

Lecture 9 Convexity of asymptotic spectra

1. Convexity questions
2. Finite-dimensional picture
3. Examples and applications
4. Log-convexity for matrix multiplication
5. Log-star-convexity for rectangular matrix multiplication

1. Convexity questions

Asymptotic spectrum $X = \{\phi \in \text{Hom}(\text{tensors}, \mathbb{R}_{\geq 0}) : \leq - \text{monotone}\}.$

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- Support functionals too in $X(\text{oblique})$.

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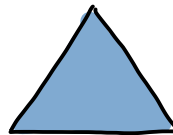
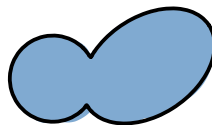
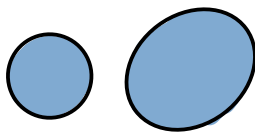
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- Quantum functionals interpolate between the flattening ranks
- Support functionals too in $X(\text{oblique})$.

We don't know if we can do this between arbitrary points in X .



Open problem Is X : disconnected connected stronger structure ?

2. Finite-dimensional picture

$$X \subseteq \mathbb{R}_{\geq 0}^{\text{tensors}}$$

infinite-dimensional

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fix tensor T of interest

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$\rightsquigarrow$

$$\bullet X(T) := \{ \phi(T) : \phi \in X \} \subseteq \mathbb{R}_{\geq 0}$$

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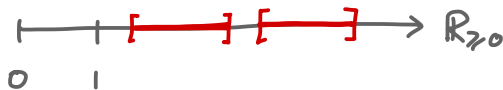
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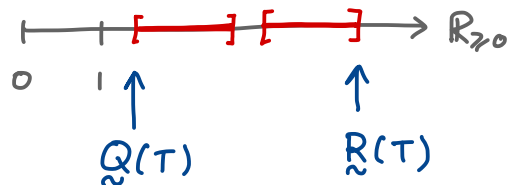
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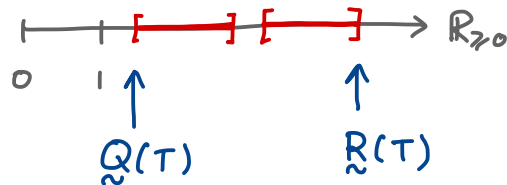
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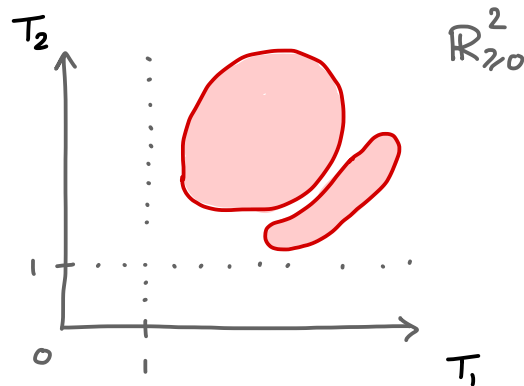
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$X(T_1, T_2)$



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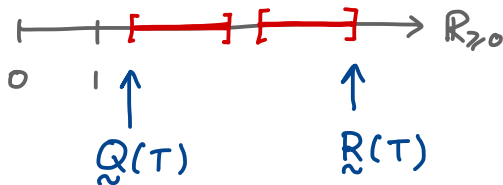
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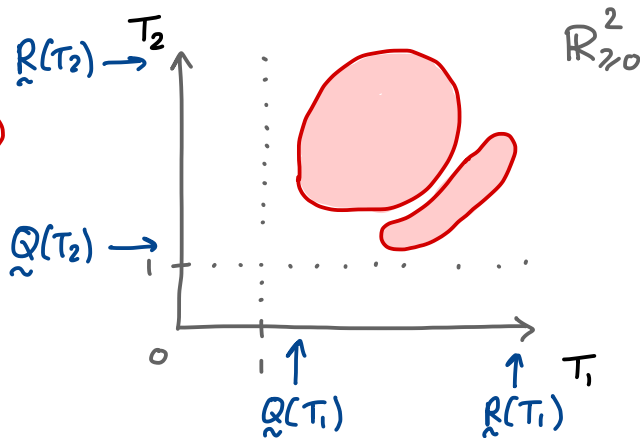
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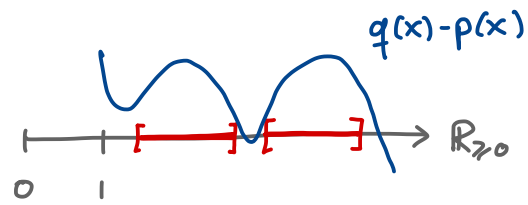
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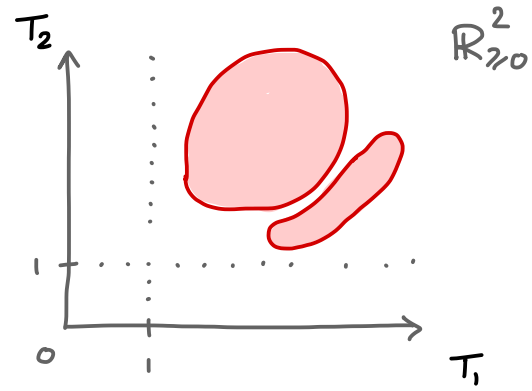
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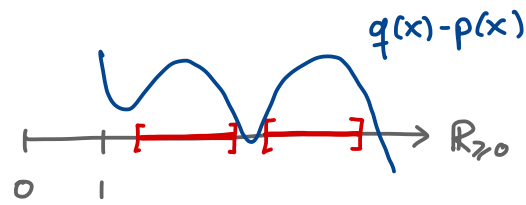


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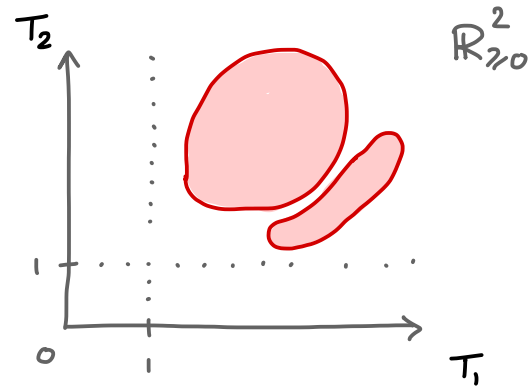


What other information do we get?

$X(\tau)$



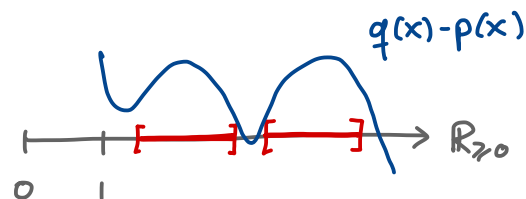
$X(\tau_1, \tau_2)$



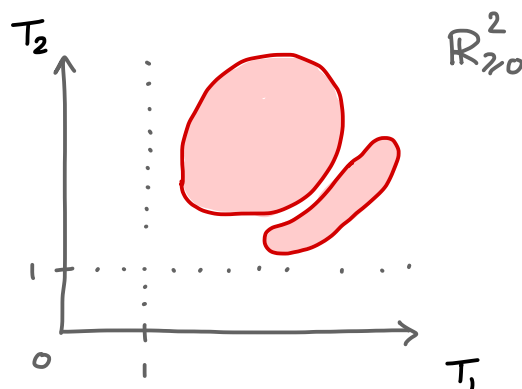
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Theorem. $\forall p, q \in \mathcal{IN}[x], \quad p(\tau) \leq q(\tau) \Leftrightarrow p(x) \leq q(x) \text{ on } X(\tau)$
 $\Leftrightarrow 0 \leq q(x) - p(x) \text{ on } X(\tau).$

$X(\tau)$



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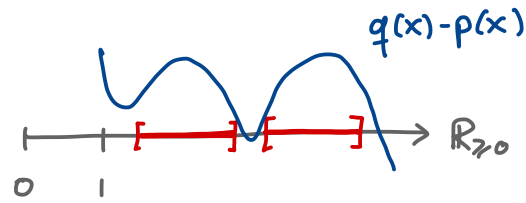


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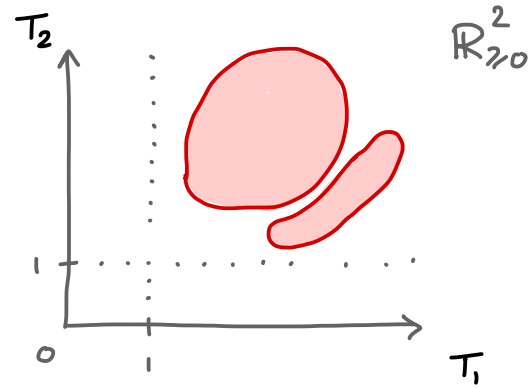
Theorem. $\forall p, q \in \mathbb{N}[x], \quad p(\tau) \preceq q(\tau) \Leftrightarrow p(x) \leq q(x) \text{ on } X(\tau)$
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Proof Duality: $p(\tau) \preceq q(\tau) \Leftrightarrow \forall \phi \in X, \quad \underbrace{\phi(p(\tau))}_{p(\phi(\tau))} \leq \underbrace{\phi(q(\tau))}_{q(\phi(\tau))}$ $\square$

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Theorem $\forall p, q \in \mathbb{N}[x_1, x_2], \quad p(\tau_1, \tau_2) \leq q(\tau_1, \tau_2) \Leftrightarrow 0 \leq q(x_1, x_2) - p(x_1, x_2)$
 etc. on $X(\tau_1, \tau_2).$

3. Examples and applications

Theorem $X(W) = [2^{h(1/3)} \approx 1.88, 2]$ proven using support functionals

$$p(w) \leq q(w) \iff 0 \leq q-p \text{ on } \text{---} \underset{1.88.}{\text{[}} \text{---} \underset{2}{\text{]}} \text{---}$$

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Theorem $\chi(W) = [2^{h(1/3)} \approx 1.88, 2]$ proven using support functionals

$p(w) \lesssim q(w) \iff 0 \leq q-p$ on $\text{---} \underset{1.80.}{\boxed{\text{---}}} \underset{2}{\boxed{\text{---}}}$

Theorem $X(\langle 2, 2, 2 \rangle)$ is connected, so equals $[4, 2^\omega]$

$$p(\langle 2,2,2 \rangle) \stackrel{\approx}{\sim} q(\langle 2,2,2 \rangle) \iff 0 \leq q-p \text{ on } \begin{array}{c} \text{---} [\text{---}] \text{---} \\ 4 \qquad 2^\omega \end{array}$$

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Actually we can do more

Theorem $X(\langle 2, 2, 2 \rangle)$ is connected, so equals $[4, 2^\omega]$

$$\rho(\langle 2,2,2 \rangle) \stackrel{\approx}{\sim} q(\langle 2,2,2 \rangle) \iff 0 \leq q - p \text{ on } [4, 2^w]$$

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$$\Leftrightarrow \forall x \in X(\langle 2, 2, 2 \rangle), \sum_i x^{\log n_i} \leq \sum_j x^{\log m_j} \quad \square.$$

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How is this useful? Upper bounds on w .



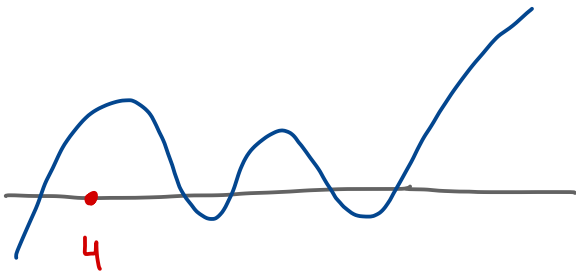
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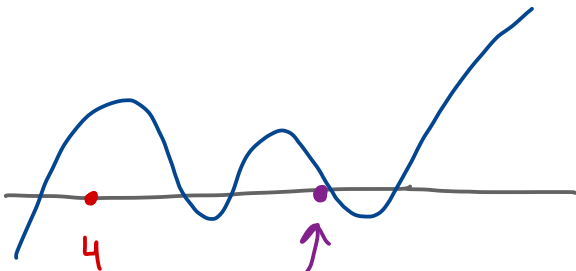
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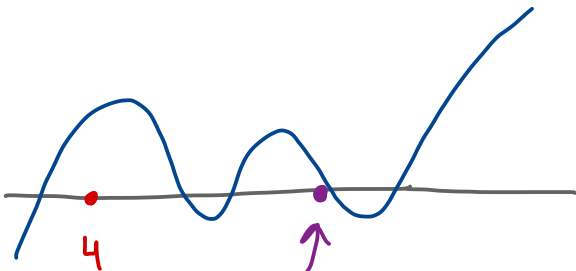
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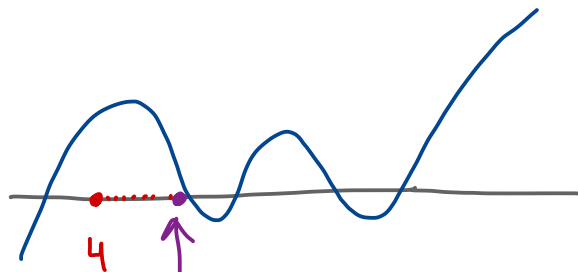
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So to prove log-convexity of say $X = X(\langle 2, 2, 2 \rangle)$ it suffices to prove $X = \tilde{X}$, so it suffices to show X is characterized by monom. ineq.

Monomials

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$$f \leq_M g \text{ on } X : \Leftrightarrow \forall \varepsilon > 0 \quad \forall N \text{ large enough}$$

$$\forall u \in \text{Mon}(f^N), \exists v \in \text{Mon}(2^{\varepsilon N} g^N), \quad u \leq v \text{ on } X$$

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Then $u \leq 2^{\varepsilon n} g^n$ on $\tilde{X}$.

Then $f^n \leq 2^{2\varepsilon n} g^n$ on $\tilde{X}$ (because f^n has $\text{poly}(n)$ many monomials)

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Then $f^n \leq 2^{2\varepsilon n} g^n$ on $\tilde{X}$ (because f^n has $\text{poly}(n)$ many monomials)

Let $\varepsilon \rightarrow 0$. $\square$

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We will prove:

Condition A

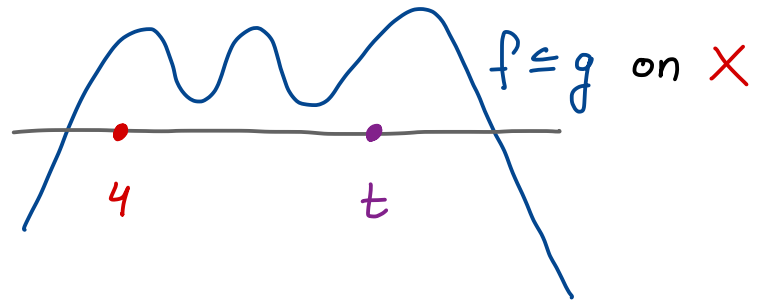
$\forall t \in X, f \leq g \text{ on } X \Rightarrow f \leq_M g \text{ on } \underbrace{\{u, t\}}_{R_1(\langle 2, 2, 2 \rangle) \in X(\langle 2, 2, 2 \rangle)}$

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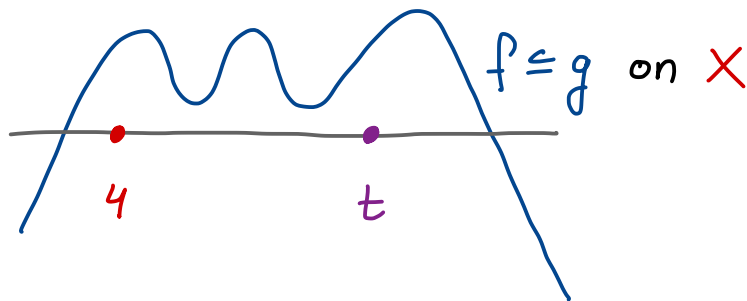
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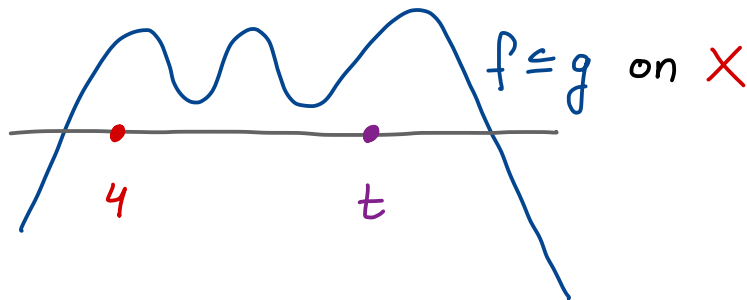
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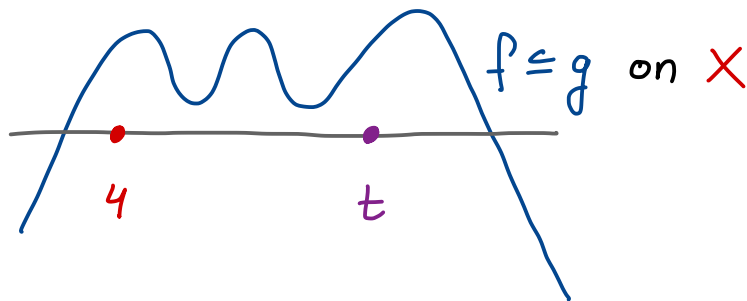
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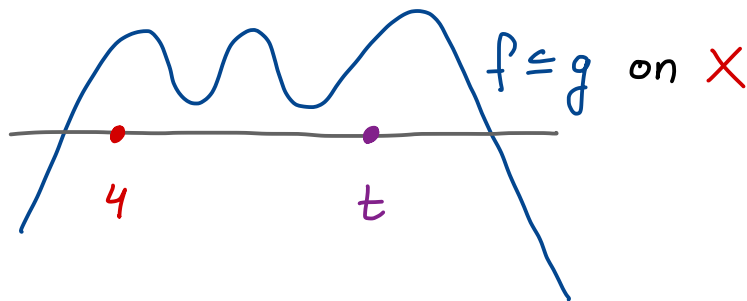
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4.2 Compression

It turns out we need to do this.

Theorem. If $\ell \cdot \langle 2, 2, 2 \rangle^{\otimes k} \subseteq S \oplus \mathcal{U}$ and $R_1(S) \subseteq \ell \cdot (2^n)^{k-2}$,
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Three ideas:

- stabilizer of $\langle a, b, c \rangle$
- local projection of $\langle a, b, c \rangle$ to $\langle a', b', c' \rangle$
- Matrix subspace lemma

Stabilizer

$$\langle n, n, n \rangle = (A, B, C) \mapsto \text{Tr}(ABC)$$

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Lemma

$\langle n, n, n \rangle \in \mathbb{F}^{n^2} \otimes \mathbb{F}^{n^2} \otimes \mathbb{F}^{n^2}$ is invariant under

$$\mathbb{F}^{n^2} \rightarrow \mathbb{F}^{n^2}: A \mapsto XAY^{-1}$$

$$\mathbb{F}^{n^2} \rightarrow \mathbb{F}^{n^2}: B \mapsto YBZ^{-1}$$

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Local projection

$$\langle a, b, c \rangle = \sum_{i \in [a]} \sum_{j \in [b]} \sum_{k \in [c]} e_{ij} \otimes e_{jk} \otimes e_{ki} \in \mathbb{F}^{ab} \otimes \mathbb{F}^{bc} \otimes \mathbb{F}^{ca}$$

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Lemma

$$(\text{Id} \otimes P_{b', c'} \otimes P_{c', a'}) \langle a, b, c \rangle = \langle a', b', c' \rangle$$

Matrix subspace lemma

$\mathbb{F}$ large enough

lemma let $V \subseteq \mathbb{F}^{n_1} \otimes \mathbb{F}^{n_2}$, $d = \dim V$.

Suppose $d \leq (b_1+1)(b_2+1)-1$ for some $0 \leq b_i \leq n_i$

Then there exist subspaces $W_i \subseteq \mathbb{F}^{n_i}$, $\dim(W_i) = n_i - b_i$

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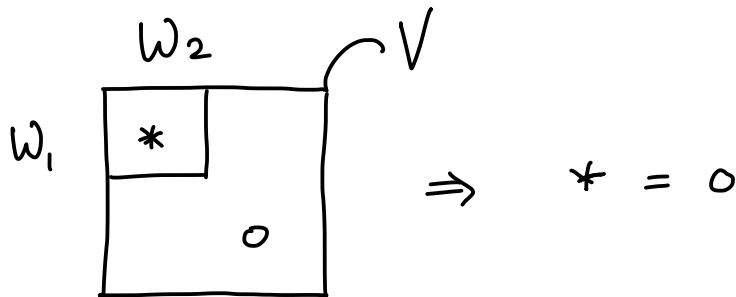
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Picture



Compression

Lemma. If $\langle n_1, n_2, n_3 \rangle \in S \oplus \mathcal{U}$ and $R_1(S) < \prod_i (n_i - m_i + 1)$
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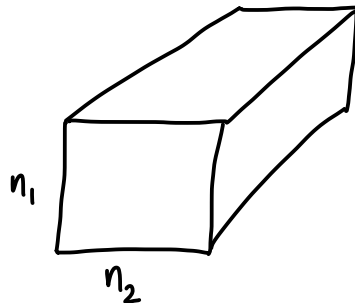
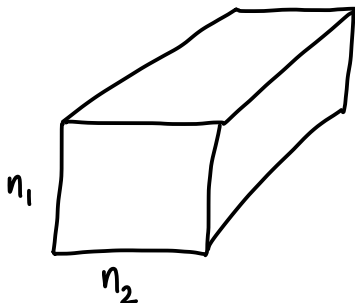
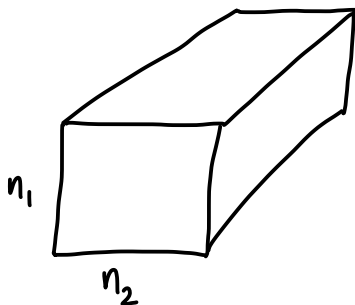
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Proof wlog $\langle n_1, n_2, n_3 \rangle = S + U \in \mathbb{F}^{n_1 n_2 \times n_2 n_3 \times n_3 n_1}$

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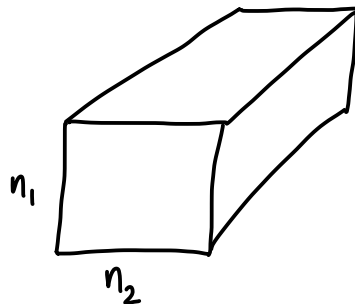
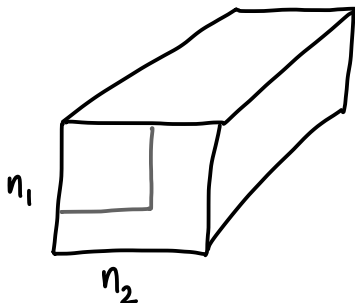
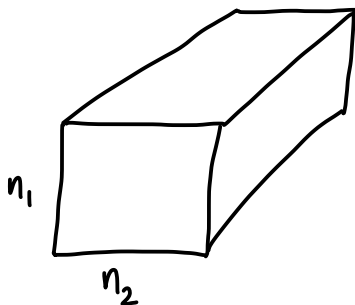
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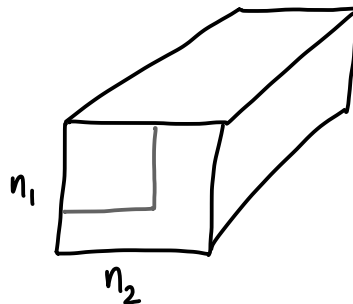
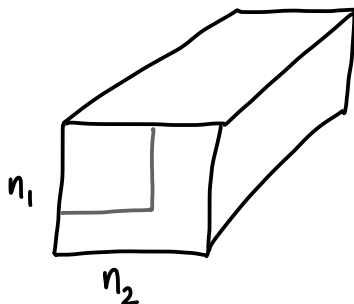
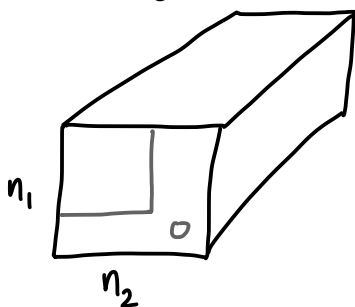
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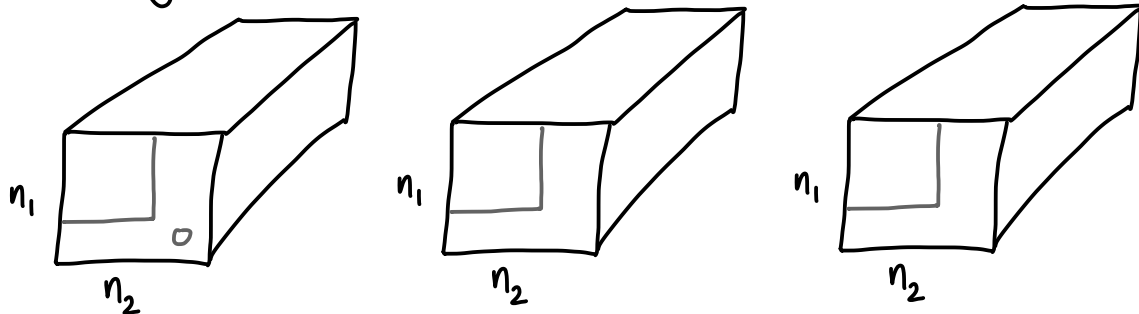
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Theorem $X(\langle 2, 2, 2 \rangle)$ is connected

Strassen proves something stronger

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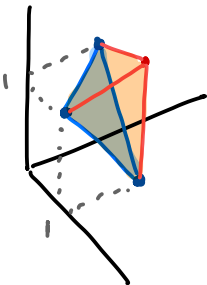
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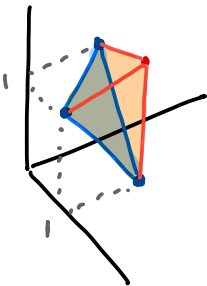
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Open problem Is it convex?

$$\bigoplus_i \langle a_i, b_i, c_i \rangle \approx \bigoplus_i \langle d_i, e_i, f_i \rangle \Rightarrow \sum_i (a_i b_i c_i)^{w/3} \leq \sum_i (d_i e_i f_i)^{w/3}$$