

# ON CONJECTURES OF REED

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Reedfest

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## BROOKS' THEOREM

Theorem (Brooks 1941)

$\chi(G) \leq \Delta(G)$  unless  $G = K_{\Delta(G)+1}$  or  $G$  is an odd cycle.

## VIZING'S PROBLEM

Vizing, "Some unsolved problems in graph theory" (1968):

Можно ли для любого наперед заданного натурального числа  $k \geq 2$  построить граф со сколь угодно большим обхватом и с хроматическим числом  $k$ ? Можно. Это доказал П. Эрдёш [39], основываясь на мощностных соображениях. Удивительно, что до сих пор нет конструктивного доказательства этого факта. В [40] указан способ построения графов с любым хроматическим числом без циклов длины  $\leq 7$ . Это лучшее, что мы имеем на сегодняшний день.

Если  $\sigma(L)$  — максимальная степень вершины графа  $L$ , то, очевидно,  $\gamma(L) \leq \sigma(L) + 1$ . В 1941 г. Р. Брукс [41] доказал, что при  $\sigma(L) \geq 3$  и  $\omega(L) \leq \sigma(L)$  справедлива оценка  $\gamma(L) \leq \sigma(L)$ . Дальнейшие исследования можно проводить, учитывая более точно соотношения между  $\sigma$  и  $\omega$ . Пожалуй, следует начать с оценки хроматического числа графа без треугольников ( $\omega = 2$ ) с данной максимальной степенью вершины.

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Given a natural number  $k \geq 2$ , is it possible to construct a graph with arbitrarily large compass and with chromatic number  $k$ ? Erdős [39] has proved this; his proof is based on counting arguments. It is astonishing that no constructive proof for this fact has yet been given. In [40] a method is given of constructing graphs of arbitrary chromatic number without having cycles of length  $\leq 7$ . This is the best we have at the present time.

If  $\sigma(L)$  is the maximum degree of a vertex in a graph  $L$ , it is clear that  $\chi(L) \leq \sigma(L) + 1$ . Brooks [41] showed in 1941 that  $\chi(L) \leq \sigma(L)$  whenever  $\sigma(L) \geq 3$  and  $\omega(L) \leq \sigma(L)$ . Further investigations could be conducted, taking into account a more exact relation between  $\sigma$  and  $\omega$ . Perhaps one should start with estimates of the chromatic number of a graph without triangles ( $\omega = 2$ ) and with given maximal degree for vertices.

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$$\text{where } \rho = \max_{\emptyset \neq H \subseteq G} \frac{|H|}{\alpha(H)}$$

(upper bounds on  $\rho$  are like lower bounds on  $\alpha$ )



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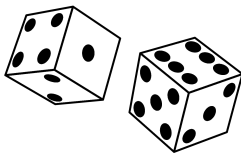
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$$\chi \leq \lceil \tfrac{1}{2}(\omega + \Delta + 1) \rceil?$$

## PROBABILISTIC METHOD



If random object has property with positive probability,  
then there exists at least one object with that property





## CHROMATIC NUMBER OF TRIANGLE-FREE GRAPHS

Further investigations could be conducted, taking into account a more exact relationship between  $\Delta$  and  $\omega$ . Perhaps one should start with estimates of the chromatic number of a graph without triangles ( $\omega = 2$ ) and with given maximal degree for vertices.

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Theorem (Shearer 1983, cf. Ajtai, Komlós, Szemerédi 1980/1)

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**NB:** these bounds are sharp up to a factor 2 by random  $\Delta$ -regular graphs.

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Bounds of Shearer (1995) and Johansson (1996+) are out by a  $\log \log \Delta$  factor.

## THE $\omega$ , $\Delta$ , $\chi$ CONJECTURE

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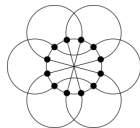
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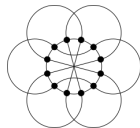
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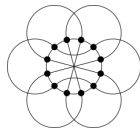
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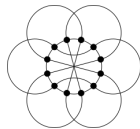
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(for  $\omega \leq \Delta^{1/101}$ ,  $\Delta$  large enough (Davies, Kang, Pirot, Sereni 2020+))
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$\eta < 1/\binom{\Delta}{2}$  means triangle-free,  $\eta = 1$  means unrestricted

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Theorem (Alon, Krivelevich, Sudakov 1999, cf. Vu 2002, Achlioptas, Iliopoulos, Sinclair 2019)

$$\chi(G) = O\left(\frac{\Delta}{\log \frac{e}{\sqrt{\eta}}}\right) \text{ for any } G \text{ of maximum degree } \Delta$$

*and local density at most  $\eta$ ,  $\frac{1}{\Delta^2} \leq \eta \leq 1$*

Theorem (Davies, Kang, Pirot, Sereni 2020+)

$$\chi(G) \lesssim \frac{\Delta}{\log \frac{e}{\sqrt{\eta}}} \text{ for any } G \text{ of maximum degree } \Delta$$

*and local density at most  $\eta$ ,  $\frac{1}{\Delta^2} \leq \eta \ll 1$*

$\eta = \frac{1}{\Delta^2}$  matches Molloy's; the bound is sharp up to a factor of between 2 and 4

**NB** :  $\eta = 1$  should match  $\Delta + 1$  bound, but neither gives this...

## CHROMATIC NUMBER UNDER BOUNDED LOCAL DENSITY

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Lower bounds on  $\varepsilon$  key to sharpest results towards Reed's conjecture  
(and sharpest results towards the Erdős-Nešetřil conjecture)

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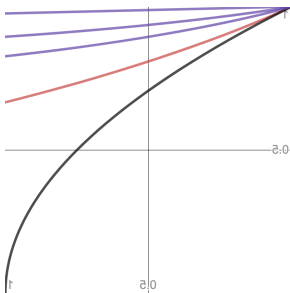
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Theorem (Hurley, de Joannis de Verclos, Kang 2021)

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    - random vertex ordering to decide in ??
    - allows better iteration of ?? and ??

## LINK UP?

Theorem (Davies, Kang, Piro, Sereni 2020+)

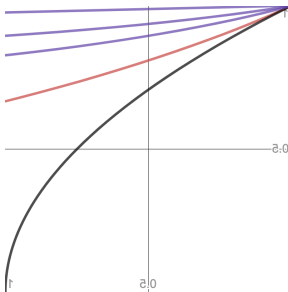
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## BEYOND TRIVIAL

$$\omega(G) \leq \rho(G) \leq \chi_f(G) \leq \chi(G) \leq \chi_\ell(G) \leq \Delta(G) + 1$$

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The list chromatic number  $\chi_\ell(G)$  is least  $k$  such that  $G$  is  $k$ -choosable.

$\chi_\ell(G) \geq \chi(G)$  by considering constant  $L$ .



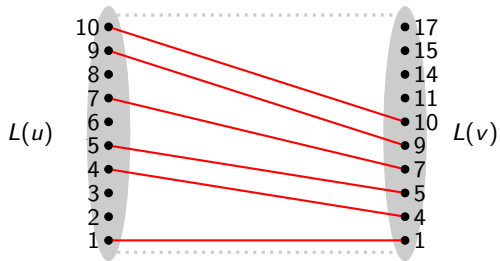
## COVER GRAPHS



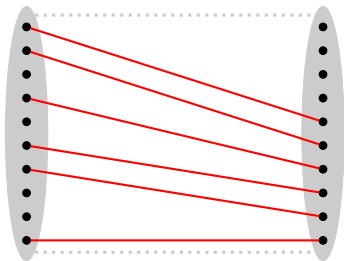
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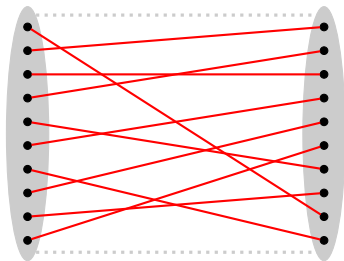
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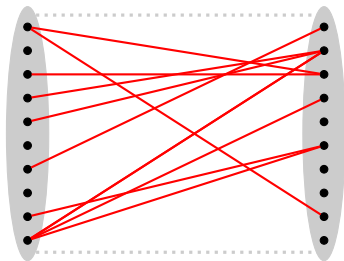
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BESz knew hypothetically the sharpness of the factor 2, but a construction of Szabó & Tardos (2005) shows moreover the statement is exactly sharp!

**Theorem (Szabó & Tardos 2005)**

*There is a graph of maximum degree  $D$  with a partition into parts of size  $2D - 1$  that does not admit an IT.*

## BRUCE DIDN'T WANT TO GO SO FAR

Trivial bound ( $\chi_\ell(G) \leq \Delta(G) + 1$ ) redux:

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Bohman & Holzman (2002): counterexample!

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Still open: could  $f(D) = D + C$  suffice?

## AND A LITTLE BIT FURTHER?

Trivial bound ( $\chi_c(G) \leq \Delta(G) + 1$ ) redux redux:

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Proposition (Dvořák, Esperet, Kang, Ozeki 2021, cf. Wanless & Wood 2022)

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NB: incidentally, this confirms a conjecture of Erdős, Gyárfás, Łuczak (1994).

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Loh–Sudakov (2007) and Kang–Kelly (2022) conjectures



## A CROSSING OF THE TWO ARCS

### Theorem (Amini & Reed 2008)

*There is  $C > 0$  such that, if  $H$  is a  $(C \cdot D / \log D)$ -fold list-cover,  $H$  is triangle-free, and  $H$  has maximum degree  $D$ , then  $H$  admits an IT.*

**NB:** Alon & Assadi (2020+) independently derived a very similar result with the aim of ‘palette sparsification’.



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### Theorem (Amini & Reed 2008)

*There is  $C > 0$  such that, if  $H$  is a  $(C \cdot D / \log D)$ -fold list-cover,  $H$  is triangle-free, and  $H$  has maximum degree  $D$ , then  $H$  admits an IT.*

**NB:** Alon & Assadi (2020+) independently derived a very similar result with the aim of 'palette sparsification'.

### Conjecture (Cambie & Kang 2020+)

*There is  $C > 0$  such that, if  $H$  is a  $(C \cdot D / \log D)$ -fold correspondence-cover,  $H$  is triangle-free, and  $H$  has maximum degree  $D$ , then  $H$  admits an IT.*

**NB:** Anderson, Bernshteyn, Dhawan (2021+) propose an even stronger form.



QUESTIONS?

