

recent paper by Pascal Koiran

Tensor rank and commuting matrices

<https://arxiv.org/abs/2006.02374v1>

Matrices generally don't commute

$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \quad \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}$$

$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \cdot \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

≠

$$\begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix} \cdot \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} = \begin{pmatrix} -1 & 0 \\ 0 & 1 \end{pmatrix}$$

But we may extend them to make them commute

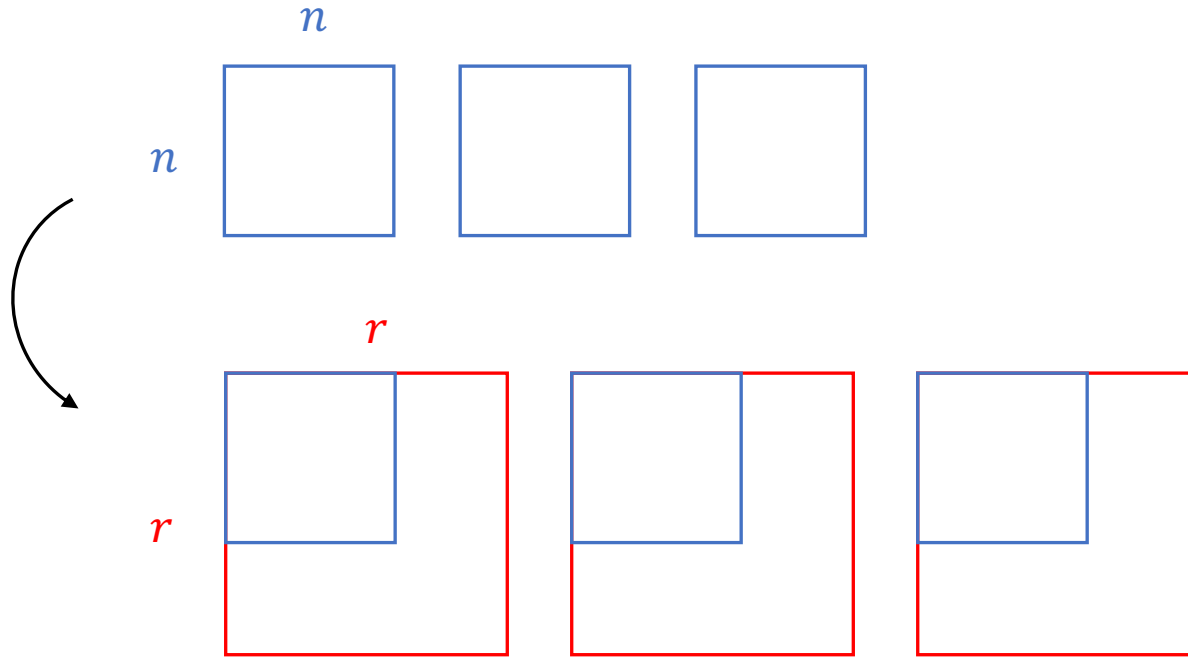
$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

$$\begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}$$

$$\begin{pmatrix} 0 & 1 & -1 \\ 1 & 0 & 2 \\ 2 & -1 & 3 \end{pmatrix}$$

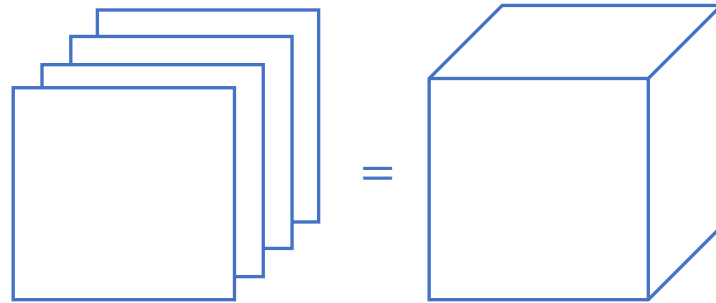
$$\begin{pmatrix} 0 & -1 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & -1 \end{pmatrix}$$

Problem 1: Extend matrices until they commute

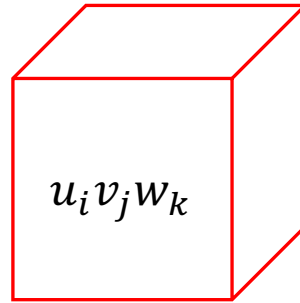


Commuting extension

Tensor

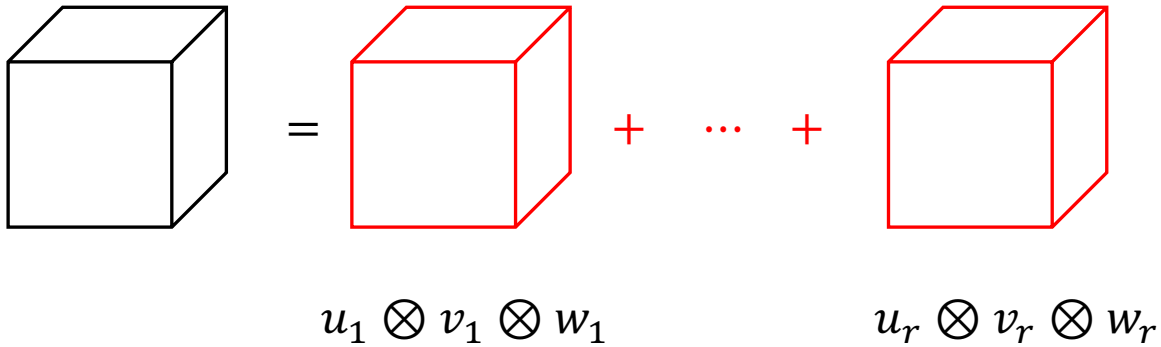


A simple tensor is the outer product of three vectors



$$u \otimes v \otimes w$$

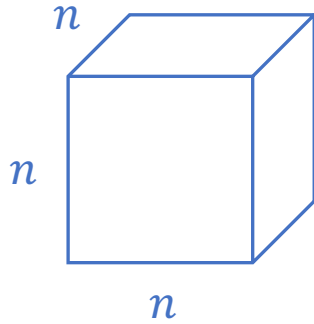
Smallest number of simple tensors that sum to T is the tensor rank



The diagram illustrates the decomposition of a tensor T into a sum of simple tensors. On the left, a black cube represents the tensor T . This is followed by an equals sign. To the right of the equals sign, there is a sequence of red cubes. The first red cube is labeled $u_1 \otimes v_1 \otimes w_1$ below it. This is followed by a plus sign, an ellipsis, another plus sign, and a final red cube labeled $u_r \otimes v_r \otimes w_r$ below it.

$$T = u_1 \otimes v_1 \otimes w_1 + \dots + u_r \otimes v_r \otimes w_r$$

Problem 2: Random tensors have high rank; find explicit ones



random tensor

$$= c n^2$$

explicit tensor

$$\geq d n$$

tensor rank

Explicit tensors of high rank, in complexity theory:

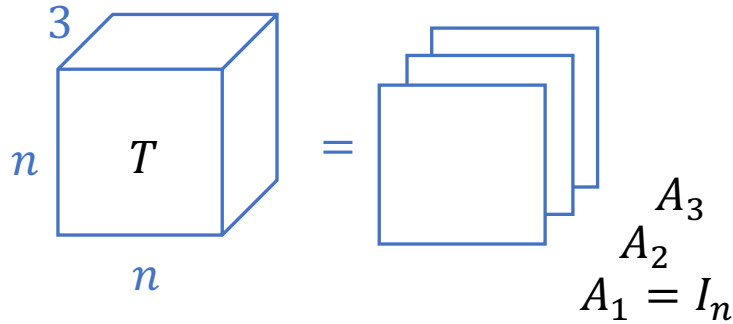
1. Explicit tensors of high rank imply lower bounds on the size of **arithmetic formulas** [Raz]
2. Tools to prove explicit tensor rank lower bounds may prove non-trivial lower bounds on the exponent of **matrix multiplication**

$$\omega > 2$$

Inspiration: Strassen lower bound on tensor rank

Theorem 1. (Strassen)

Given:



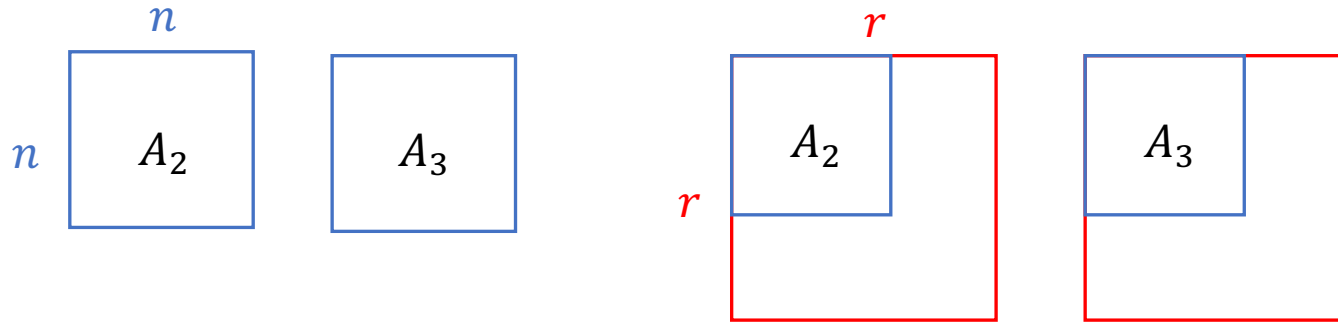
Then:

$$\text{rank}(T) \geq n + \frac{1}{2} \text{rank}(A_2 A_3 - A_3 A_2)$$

Koiran interpretation: Lower bound on commuting extensions

Theorem 2. (Strassen, Koiran)

Given:



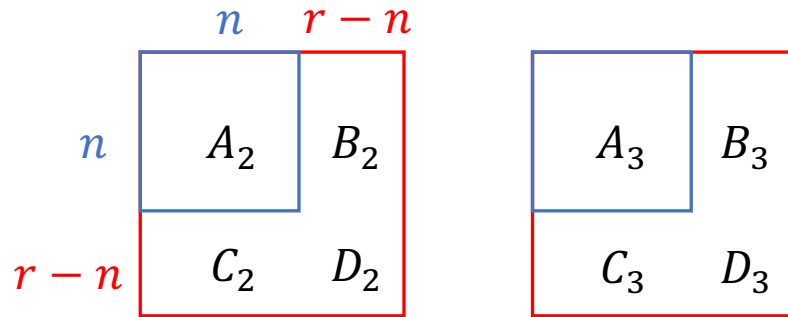
Commuting extension

Then:

$$r \geq n + \frac{1}{2} \text{rank}(A_2 A_3 - A_3 A_2)$$

Proof of the Koiran interpretation

Given: Commuting extension

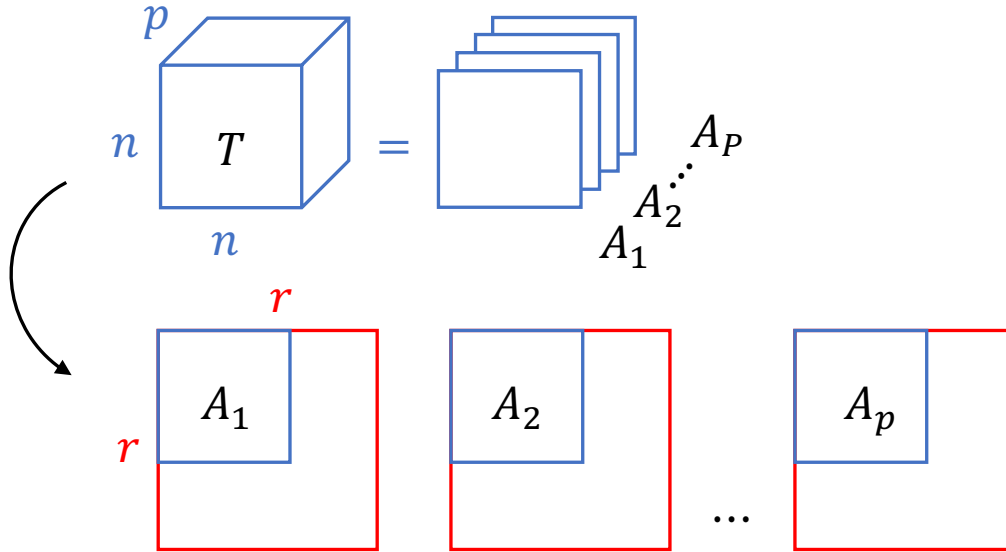


1. **Commutativity:** $A_2A_3 - A_3A_2 = B_3C_2 - B_2C_3$
2. $2(r - n) \geq \text{rank}(B_3C_2 - B_2C_3)$
3. $r \geq n + \frac{1}{2}\text{rank}(A_2A_3 - A_3A_2)$ ■

Matrix extensions characterize tensor rank

Theorem 3. (Koiran)

Given:



Commuting and diagonalizable extension

Then:

$$\text{rank}(T) + n \geq r_{\min} \geq \text{rank}(T)$$

Some words on commuting and diagonalizable matrices

B_1, \dots, B_p $n \times n$ matrices

- **Simultaneously** diagonalizable \Rightarrow Diagonalizable

$$B_k = V^{-1} D_k V \qquad B_k = V_k^{-1} D_k V_k$$

- Simultaneously diagonalizable \Rightarrow Commuting

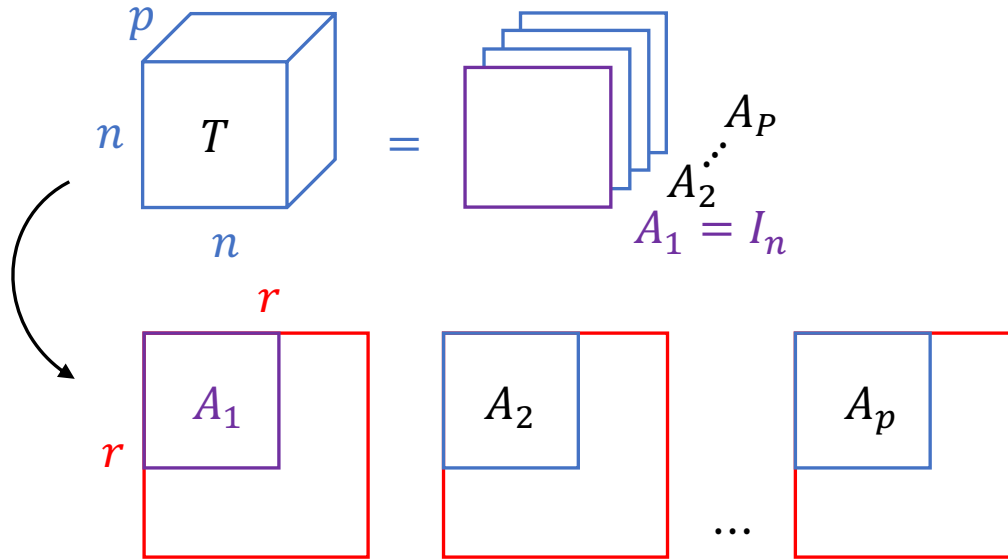
$$\begin{aligned} B_k B_l &= V^{-1} D_k V V^{-1} D_l V \\ &= V^{-1} D_k D_l V \\ &= V^{-1} D_l D_k V \end{aligned}$$

- Simultaneously diagonalizable \Leftrightarrow Commuting and diagonalizable

Simpler version of the theorem

Theorem 3'. (Koiran)

Given:

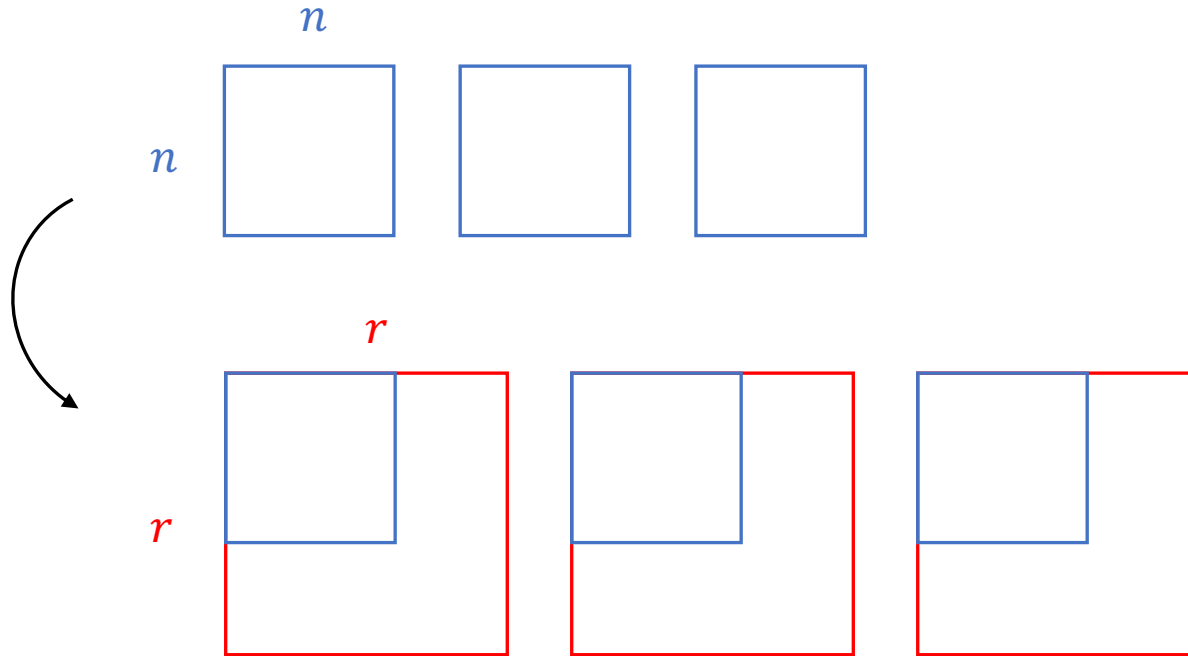


Commuting and diagonalizable extension

Then:

$$r_{\min} = \text{rank}(T)$$

Problem 1: Extend matrices until they commute



Commuting extension

(or: Commuting and diagonalizable extension)

Example: W-tensor

$$\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \quad \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}$$

$$\begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \quad \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{pmatrix}$$

Example: matrix multiplication tensor

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix} \begin{pmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

$$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 \\ 2 & -2 & 0 & 0 & 2 & 0 & 2 \\ -2 & 1 & 0 & 0 & -2 & -1 & -2 \\ -1 & 1 & 0 & 0 & -1 & 0 & -1 \end{pmatrix} \begin{pmatrix} 0 & 0 & 1 & 0 & 0 & 1 & -1 \\ 0 & 0 & 0 & 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & -1 & -1 & -1 & 0 \\ 0 & 0 & -1 & 1 & 1 & 0 & 1 \end{pmatrix} \begin{pmatrix} 0 & 0 & 0 & 0 & -\frac{1}{2} & 0 & -1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & \frac{1}{2} & 0 & 0 \\ -1 & 0 & 0 & -1 & -\frac{3}{2} & -1 & -1 \\ 1 & 0 & 0 & 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 & 2 & 1 & 2 \end{pmatrix}$$

Some questions

How can we lower bound commutative extensions?

commutative and diagonalizable extensions?