

Formalizing and Proving Theorems in Coq — Lecture 6

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Today's lecture

You will learn about. . .

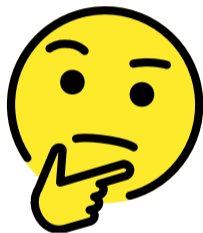
- ▶ the Curry-Howard isomorphism as an idea
- ▶ how this correspondence manifests in Coq

Making you wonder

Interesting coincidences:

- ▶ “*is of type*” and “*is a proof of*” are denoted by :
- ▶ *implication* and *functions* both use the symbol \rightarrow

This is completely intentional!



A logical pun

In Coq's underlying logic

- ▶ Proofs are expressions (programs)
- ▶ Propositions are types
- ▶ Simplification is execution



Demo

Exploiting the correspondence

We have just. . .

- ▶ constructed proofs via syntax.
- ▶ constructed definitions in proof mode.

Computational content

The Brouwer-Heyting-Kolmogorov interpretation of logic:

- ▶ A proof of $P \rightarrow Q$ *transforms* a proof of P into a proof of Q .
- ▶ A proof of $P \wedge Q$ is a proof of P *and* a proof of Q .
- ▶ A proof of $P \vee Q$ is a proof of P *or* a proof of Q .

Let's see how this works in Coq!

One step further

What about the first-order connective \forall ?

For instance, what is $\forall n : \text{nat}. P(n)$ as a type?

Alternatively, what is a value of type $\forall n : \text{nat}. P(n)$?

A function that, for every n , produces a value of type $P(n)$!

One step further

What about the first-order connective \exists ?

For instance, what is $\exists n : \text{nat}. P(n)$ as a type?

Alternatively, what is a value of type $\exists n : \text{nat}. P(n)$?

A pair of an n and a value of type $P(n)$!

It's up to you now

Start working on your projects!