_{\{N_i \leq m\}} \frac{\pi(x)(1 - \pi(x))^{N_i-1}}{\tilde{\pi}(1 - \tilde{\pi})^{N_i-1}},$$

The book shows that

$$\limsup_{x \rightarrow \infty} \frac{\text{Var}[Z(x)]}{z(x)^2} \leq \limsup_{x \rightarrow \infty} \frac{\mathbb{E}[Z(x)^2]}{z(x)^2} = e - 1.$$

Hence, logarithmically efficient and bounded relative error!

Rare-Event Simulation

We now proceed to the study of sums of light-tailed random variables (i.e. the relevant tails decay at least at an exponential rate).

Recall the idea of exponential tilting: suppose that X_1, \dots, X_n are i.i.d with common density $f(x)$. The importance distribution then preserves the i.i.d. property but changes $f(x)$ to

$g_\theta(x) = \frac{e^{\theta x}}{\mathbb{E}[e^{\theta X}]} f(x)$. Then,

$$L_{n,\theta} = \prod_{i=1}^n \frac{f(X_i)}{g_\theta(X_i)} = e^{-\theta S_n} \hat{F}[\theta]^n = e^{-\theta S_n + n\kappa(\theta)},$$

where $S_n = \sum_{i=1}^n X_i$, $\hat{F}[\theta] = \mathbb{E}[e^{\theta X}]$, $\kappa(\theta) = \log \hat{F}[\theta]$.

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where $S_n = \sum_{i=1}^n X_i$, $\hat{F}[\theta] = \mathbb{E}[e^{\theta X}]$, $\kappa(\theta) = \log \hat{F}[\theta]$. Moreover,

$$\begin{aligned} \mathbb{E}[h(X_1, \dots, X_n)] &= \mathbb{E}_\theta [h(X_1, \dots, X_n) L_{n,\theta}] \\ &= \mathbb{E}_\theta \left[h(X_1, \dots, X_n) e^{-\theta S_n + n\kappa(\theta)} \right]. \end{aligned}$$

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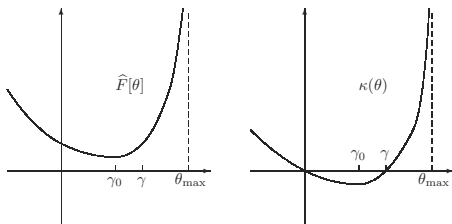
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What do $\hat{F}[\theta]$ and $\kappa(\theta)$ typically look like?

Rare-Event Simulation

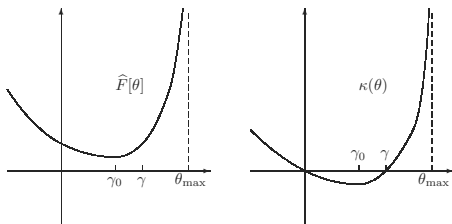
These are $\hat{F}[\theta]$ and $\kappa(\theta)$ for a distribution F with negative mean and $F(0) < 1$:



- ▶ For $\theta \geq \theta_{max}$, $\hat{F}(\theta) = \infty$.
- ▶ γ_0 is the solution of $\hat{F}'[\theta] = \kappa'(\theta) = 0$.
- ▶ γ is the solution to $\hat{F}[\gamma] = 1 - \kappa(\gamma) = 1$.

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We also define the notion of changed drift:

$$\mu_\theta = \mathbb{E}_\theta [X] = \mathbb{E} \left[\frac{Xe^{\theta X}}{\hat{F}[\theta]} \right] = \frac{\hat{F}'[\theta]}{\hat{F}[\theta]} = \kappa'(\theta).$$

Thus, $\mu_\theta < 0$ when $\theta < \gamma_0$, $\mu_{\gamma_0} = 0$ and $\mu_\theta > 0$ when $\theta > \gamma_0$.

Rare-Event Simulation

Suppose that the support of X is not contained in $(-\infty, 0]$, but that still, $\mathbb{E}[X] < 0$.

Now: **Siegmund's algorithm** to consider the problem of estimating, for large x ,

$$z(x) = \mathbb{P}(\tau(x) < \infty), \text{ where } \tau(x) := \inf\{n : S_n > x\}.$$

Idea: use exponential tilting to obtain

$$\begin{aligned} z(x) &= \mathbb{E}[\mathbf{1}_{\{\tau(x) < \infty\}}] = \mathbb{E}_\theta[\mathbf{1}_{\{\tau(x) < \infty\}} L_{\tau(x), \theta}] \\ &= \mathbb{E}_\theta[\mathbf{1}_{\{\tau(x) < \infty\}} e^{-\theta S_{\tau(x)} + \tau(x)\kappa(\theta)}]. \end{aligned}$$

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Q: How to choose θ ?

A: Well, should have at the very least that $\mathbb{P}_\theta(\tau(x) < \infty) = 1$, which happens when $\mathbb{E}_\theta[X] > 0$, or $\theta > \gamma_0$. Then, we have that

$$z(x) = \mathbb{E}_\theta\left[e^{-\theta S_{\tau(x)} + \tau(x)\kappa(\theta)}\right].$$

Rare-Event Simulation

Actually, when a positive solution γ to $\hat{F}[\gamma] = 1 - \kappa(\theta) = 1$ exists, $\theta = \gamma$ turns out to be optimal. Then, we have that

$$z(x) = \mathbb{E}_\theta \left[e^{-\theta S_{\tau(x)} + \tau(x)\kappa(\theta)} \right] = \mathbb{E}_\gamma \left[e^{-\gamma S_{\tau(x)}} \right] = e^{-\gamma x} \mathbb{E}_\gamma \left[e^{-\gamma \xi(x)} \right],$$

where $\xi(x) = S_{\tau(x)} - x$ is the overshoot.

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Theorem: Siegmund's algorithm given by

$$Z(x) = e^{-\gamma x} e^{-\gamma \xi(x)}$$

(simulated in a setting where the distribution of the X 's are exponentially tilted at rate γ) has bounded relative error.

Rare-Event Simulation

Example of this particular change of measure: let $X = U - T$, where U, T are independent exponentially distributed with rate δ, β , where $\beta < \delta$. Then, γ is found by solving

$$1 = \hat{F}[\gamma] = \mathbb{E} \left[e^{\gamma U} \right] \mathbb{E} \left[e^{-\gamma T} \right] = \frac{\delta}{\delta - \gamma} \frac{\beta}{\beta + \gamma}.$$

This leads to $\gamma = \delta - \beta$.

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This leads to $\gamma = \delta - \beta$. Note that

$$\mathbb{E} \left[e^{sX} \right] = \mathbb{E}_{\gamma} \left[e^{sX} L_{1,\gamma} \right] = \mathbb{E}_{\gamma} \left[e^{sX} e^{-\gamma X} \right] = \mathbb{E}_{\gamma} \left[e^{(s-\gamma)X} \right].$$

This implies that

$$\mathbb{E}_{\gamma} \left[e^{sX} \right] = \mathbb{E} \left[e^{(s+\gamma)X} \right] = \mathbb{E} \left[e^{(s+\gamma)U} \right] \mathbb{E} \left[e^{-(s+\gamma)V} \right] = \frac{\beta}{\beta - s} \frac{\delta}{\delta + s}.$$

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This represents the distribution of the difference between two exponential random variables with rates β and δ !

Conclusion: $z(x)$ can be computed or estimated by interchanging parameters. Applications in M/M/1 queue (see book).

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Theorem: Siegmund's algorithm given by

$$Z(x) = e^{-\gamma x} e^{-\gamma \xi(x)}$$

(simulated in a setting where the X 's are exponentially tilted at rate γ) has bounded relative error.

Sketch of proof:

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Sketch of proof:

- ▶ One can reason that under the exponentially tilted measure, $\xi(x)$ is a regenerative process, regenerating each time $S_{\tau(x)}$ has a partial maximum.
- ▶ From this, we can derive that $\{\xi(x), x \geq 0\}$ has a stationary distribution, i.e.

$$\lim_{x \rightarrow \infty} \mathbb{E}_\gamma \left[e^{-\gamma \xi(x)} \right] = \mathbb{E}_\gamma \left[e^{-\gamma \xi(\infty)} \right] =: C$$

so that

$$z(x) = \mathbb{P}(\tau(x) < \infty) \sim C e^{-\gamma x},$$

which is called the *Cramér-Lundberg* approximation.

Rare-Event Simulation

Proof cntd.: We furthermore have that

$$\mathbb{E}_\gamma [Z^2(x)] = e^{-2\gamma x} \mathbb{E}_\gamma [e^{-2\gamma \xi(x)}] \sim C_1 e^{-2\gamma x},$$

where $C_1 = \mathbb{E}_\gamma [e^{-2\gamma \xi(x)}]$. Thus,

$$\text{Var}_\gamma [Z(x)] \sim C_1 e^{-2\gamma x} - C^2 (e^{-2\gamma x}) = C_2 e^{-2\gamma x}$$

with $C_2 = C_1 - C^2$. Thus, we have that the relative error satisfies

$$\frac{\text{Var}_\gamma [Z(x)]}{z(x)} \sim \frac{C_2 e^{-2\gamma x}}{C^2 e^{-2\gamma x}} = C_3 < \infty.$$

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A: Yes, actually, it's the only one that even admits logarithmic efficiency!

To see this, suppose that we regard another importance sampling density $\tilde{h}(x)$ such that $\mathbb{E}_{\tilde{h}}[X] > 0$ (why?). Then, we would have

that $z(x) = \mathbb{E}[\mathbb{1}_{\{\tau(x) < \infty\}}] = \mathbb{E}_{\tilde{h}}[L_{\tau(x)}(f|\tilde{h})]$, so that

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Theorem: This importance sampling algorithm is logarithmically efficient if and only if $\tilde{h}(x) = g_\gamma(x)$.

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that $z(x) = \mathbb{E}[\mathbb{1}_{\{\tau(x) < \infty\}}] = \mathbb{E}_{\tilde{h}}[L_{\tau(x)}(f|\tilde{h})]$, so that

$$Z(x) := L_{\tau(x)}(f|\tilde{h}) = \prod_{i=1}^{\tau(x)} \frac{f(X_i)}{\tilde{h}(X_i)}.$$

Theorem: This importance sampling algorithm is logarithmically efficient if and only if $\tilde{h}(x) = g_\gamma(x)$.

Proof: We have already shown sufficiency.

Rare-Event Simulation

Q: Is $g_\gamma(x)$ the only importance sampling density which admits a bounded relative error in Siegmund's algorithm?

A: Yes, actually, it's the only one that even admits logarithmic efficiency!

To see this, suppose that we regard another importance sampling density $\tilde{h}(x)$ such that $\mathbb{E}_{\tilde{h}}[X] > 0$ (why?). Then, we would have

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Theorem: This importance sampling algorithm is logarithmically efficient if and only if $\tilde{h}(x) = g_{\gamma(x)}$.

Proof: We have already shown sufficiency. So we assume now that $\tilde{h}(x) \neq g_{\gamma(x)}$ and show that the algorithm can not be logarithmically efficient in that case.

Rare-Event Simulation

Suppose $\tilde{h}(x) \neq g_{\gamma(x)}$. Then,

$$\begin{aligned}\mathbb{E}_{\tilde{h}} [Z^2(x)] &= \mathbb{E}_{\tilde{h}} \left[L_{\tau(x)}^2(f|\tilde{h}) \right] = \mathbb{E}_{\tilde{h}} \left[L_{\tau(x)}^2(f|g_{\gamma}) L_{\tau(x)}^2(g_{\gamma}|\tilde{h}) \right] \\ &= \mathbb{E}_{\gamma} \left[L_{\tau(x)}^2(f|g_{\gamma}) L_{\tau(x)}(g_{\gamma}|\tilde{h}) \right] = \mathbb{E}_{\gamma} \left[e^{\sum_{i=1}^{\tau(x)} \kappa_i} \right],\end{aligned}$$

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where

$$K_i = \log \left(\left(\frac{f(X_i)}{g_{\gamma}(X_i)} \right)^2 \frac{g_{\gamma}(X_i)}{\tilde{h}(X_i)} \right) = -2\gamma X_i - \log \left(\frac{\tilde{h}(X_i)}{g_{\gamma}(X_i)} \right).$$

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Further,

$$\mathbb{E}_{\gamma} [K_i] = -2\gamma \mathbb{E}_{\gamma} [X_i] + \epsilon',$$

where $\mathbb{E}_{\gamma} [X_i] > 0$ and $\epsilon' = -\mathbb{E}_{\gamma} \left[\log \left(\frac{\tilde{h}(X_1)}{g_{\gamma}(X_1)} \right) \right]$. Further, we can prove that $\epsilon' > 0$.

Rare-Event Simulation

Further, due to Jensen's inequality and Wald's equality,

$$\begin{aligned}\mathbb{E}_{\tilde{h}} [Z^2(x)] &= \mathbb{E}_{\gamma} \left[e^{\sum_{i=1}^{\tau(x)} K_i} \right] \\ &\geq e^{\mathbb{E}_{\gamma} \left[\sum_{i=1}^{\tau(x)} K_i \right]} \\ &= e^{\mathbb{E}_{\gamma}[\tau(x)](\epsilon' - 2\gamma\mathbb{E}_{\gamma}[X_i])}\end{aligned}$$

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Note that $\mathbb{E}_{\gamma} \left[\frac{\tau(x)}{x} \right] \rightarrow \frac{1}{\mathbb{E}_{\gamma}[X_1]}$ a.s. as $x \rightarrow \infty$. Thus, using $z(x) \sim Ce^{-\gamma x}$, we have that for $0 < \epsilon < \frac{\epsilon'}{\gamma\mathbb{E}_{\gamma}[X_1]}$,

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Conclusion: when $\tilde{h} \neq g_{\gamma}$, we will have no logarithmic efficiency... theorem proved!

Rare-Event Simulation

Well, not entirely just yet. In passing, we assumed that

$$\epsilon' = -\mathbb{E}_\gamma \left[\log \left(\frac{\tilde{h}(X_1)}{g_\gamma(X_1)} \right) \right] > 0.$$

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$$\epsilon' = -\mathbb{E}_\gamma \left[\log \left(\frac{\tilde{h}(X_1)}{g_\gamma(X_1)} \right) \right] > 0.$$

To see this, note that, again due to Jensen's inequality

$$\begin{aligned} \mathbb{E}_\gamma \left[\log \left(\frac{\tilde{h}(X_1)}{g_\gamma(X_1)} \right) \right] &< \log \mathbb{E}_\gamma \left[\frac{\tilde{h}(X_1)}{g_\gamma(X_1)} \right] \\ &= \log \int_{x: g_\gamma(x) > 0} \frac{\tilde{h}(x)}{g_\gamma(x)} g_\gamma(x) dx \\ &= \log \int_{x: g_\gamma(x) > 0} \tilde{h}(x) dx \\ &\leq \log 1 = 0. \end{aligned}$$

Rare-Event Simulation

Another problem. Consider again a random walk $S_n := \sum_{i=1}^n X_i$, where the summands are independently distributed with common distribution F and mean μ (sign unimportant this time).

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We are now interested in the family of rare events

$$A(n) = \{S_n > n(\mu + \delta)\},$$

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We are now interested in the family of rare events

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where $\delta > 0$.

Due to the weak law of large numbers we have that

$$z(n) = \mathbb{P}(A(n)) \rightarrow 0$$

The main question: how to simulate the value of $z(n)$ efficiently as n gets large?

Rare-Event Simulation

To efficiently estimate $z(n) = \mathbb{P}(A(n))$, we again employ exponential change of measure, so that our algorithm becomes

$$Z(n) = e^{-\theta S_n + n\kappa(\theta)} \mathbb{1}_{\{S_n > n(\mu + \delta)\}},$$

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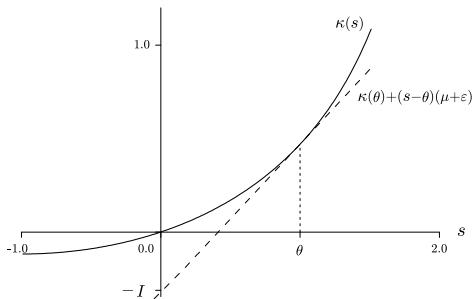
where the X_i are sampled from the tilted distribution. It turns out that θ should be chosen such that

$$\mathbb{E}_\theta [X] = \kappa'(\theta) = \mu + \delta.$$

Rare-Event Simulation

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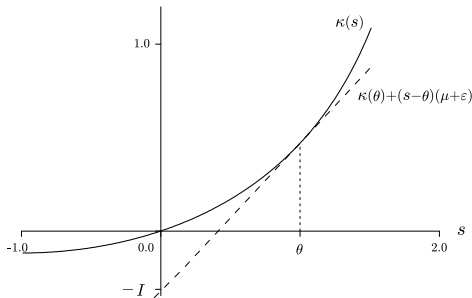
Under this constraint, $\theta > 0$, since $\kappa'(0) = \mu$ and κ is a convex function. Furthermore, we have that $I = \theta(\mu + \delta) - \kappa(\theta) > 0$.



Rare-Event Simulation

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Theorem: The algorithm using the exponentially (θ) tilted measure is logarithmically efficient, and this measure is the only one to make the algorithm logarithmically efficient.

Rare-Event Simulation

Proof: Structure similar to Siegmund's algorithm.

First, note that, since $\theta > 0$

$$\begin{aligned}z(n) &= \mathbb{P}(A(n)) = \mathbb{E}_\theta \left[e^{-\theta S_n + n\kappa(\theta)} \mathbf{1}_{\{S_n > n(\mu + \delta)\}} \right] \\&= e^{-nI} \mathbb{E}_\theta \left[e^{-\theta(S_n - n(\mu + \delta))} \mathbf{1}_{\{S_n > n(\mu + \delta)\}} \right] \\&\leq e^{-nI}\end{aligned}$$

Furthermore,

$$\begin{aligned}\text{Var}_\theta [Z(n)] &\leq \mathbb{E}_\theta [Z^2(n)] = \mathbb{E}_\theta \left[e^{-2\theta S_n + 2n\kappa(\theta)} \mathbf{1}_{\{S_n > n(\mu + \delta)\}} \right] \\&= e^{-2nI} \mathbb{E}_\theta \left[e^{-2\theta(S_n - n(\mu + \delta))} \mathbf{1}_{\{S_n > n(\mu + \delta)\}} \right] \\&\leq e^{-2nI}.\end{aligned}$$

Rare-Event Simulation

Thus, to establish logarithmic efficiency, we will require a lower bound of some sort on the $z(n)$.

Note that, when the density of X_n is exponentially tilted with parameter θ ,

$$\frac{S_n - n(\mu + \delta)}{\sqrt{n}} \rightarrow \mathcal{N}(0, \sigma_\theta^2) \text{ and}$$

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Thus,

$$\begin{aligned} \liminf_{n \rightarrow \infty} (e^{nI + \theta\sqrt{n}} z(n)) &= \liminf_{n \rightarrow \infty} e^{\theta\sqrt{n}} \mathbb{E}_\theta \left[e^{-\theta(S_n - n(\mu + \delta))} \mathbf{1}_{\{S_n > n(\mu + \delta)\}} \right] \\ &\geq \liminf_{n \rightarrow \infty} e^{\theta\sqrt{n}} \mathbb{E}_\theta \left[e^{-\theta(S_n - n(\mu + \delta))} \mathbf{1}_{\left\{ \frac{S_n - n(\mu + \delta)}{\sqrt{n}} \in (0, 1) \right\}} \right] \\ &\geq \liminf_{n \rightarrow \infty} e^{\theta\sqrt{n}} e^{-\theta\sqrt{n}} \mathbb{P}_\theta \left(\mathbf{1}_{\left\{ \frac{S_n - n(\mu + \delta)}{\sqrt{n}} \in (0, 1) \right\}} \right) = c > 0. \end{aligned}$$

Rare-Event Simulation

Thus, to show logarithmic efficiency, note that

$$\begin{aligned}\limsup_{n \rightarrow \infty} \frac{\text{Var}_\theta [Z(n)]}{z(n)^{2-\epsilon}} &= \limsup_{n \rightarrow \infty} \frac{\text{Var}_\theta [Z(n)] e^{(2-\epsilon)(nl+\theta\sqrt{n})}}{z(n)^{2-\epsilon} e^{(2-\epsilon)(nl+\theta\sqrt{n})}} \\ &\leq \frac{\limsup_{n \rightarrow \infty} e^{-\epsilon nl + (2-\epsilon)\theta\sqrt{n}}}{c^{2-\epsilon}} = 0.\end{aligned}$$

To show that there is no other density which allows logarithmic efficiency, a proof using similar arguments as the one for Siegmund's algorithm can be given.