

Lecture 3. Cap set problem and slice rank

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5. Asymptotic subrank
6. Laser method

1. Cap set problem

- $x, y, z \in (\mathbb{Z}/3\mathbb{Z})^n$ form a **line** if $(x, y, z) = (u, u+v, u+2v)$ for some $u, v \in (\mathbb{Z}/3\mathbb{Z})^n$, $v \neq 0$.

$$\begin{array}{c} \bullet u+2v \\ \bullet u+v \\ \bullet u \end{array}$$

- How large can a line-free subset $A \subseteq (\mathbb{Z}/3\mathbb{Z})^n$ be?

$$\bullet 2.2202^{n-o(n)} \leq |A| \leq 2.755^{n+o(n)}$$

[Edel, Tyrell,
DeepMind]



construction



tensor obstruction
via subrank

[Ellenberg-Gryswald]

Lines in $\mathbb{Z}/3\mathbb{Z}$ and $(\mathbb{Z}/3\mathbb{Z})^n$

$$(x, y, z) = (u, u+v, u+2v), v \neq 0 \iff \{x, y, z\} = \{1, 2, 3\}$$

$$I, J, K \in (\mathbb{Z}/3\mathbb{Z})^n \text{ are a line} \iff \forall i \in [n], I_i = J_i = K_i \text{ or} \quad (*)$$
$$\{I_i, J_i, K_i\} = \{1, 2, 3\}$$

Cap set tensor

$$C := e_1 \otimes e_2 \otimes e_3 + \text{permutations}$$

$$+ e_1 \otimes e_1 \otimes e_1 + e_2 \otimes e_2 \otimes e_2 + e_3 \otimes e_3 \otimes e_3 \in \mathbb{F}_3^3 \otimes \mathbb{F}_3^3 \otimes \mathbb{F}_3^3$$

$$C^{\otimes n} = \sum \left\{ e_I \otimes e_J \otimes e_K : I, J, K \in [3]^n, (*) \right\}$$

Lemma If $A \subseteq (\mathbb{Z}/3\mathbb{Z})^n$ is line-free, then $C^{\otimes n}|_{A \times A \times A}$ is diagonal.

So, $|A| \leq \mathcal{Q}(C^{\otimes n})$. How can we upper bound this?

Matching

$\Psi \subseteq [d] \times [d] \times [d]$ is called a **matching** if $\forall x \neq y \in \Psi, \forall i \in [3], x_i \neq y_i$.

Example $\Psi = \text{supp}(\langle d \rangle) = \{(i, i, i) : i \in [d]\}$ is a matching.

Induced matchings

Let $\Phi \subseteq [d] \times [d] \times [d]$. We call $\Psi \subseteq \Phi$ **induced** if $\exists A_1, A_2, A_3 \subseteq [d], \Psi = \Phi \cap (A_1 \times A_2 \times A_3)$.

Example $\{(0,0,0), (1,1,1)\} \subseteq \{(0,0,0), (1,1,1), (2,2,2), (1,2,3) + \text{perm}\}$ is induced

Lemma Let $\Psi \subseteq \Phi$ be an induced matching. Then $|\Psi| \leq Q(T)$ for any tensor T with support Φ .

Proof • $\Psi = \Phi \cap (A_1 \times A_2 \times A_3)$ is a matching

• $\Psi = \text{supp}(T|_{A_1 \times A_2 \times A_3})$ so $\langle |\Psi| \rangle \leq T$. $\square$

2. Slice rank, asymptotic analysis

$$C = 111 + 222 + 333 + 123 + \text{perm.} \in \mathbb{F}^3 \otimes \mathbb{F}^3 \otimes \mathbb{F}^3$$

How can we upper bound $Q(C^{\otimes n})$?

(exercise)

Example For $\mathbb{F} = \mathbb{C}$, $C \cong \langle 3 \rangle$. Then $Q(C^{\otimes n}) = 3^n$, bad bound!

We take $\mathbb{F} = \mathbb{F}_3$!

Lemma $C \cong 110 + 101 + 011 + 002 + 020 + 200$ (over $\mathbb{F}_3$)

Hint: $C : [3] \times [3] \times [3] \rightarrow \mathbb{F}_3 \cdot (x, y, z) \mapsto \delta_0(x+y+z)$

$$\delta_0(x+y+z) = 1 - (x+y+z)^2 = 1 - x^2 - y^2 - z^2 + xy + yz + xz$$

This will be the right tensor to analyse, but we need a new tool.

Recall: tensor rank one: $u \otimes v \otimes w$

Slice rank

- Slice rank one: $\sum_i u \otimes v_i \otimes w_i$, $\sum_i u_i \otimes v \otimes w_i$ or $\sum_i u_i \otimes v_i \otimes w$
- $SR(T) := \min \left\{ r : T = \sum_{j=1}^r T_j, T_j \text{ has slice rank one} \right\}$

Alt:

- $SR(T) = \min \{ a+b+c : T \cong S, S \in \mathbb{F}^a \otimes \mathbb{F}^{n_2} \otimes \mathbb{F}^{n_3} + \mathbb{F}^{n_1} \otimes \mathbb{F}^b \otimes \mathbb{F}^{n_3} + \mathbb{F}^{n_1} \otimes \mathbb{F}^{n_2} \otimes \mathbb{F}^c \}$
- $SR(T) = \min \{ R_{(1)}(S_1) + R_{(2)}(S_2) + R_{(3)}(S_3) : T = S_1 + S_2 + S_3 \}$

Lemma

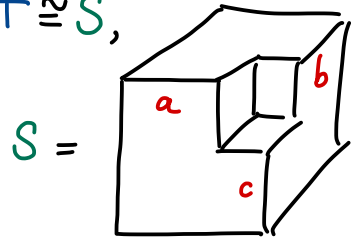
$$(1) \quad SR(\langle r \rangle) = r \quad [\text{Tao}]$$

$$(2) \quad S \leq T \Rightarrow SR(S) \leq SR(T).$$

Corollary $Q(T) \leq SR(T)$

Now we need to upper bound $SR(C^{\otimes n})$.

Visually: $T \cong S$,



Asymptotic analysis:

$$C \cong 110 + 101 + 011 + 002 + 020 + 200$$

$$C^{\otimes n} \cong \sum \left\{ IJK : \forall i \in [n], \overbrace{I_i + J_i + K_i} = 2 \right\} \quad \sum_i I_i \leq \frac{2}{3}n$$

$$= \sum \text{---} + \sum \text{---} + \sum \text{---}$$

$$I: \sum_i I_i \leq \frac{2}{3}n \quad J: \sum_i J_i \leq \frac{2}{3}n \quad K: \sum_i K_i \leq \frac{2}{3}n$$

$$\text{or } \sum_i J_i \leq \frac{2}{3}n$$

$$\text{or } \sum_i K_i \leq \frac{2}{3}n$$

$\underbrace{\hspace{10em}}$

$$R_{(1)} \leq \sum \left\{ \binom{n}{a,b,c} : a+b+c=n, b+2c \leq 2n/3 \right\} \leq O(2.755^n)$$

$$\text{So } SR(C^{\otimes n}) \leq O(2.755^n).$$

We will later see structured ways to determine $\lim_{n \rightarrow \infty} SR(T^{\otimes n})^{1/n}$.

3. Geometric rank, analytic rank, G-stable rank

We have $Q \leq SR \leq R(\ell)$ What other upper bounds are there on Q ?

Geometric rank $T: \mathbb{C}^{n_1} \times \mathbb{C}^{n_2} \rightarrow \mathbb{C}^{n_3}$ [Moshkowitz-Kopparty-Z]

$$GR(T) = \text{codim} \{ (x, y) : T(x, y) = 0 \}$$

Analytic rank $T: \mathbb{F}_p^{n_1} \times \mathbb{F}_p^{n_2} \rightarrow \mathbb{F}_p^{n_3}$ [Gowers-Wolf, Lovett]

$$AR(T) = -\log_p \frac{|\{ (x, y) : T(x, y) = 0 \}|}{p^{n_1 + n_2}} = -\log_p \mathbb{E}_{x, y, z} \omega_p^{T(x, y, z)}$$

G-stable rank [Derksen]

$$R^G(T) = \sup_{S \cong T} \min_i \frac{\|S\|_2^2}{\|S_{(i)}\|_\sigma^2} \leftarrow \begin{array}{l} \text{spectral norm} = \text{largest singular value} \\ \uparrow \\ \text{ith flattening matrix} \end{array}$$

Theorem SR, GR, AR, R^G are related up to a constant
(over appropriate fields). [Moshkovitz, Cohen, Zhu]

Still, they each have their own applications, e.g.

Theorem $\underline{Q}(\langle n, n, n \rangle) = \lceil \frac{3}{4} n^2 \rceil = GR(\langle n, n, n \rangle)$ ($< SR(\langle n, n, n \rangle) = n^2$)

Yet, they can be quite different from subrank:

Theorem For a random $n \times n \times n$ tensor T ,

- $\underline{Q}(T) \sim \sqrt{n}$

- $SR(T) = n$.

Theorem. $\underline{Q} \gtrsim \sim \sqrt{SR}$ [Qiyuan Chen - Ke Ye 2025]

4. Asymptotic subrank

For several applications (including cap sets and matrix mult.) the following parameter will be interesting.

Recall: $Q(T) := \max \{r : \langle r \rangle \leq T\}$

Asymptotic subrank $\tilde{Q}(T) := \lim_{n \rightarrow \infty} Q(T^{\otimes n})^{1/n} = \sup_n Q(T^{\otimes n})^{1/n}$.

Lemma $\tilde{Q}(T) \leq \liminf_{n \rightarrow \infty} SR(T^{\otimes n})^{1/n}$ (We don't know if the limit always exists...)

Theorem $\tilde{Q}(C) = 2.755\dots$

Thus analyzing slice rank of $C^{\otimes n}$ will not give any better bound on cap sets...

Where does the $\tilde{Q}$ lower bound come from?

5. Tight tensors

There is a general lower bound on $\underline{\alpha}$ for so-called tight tensors, and this lower bound is optimal (Thursday).

Tightness.

- $\Phi \subseteq [n] \times [n] \times [n]$ is **tight** if there are injective maps $f_i : [n] \rightarrow \mathbb{Z}$ such that $\forall u \in \Phi, \sum_i f_i(u_i) = 0$.
- $T \in \mathbb{F}^{n \times n \times n}$ is **tight** if $\exists S \cong T$, $\text{supp}(S)$ is tight.

Example These are tight:

- Example These are tight:
- Any matrix Pf: Any matrix is isomorphic to $\sum_{i=1}^r e_i \otimes e_i$. $f_1(i) = i$
 - $\langle r \rangle$ Pf: $\langle r \rangle = \sum_{i=1}^r e_i \otimes e_i \otimes e_i$. $f_2(i) = -i$
 - $\langle a, b, c \rangle$ (exercise) $f_1(i) = i, f_2(i) = i, f_3(i) = -2i$

Theorem Let $T \in \mathbb{F}^{n \times n \times n}$ have tight support $\Phi \subseteq [n] \times [n] \times [n]$.

$$\text{Then } \log_2 \mathcal{Q}(T) \geq \max_{P \in \mathcal{P}(\Phi)} \min_{i \in [3]} H(P_i)$$

Where $\mathcal{P}(\Phi)$ are the prob. distr. on Φ , P_i the i th marginal, and H the Shannon entropy.

Given Φ this is usually easy to compute.

- Example. • $W = 001 + 010 + 100$ is tight. $\mathcal{Q}(W) \stackrel{(=)}{\geq} 2^{h(1/3)} \stackrel{(=)}{\approx} 1.88$
- $C \cong 011 + 101 + 110 + 002 + 020 + 200$ is tight. $\mathcal{Q}(C) \geq 2.755\dots$
 - $\langle a, b, c \rangle$ is tight. $\mathcal{Q}(\langle a, b, c \rangle) \geq \min\{ab, bc, ac\}$.

In fact, for tight Φ , above can be obtained by induced matchings in $\Phi^{\otimes n}$.

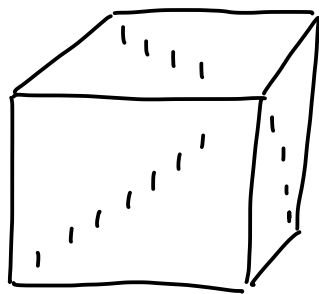
Proof relies on arithmetic progression-free sets. Also: the theorem is optimal!

6. Laser method

CW tensors $CW_q = \sum_{i=1}^q oii + ioi + iio \in \mathbb{F}^{q+1} \otimes \mathbb{F}^{q+1} \otimes \mathbb{F}^{q+1}$.

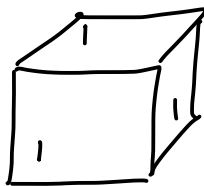
Lemma $\underline{R}(CW_q) \leq q+2$ Cheap! How to get matrix mult out of this?

CW_q is not tight (for $q \geq 2$), but it is tight if we squint... and consider the block tensor, grouping $\{0\}, \{1, \dots, q\}$ in each coordinate:



CW_q

blocked



W

$$W^{\otimes n} \geq \langle 1.88^{n-o(n)} \rangle = r(n)$$

$$CW_q^{\otimes n} \geq \bigoplus_{i=1}^{r(n)} \langle a_i, b_i, c_i \rangle$$

Asymptotic sum inequality: $\omega \leq 2.41$
 $\dots \omega \leq 2.37$