### COLOURING GRAPHS OF BOUNDED LOCAL DENSITY

Ross J. Kang\*



Radboud University Nijmegen

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<sup>\*</sup>With Davies, Hurley, de Joannis de Verclos, Pirot, Sereni. Support from NWO.



#### VIZING'S PROBLEM

#### Vizing, "Some unsolved problems in graph theory" (1968):

Можно ли для любого наперед заданного натурального числа  $k \ge 2$ построить граф со сколь угодно большим обхватом и с хроматическим числом k? Можно. Это доказал П. Эрдёш [39], основываясь на мощностных соображениях. Удивительно, что до сих пор нет конструктивного доказательства этого факта. В [40] указан способ построения графов с любым хроматическим числом без циклов длины  $\leqslant 7$ . Это лучшее, что мы имеем на сегодняшний день.

Если  $\sigma(L)$  — максимальная степень вершины графа L, то, очевидно,  $\gamma(L) \leqslant \sigma(L) + 1$ . В 1941 г. Р. Брукс [41] доказал, что при  $\sigma(L) \geqslant 3$ и  $\omega(L) \leqslant \sigma(L)$  справедлива оценка  $\gamma(L) \leqslant \sigma(L)$ . Дальнейшие исследования можно проводить, учитывая более точно соотношения между  $\sigma$  и  $\omega$ . Пожалуй, следует начать с оценки хроматического числа графа без треугольников ( $\omega = 2$ ) с данной максимальной степенью вершины.

#### VIZING'S PROBLEM

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Given a natural number  $k \ge 2$ , is it possible to construct a graph with arbitrarily large compass and with chromatic number k? Erdős [39] has oroved this; his proof is based on counting arguments. It is astonishing that no constructive proof for this fact has yet been given. In [40] a method is given of constructing graphs of arbitrary chromatic number without having cycles of length < 7. This is the best we have at the present time.

If  $\sigma(L)$  is the maximum degree of a vertex in a graph L, it is clear that  $\gamma(L) < \sigma(L) + 1$ . Brooks [41] showed in 1941 that  $\gamma(L) \leq \sigma(L)$  whenever  $\sigma(L) > 3$  and  $\omega(L) < \sigma(L)$ . Further investigations could be conducted, taking into account a more exact relation between  $\sigma$  and  $\omega$ . Perhaps one should start with estimates of the chromatic number of a graph without triangles ( $\omega = 2$ ) and with given maximal degree for vertices.

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Distinguished origins:

- Mantel (1907), Turán (1941)
- Ramsey (1930), Erdős & Szekeres (1935)
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Elegant, modern challenges!

 $\omega(G)$ 

$$\omega(G) \leq \max_{\emptyset \neq H \subseteq G} \frac{|H|}{\alpha(H)}$$

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In general, each inequality can be strict<sup>†</sup>

<sup>&</sup>lt;sup>†</sup>On strictness of the second, see Blumenthal, Lidický, Martin, Norin, Pfender, Volec (2018+), and Dvořák, Ossona de Mendez, Wu (2020); nice open question in the triangle-free case.

### PROBABILISTIC METHOD



If random object has property with positive probability, then there exists *at least one* object with that property













TRIANGLE-FREE GRAPHS

# Off-diagonal Ramsey $\operatorname{numbers}^{\ddagger}$



<sup>‡</sup>Picture credit: Soifer 2009

### Off-diagonal Ramsey numbers

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Theorem (Bollobás 1981)  $\max_{\emptyset \neq H \subseteq G} \frac{|H|}{\alpha(H)} \geq \frac{\Delta}{2 \log \Delta} \text{ and } G \text{ has arbitrarily large girth whp for } G \sim G_{n,\Delta}$ 

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$$\omega(G) \leq \max_{\emptyset \neq H \subseteq G} \frac{|H|}{\alpha(H)} \leq \chi_f(G) \leq \chi(G) \leq \mathsf{ch}(G) \leq \Delta(G) + 1$$

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 $\begin{array}{ll} \mbox{Theorem (} & \mbox{Johansson 1996+)} \\ \mbox{ch}(G) = O\left(\frac{\Delta}{\log\Delta}\right) \mbox{ for any triangle-free $G$ of maximum degree $\Delta$} \end{array}$ 

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Theorem (Davies, de Joannis de Verclos, Kang, Pirot 2021)  $\chi_f(G) \lesssim \frac{\Delta}{\log \Delta}$  for any triangle-free G of maximum degree  $\Delta \rightsquigarrow$ 

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## Two more Ramsey-type problems

$$\omega(G) \leq \max_{\substack{0 \neq H \subseteq G}} \frac{|H|}{\alpha(H)} \leq \chi_f(G) \leq \chi(G) \leq ch(G) \leq \Delta(G) + 1$$

<sup>§</sup>Picture credit: Wikipedia/David Eppstein

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- $\omega = 2$ ,  $\Delta$  large enough (Johansson 1996+)
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Chvátal graph (1970)§ has  $\omega=$  2,  $\Delta=$  4,  $\chi=$  4

Bound holds for:

- $\omega = 2$ ,  $\Delta$  large enough (Johansson 1996+) ( $\omega \leq \Delta^{1/100}$ ,  $\Delta$  large enough (Davies, Kang, Pirot, Sereni 2020+))
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#### Corollary

 $\chi(G) \leq [\varepsilon_2 \omega + (1 - \varepsilon_2)(\Delta + 1)]$  for any G of clique number  $\omega$  and maximum degree  $\Delta$  for some  $\varepsilon_2 > 0$  provided  $\Delta$  is large enough

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 $\chi(G) \leq \lceil \varepsilon_2 \omega + (1 - \varepsilon_2)(\Delta + 1) \rceil$  for any G of clique number  $\omega$  and maximum degree  $\Delta$  for some  $\varepsilon_2 > 0$  provided  $\Delta$  is large enough

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NB:  $\varepsilon_2$  may not be larger than 0.5

two edges "close": incident or connected by an edge induced matching: edge subset with no two edges close

strong edge-colouring of G: partition of the edges into induced matchings strong chromatic index  $\chi(L(G)^2)$  of G: least number of parts needed

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Conjecture (Erdős & Nešetřil 1985)  $\chi(L(G)^2) \leq \frac{5}{4}\Delta^2$  for any G of maximum degree  $\Delta$ 

So far only confirmed for  $\Delta \leq 3$ (Andersen 1992, Horák, He, Trotter 1993)



 $\chi(L(G)^2) < 2\Delta^2$  since  $\Delta(L(G)^2) \le 2\Delta(\Delta-1)$  trivially

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Theorem (Molloy & Reed 1997)  $\chi(L(G)^2) \leq (2 - \varepsilon_3)\Delta^2$  for any G of maximum degree  $\Delta$  for some  $\varepsilon_3 > 0$  provided  $\Delta$  is large enough

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• Molloy & Reed (1997):  $\varepsilon_3 > 0.001$ • Bruhn & Joos (2018):  $\varepsilon_3 > 0.070$ 

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NB:  $\varepsilon_3 (\geq \varepsilon_4)$  may not be larger than 0.75

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G of maximum degree  $\Delta$ has local density at most  $\eta$  or, alternatively, is (locally)  $(1 - \eta)$ -sparse if  $\leq \eta {\Delta \choose 2}$  edges per neighbourhood no edge is induced in any neighbourhood  $\iff$  triangle-free

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 $\eta < 1/{\Delta \choose 2}$  corresponds to triangle-free,  $\eta = 1$  corresponds to general case

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Best chromatic number bounds in graphs of local density at most  $\eta$ ?

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Theorem (Alon, Krivelevich, Sudakov 1999, cf. Vu 2002, Achlioptas, lliopoulos, Sinclair 2019)

 $\chi(G) = O\left(\frac{\Delta}{\log \frac{e}{\sqrt{\eta}}}\right) \text{ for any } G \text{ of maximum degree } \Delta$ and local density at most  $\eta$ ,  $\frac{1}{\Delta^2} \le \eta \le 1$  Best chromatic number bounds in graphs of local density at most  $\eta$ ?

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#### Theorem (Molloy & Reed 1997)

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- Hurley, de Joannis de Verclos, Kang (2021):
  - random vertex ordering to decide in 2
    - allows better iteration of 1 and 2

## NAÏVE INDEPENDENT SETS

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#### Lemma

For  $\iota > 0$ , generating I as above from a  $\sigma$ -sparse  $\Delta$ -regular graph satisfies for large enough  $\Delta$  and  $\gamma$ , and for every v

$$\left| \mathbb{P}(\mathbf{v} \in \mathbf{I}) - \frac{1 - e^{-\gamma}}{\Delta} \right| \le \frac{2}{\Delta^2},$$
$$\frac{\mathbb{P}(\mathbf{N}(\mathbf{v}) \cap \mathbf{I} \neq \emptyset)}{\mathbb{E}|\mathbf{N}(\mathbf{v}) \cap \mathbf{I}|} \le 1 - \frac{\sigma}{2} + \frac{\sigma^{3/2}}{6} + \iota$$

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