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Generalised powerlocales via relation lifting

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This paper introduces an endofunctor V_T on the category of frames that is parametrised by an endofunctor T on the category Set that satisfies certain constraints. This generalises Johnstone's construction of the Vietoris powerlocale in the sense that his construction is obtained by taking for T the finite covariant power set functor. Our construction of the T-powerlocale $V_T \mathbb{L}$ out of a frame $\mathbb{L}$ is based on ideas from coalgebraic logic and makes explicit the connection between the Vietoris construction and Moss's coalgebraic cover modality.

We show how to extend certain natural transformations between set functors to natural transformations between T-powerlocale functors. Finally, we prove that the operation V_T preserves some properties of frames, such as regularity, zero-dimensionality and the combination of zero-dimensionality and compactness.

1. Introduction

The aim of this paper is to show how coalgebraic modal logic can be used to understand, study and generalise the point-free topological construction of taking Vietoris powerlocales.

1.1. *Hyperspaces and powerlocales*

The *Vietoris hyperspace construction* is a topological construction on compact Hausdorff spaces, which was introduced in Vietoris (1922) as a generalisation of the Hausdorff metric. Given a topological space X, one defines a new topology τ_X on K X, which is the set of compact subsets of X. This new topology τ_X has as its basis all sets of the form

$$\nabla\{U_1, \ldots, U_n\} := \{F \in \text{K } X \mid F \subseteq \bigcup_{i=1}^{n} U_i \text{ and } \forall i \leqslant n, F \between U_i\},$$

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where $U_1, \ldots, U_n \subseteq X$ is a finite collection of open sets and $F \between U$ is notation to indicate that $F \cap U \neq \varnothing$. Alternatively, one can use a subbasis to generate τ_X consisting of subbasic open sets of the shape

$$\Box U := \{F \in K\,X \mid F \subseteq U\}$$

and

$$\Diamond U := \{F \in K\,X \mid F \between U\}.$$

To generate the basic open sets $\nabla\{U_1, \ldots, U_n\}$ from $\Box U$ and $\Diamond U$, one can use the expression

$$\nabla\{U_1, \ldots, U_n\} = \Box\left(\bigcup_{i=1}^n U_i\right) \cap \bigcap_{i=1}^n \Diamond U_i.$$

In the field of *point-free topology*, a considerable amount of general topology has been recast in a way that makes it more compatible with constructive mathematics and topos theory (standard references are Johnstone (1982) and Vickers (1989)). The main idea is to study the lattices of open sets of topological spaces, rather than their associated sets of points. In other words, it is an approach to topology through algebra, where rather than categories of topological spaces, one studies categories of *locales*, or their algebraic counterparts, *frames*. Frames are complete lattices in which finite meets distribute over arbitrary joins, and can be seen as the algebraic models of propositional geometric logic, a branch of logic where finite conjunctions are studied in combination with infinite disjunctions. Substantial parts of this paper arose out of the direct application of techniques from coalgebraic logic to frames/locales. This has led to two consequences. The first is that most results are stated in terms of frames rather than locales, since frames are closer to the Boolean algebras predominantly used in coalgebraic logic. The second consequence is that we have given little consideration to issues of constructivity, in order to be able to apply coalgebraic logic techniques directly. We will briefly revisit these matters in Section 5. However, despite our bias towards frames, we have favoured the name 'powerlocale' over 'powerframe'.

Johnstone (1982) defines a point-free, syntactic version of the Vietoris powerlocale using an extension of geometric logic with two unary operators $\Box$ and $\Diamond$. However, he quickly also introduces expressions of the shape

$$\Box\left(\bigvee A\right) \wedge \bigwedge_{b \in B} \Diamond b,$$

where A and B are finite sets, which is reminiscent of the expression for $\nabla\{U_1, \ldots, U_n\}$ above. Nevertheless, the description of the Vietoris powerlocale is usually given with $\Box$ and $\Diamond$ as primitive, and not without good reason: one may obtain the Vietoris powerlocale by first constructing one-sided locales corresponding to the $\Box$-generators on the one hand and the $\Diamond$-generators on the other, and then joining these two one-sided powerlocales to obtain the Vietoris powerlocale (Vickers and Townsend 2004). However, the question remains as to whether one can describe the Vietoris powerlocale directly in terms of its basic opens, corresponding to $\nabla\{U_1, \ldots, U_n\}$, rather than the subbasic opens expressed in terms of $\Box$ and $\Diamond$. One of the main contributions of this paper is to show that this is indeed possible.

1.2. *The cover modality and coalgebraic modal logic*

The notation using $\Box$ and $\Diamond$ above is highly suggestive of modal logic. This is no coincidence: Johnstone's presentation of the Vietoris powerlocale in terms of generators and relations extends the axioms of positive (that is, negation-free) modal logic to the geometric setting.

In Boolean-based modal logic, one can define a ∇-modality which is applied to finite *sets* of formulas. This ∇-modality then has the following semantics. If $\mathfrak{M} = \langle W, R, V \rangle$ is a Kripke model and α is a finite set of formulas, then for any state $w \in W$,

$$\mathfrak{M}, w \Vdash \nabla\alpha \text{ iff } \forall a \in \alpha, \exists v \in R[w], \mathfrak{M}, v \Vdash a \text{ and}$$

$$\forall v \in R[w], \exists a \in \alpha, \mathfrak{M}, v \Vdash a.$$

In classical modal logic, the ∇-modality is equi-expressive with the $\Box$- and $\Diamond$-modalities using the following translations:

$$\nabla\alpha \equiv \Box \left(\bigvee\alpha \right) \wedge \bigwedge_{a \in \alpha} \Diamond a,$$

and, in the other direction,

$$\Box a \equiv \nabla\{a\} \vee \nabla\varnothing, \text{ and } \Diamond a \equiv \nabla\{a, \top\}.$$

As a primitive modality, ∇ was first introduced in Barwise and Moss (1996) in the study of circularity and in Janin and Walukiewicz (1995) in the study of the modal μ-calculus. It was in Moss's work (Moss 1999), however, that the ∇-modality stepped into the spotlight as a modality suitable for generalisation to the abstraction level of *coalgebras*.

The theory of Coalgebra aims to provide a general mathematical framework for the study of state-based evolving systems. Given an endofunctor T on the category Set of sets with functions, we have a coalgebra of type T, or briefly: a T-coalgebra is simply a function $\sigma: X \to TX$, where X is the underlying set of states of the coalgebra, and a T-coalgebra morphism between coalgebras $\sigma: X \to TX$ and $\sigma: X' \to TX'$ is simply a function $f: X \to X'$ such that $Tf \circ \sigma = \sigma' \circ f$. Aczel (1988) introduced T-coalgebras as a means to study transition systems. A natural example of such transition systems is provided by the *Kripke frames* and *Kripke models* used in the model theory of propositional modal logic: the category of Kripke frames and bounded morphisms is isomorphic to the category of P-coalgebras, where $P: \text{Set} \to \text{Set}$ is the covariant powerset functor. Universal coalgebra was later introduced in Rutten (2000) as a theoretical framework for modelling the behaviour of set-based transition systems that are parametric in their transition functor $T: \text{Set} \to \text{Set}$.

Coalgebraic logics are designed and studied in order to reason formally about coalgebras and their behaviour; one of the main applications of this approach is the design of specification and verification languages for coalgebras. The most influential approach to coalgebraic logic, known as *coalgebraic modal logic* (Cîrstea *et al.* 2009), is to try to generalise propositional modal logic from Kripke structures to the setting of arbitrary set-based coalgebras. Seminal for this approach was the observation by L. Moss in the paper mentioned earlier (Moss 1999) that the semantics of the cover modality ∇ can be described using the categorical technique of *relation lifting*. This observation paved

the way for generalisations to other functors that admit a reasonable notion of relation lifting: Moss introduced a modality ∇_T, which is parametric in the transition type functor T and can be interpreted in T-coalgebras through relation lifting.

While Moss's perspective was entirely semantic, his work naturally raised the question of whether good derivation systems could be developed for the coalgebraic cover modality ∇_T that is parametric in the coalgebra functor T. Building on earlier work by Bílková, Palmigiano and Venema (Palmigiano and Venema 2007; Bílková *et al.* 2008) for the power set case, Kupke *et al.* (2008, 2010) proved the soundness and completeness of such a derivation system $\mathbf{M}_T$. The latter paper also introduced an associated functor $\mathbb{M}_T$ on the category of Boolean algebras, which can be regarded as the algebraic correspondent of the topological Vietoris functor on the dual category of Stone spaces.

1.3. *Contribution*

In this paper we translate the coalgebraic modal derivation system $\mathbf{M}_T$ from its Boolean origins (Bílková *et al.* 2008; Kupke *et al.* 2008) to the setting of geometric logic. Basically, this means we take some initial steps towards developing a *geometric* coalgebraic modal logic, that is, a logic with finite conjunctions, infinite disjunctions and the coalgebraic cover modality ∇_T.

The main conceptual contribution of this paper is the introduction of a *generalised powerlocale construction* V_T that is parametric in a functor $T : \mathsf{Set} \to \mathsf{Set}$ satisfying some categorical conditions. Given a frame $\mathbb{L}$, we define its T-powerlocale $V_T\mathbb{L}$ using a presentation that takes the set $\{\nabla_T \alpha \mid \alpha \in TL\}$ as generators and the geometric version of the ∇-axioms as relations.

As we will see, the classical Vietoris powerlocale construction is an instantiation of the T-powerlocale, where we take $T = P_\omega$, the covariant finite power set functor. This reveals that the connection between the Vietoris construction and the cover modality, which was already implicit in *semantic* form in Vietoris (1922), can also be made explicit *syntactically* using coalgebraic modal logic. Our approach shows how to describe the Vietoris constructions syntactically using the ∇-expressions as primitives, rather than as expressions derived from $\square$- and $\diamond$-primitives, as introduced in Johnstone (1982).

In addition, we prove some technical results concerning the T-powerlocale construction. To start with, we discuss some *functorial properties*; in particular, we show that we are in fact dealing with a *functor*

$$V_T : \mathsf{Fr} \to \mathsf{Fr}$$

on the category of frames with algebraic frame homomorphisms. Furthermore, we show how to extend certain natural transformations between transition functors to natural transformations between T-powerlocale functors; this generalises, for instance, the frame homomorphism from the Vietoris locale onto the original frame. We also give an alternative *flat-site presentation* of the T-powerlocale $V_T\mathbb{L}$, showing that each element of a T-powerlocale has a disjunctive normal form. Finally, we prove some first *preservation results*; in particular, we show that the operation V_T preserves some important

properties of frames, such as regularity, zero-dimensionality and the combination of zero-dimensionality and compactness.

1.4. *Overview*

In Section 2 we introduce preliminaries on category theory, relation lifting, frame presentations and the classical point-free presentation of the powerlocale. In Section 3 we introduce the T-powerlocale construction V_T. We then show that the P_ω-powerlocale is isomorphic to the classical Vietoris powerlocale and discuss some functorial properties of the construction. We conclude this section by providing the above-mentioned flat-site presentation of T-powerlocales. In Section 4 we prove our preservation results and provide a new and constructively valid proof of the preservation of compactness for the 'classical' Vietoris construction. We conclude in Section 5 with some possibilities for future work.

2. Preliminaries

2.1. *Basic mathematics*

We begin by fixing some mathematical notation and terminology. Let $f : X \to X'$ be a function. Then the *graph* of f is the relation

$$Gr\, f ::= \{(x, f(x)) \in X \times X' \mid s \in X\}.$$

Given a relation $R \subseteq X \times X'$, we denote the *domain* and *range* of R by $\mathsf{dom}(R)$ and $\mathsf{rng}(R)$, respectively. Given subsets $Y \subseteq X$, $Y \subseteq X'$, the *restriction* of R to Y and Y' is given by

$$R{\restriction}_{Y \times Y'} ::= R \cap (Y \times Y').$$

The composition of two relations $R \subseteq X \times X'$ and $R' \subseteq X' \times X''$ is denoted by $R \mathbin{;} R'$, whereas the composition of two functions $f : X \to X'$ and $f' : X' \to X''$ is denoted by $f' \circ f$. Thus, we have $Gr\,(f' \circ f) = Gr\, f \mathbin{;} Gr\, f'$.

We will use $P(X)$ and $P_\omega(X)$ to denote the *power set* and *finite power set* of a given set X. The *diagonal on* X is the relation $\Delta_X = \{(x, x) \mid x \in X\}$. Given two sets X, Y, we say that X *meets* Y (notation, $X \between Y$) if $X \cap Y$ is inhabited (that is, non-empty).

A *pre-order* is a pair (X, R) where R is a reflexive and transitive relation on X. Given such a pre-order, we define the operations $\downarrow_{(X,R)}, \uparrow_{(X,R)} : PX \to PX$ by

$$\downarrow_{(X,R)}(Y) := \{x \in X \mid x\, R\, y \text{ for some } y \in Y\}$$
$$\uparrow_{(X,R)}(Y) := \{x \in X \mid y\, R\, x \text{ for some } y \in Y\}.$$

If no confusion is likely, we will write $\downarrow_X$ or $\downarrow$ rather than $\downarrow_{(X,R)}$.

2.2. *Category theory*

We will assume familiarity with the basic notions from category theory, including those of categories, functors, natural transformations and (co-)monads – Mac Lane (1998), for instance, may be consulted as a reference text.

We let Set denote the category with sets as objects and functions as morphism; endofunctors on this category will simply be called *set functors*. The most important set functor we shall use is the covariant power set functor P, which is in fact (part) of a monad (P, μ, η), with $\eta_X : X \to P(X)$ denoting the singleton map $\eta_X : x \mapsto \{x\}$ and $\mu_X : PPX \to PX$ denoting union, $\mu_X(\mathscr{A}) := \bigcup \mathscr{A}$. The contravariant power set functor will be denoted by $\check{P}$.

We will restrict our attention to set functors satisfying certain properties, of which the first is crucial. In order to define it, we need to recall the notion of a (weak) pullback. Given two functions $f_0 : X_0 \to X$ and $f_1 : X_1 \to X$, a *weak pullback* is a set P, together with two functions $p_i : P \to X_i$ such that $f_0 \circ p_0 = f_1 \circ p_1$, and, in addition, for every triple (Q, q_0, q_1) also satisfying $f_0 \circ q_0 = f_1 \circ q_1$, there is an arrow $h : Q \to P$ such that $q_0 = h \circ p_0$ and $q_1 = h \circ p_0$: diagrammatically:
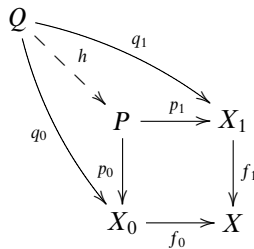
For (P, p_0, p_1) to be a *pullback*, we also require that the arrow h is unique.

A functor T *preserves weak pullbacks* if it transforms every weak pullback (P, p_0, p_1) for f_0 and f_1 into a weak pullback (TP, Tp_0, Tp_1) for Tf_0 and Tf_1. An equivalent characterisation is to require T to *weakly preserve pullbacks*, that is, to turn pullbacks into weak pullbacks. We will see yet another, and motivating, characterisation of this property in Section 2.3.

The second property we will impose on our set functors is that of standardness. Given two sets X and X' such that $X \subseteq X'$, let $\iota_{X,X'}$ denote the inclusion map from X into X'. A weak pullback-preserving set functor T is *standard* if it *preserves inclusions*, that is, $T\iota_{X,X'} = \iota_{TX,TX'}$ for every inclusion map $\iota_{X,X'}$.

Remark 2.1. Unfortunately, the definition of standardness is not uniform throughout the literature. Our definition of standardness is taken from Moss (1999), while, for instance, Adámek and Trnková (1990) has an additional condition involving so-called distinguished points. Fortunately, the two definitions are equivalent when the functor preserves weak pullbacks – see Kupke (2006, Lemma A.2.12).

The restriction to standard functors is not essential, since every set functor is 'almost standard' (Adámek and Trnková 1990, Theorem III.4.5): given an arbitrary set functor T, we may find a standard set functor T' such that the restriction of T and T' to all non-empty sets and non-empty functions are naturally isomorphic.

Finally, we shall require that our functors are determined by their behaviour on finite sets. We say a standard set functor T is *finitary* if $TX = \bigcup \{TX' \mid X' \subseteq_\omega X\}$. Our focus on finitary functors is not so much a restriction as a convenient way to express the fact that we are interested in the *finitary version* of an arbitrary set functor in the sense that

P_ω is the finitary version of P. Generally, we may define, for a standard functor T, the functor T_ω that on objects X is defined by $T_\omega X = \bigcup \{TX' \mid X' \subseteq X\}$, while on arrows f we simply put $T_\omega f := Tf$.

Since there are many set functors that are standard, finitary and weak pullback-preserving, the results in this paper have a wide scope.

Example 2.2. The identity functor Id, the finitary power set functor P_ω and, for each set Q, the constant functor C_Q (given by $C_Q X = Q$ and $C_Q f = id_Q$) are standard, finitary, and preserve weak pullbacks.

For a slightly more involved example, consider the finitary *multiset* functor M_ω. This functor takes a set X to the collection $M_\omega X$ of maps $\mu : X \to \mathbb{N}$ of finite support (that is, for which the set $Supp(\mu) := \{x \in X \mid \mu(x) > 0\}$ is finite), and its action on arrows is defined as follows. Given an arrow $f : X \to X'$ and a map $\mu \in M_\omega X$, we define $(M_\omega f)(\mu) : X' \to \mathbb{N}$ by putting

$$(M_\omega f)(\mu)(x') := \sum \{\mu(x) \mid f(x) = x'\}.$$

With this definition, the functor is not standard, but we may 'standardise' it by representing any map $\mu : X \to \mathbb{N}$ of finite support by its 'support graph' $\{(x, \mu x) \mid \mu x > 0\}$. As a variant of M_ω, consider the finitary probability functor D_ω, where

$$D_\omega X = \left\{ \delta : X \to [0,1] \mid Supp(\delta) \text{ is finite and } \sum_{x \in X} \delta(x) = 1 \right\},$$

while the action of D_ω on arrows is just like that of M_ω.

Perhaps more importantly, the class of finitary, standard functors that preserve weak pullbacks is closed under the following operations: composition $(\circ)$, product $(\times)$, co-product $(+)$ and exponentiation with respect to some set D $((\cdot)^D)$. As a corollary, we can inductively define the following class $EKPF_\omega$ of *extended finitary Kripke polynomial functors*:

$$T ::= Id \mid P_\omega \mid C_Q \mid M_\omega \mid D_\omega \mid T_0 \circ T_1 \mid T_0 + T_1 \mid T_0 \times T_1 \mid T^D.$$

Hence, all extended Kripke polynomial functors falls within the scope of the work in this paper.

As running examples in this paper, we will often take the *binary tree functor* $B = Id \times Id$ and the finitary power set functor P_ω.

An interesting property of standard functors is that they preserve finite intersections (Adámek and Trnková 1990, Theorem III.4.6), that is,

$$T(X \cap Y) = TX \cap TY.$$

As a consequence, if T is finitary, for any object $\xi \in TX$, we may define

$$Base_X^T(\xi) := \bigcap \{X' \in P_\omega(X) \mid \xi \in TX'\},$$

and show that $Base_X^T(\xi)$ is the *smallest* set X' such that $\xi \in TX'$ (Venema 2006). In fact, the base maps provide a natural transformation $Base^T : T \to P_\omega$; we will write this fact down explicitly in the next section so that we can refer to it.

To facilitate the reasoning in this paper, which will involve objects of various different types, we use the following variable naming convention.

Convention 2.3. Let X be a set and let $T : \mathsf{Set} \to \mathsf{Set}$ be a functor. We use the following naming convention:

| Set	Elements
X	$a, b, \ldots, x, y, \ldots$
TX	$\alpha, \beta, \ldots$
PX	$A, B, \ldots$
PTX	$\Gamma, \Delta, \ldots$
TPX	$\Phi, \Psi, \ldots$

2.3. *Relation lifting*

In Section 1, we mentioned that coalgebraic modal logic using the cover modality, as introduced by Moss, crucially uses relation lifting, both for its syntax and semantics. Relation lifting is a technique that allows us to extend a functor $T : \mathsf{Set} \to \mathsf{Set}$ defined on the category of sets to a functor $\overline{T} : \mathsf{Rel} \to \mathsf{Rel}$ on the category of sets and relations in a natural way. In this section, we will introduce some of the basic facts and definitions about relation lifting.

Let T be a set functor. Given two sets X and X', and a binary relation R between $X \times X'$, we define the *lifted relation* $\overline{T}(R) \subseteq TX \times TX'$ by

$$\overline{T}(R) := \{((T\pi)(\rho), (T\pi')(\rho)) \mid \rho \in TR\},$$

where $\pi : R \to X$ and $\pi' : R \to X'$ are the projection functions given by $\pi(x, x') = x$ and $\pi'(x, x') = x'$. Diagrammatically, we have

$$X \xleftarrow{\;\;\pi\;\;} R \xrightarrow{\;\;\pi'\;\;} X'$$
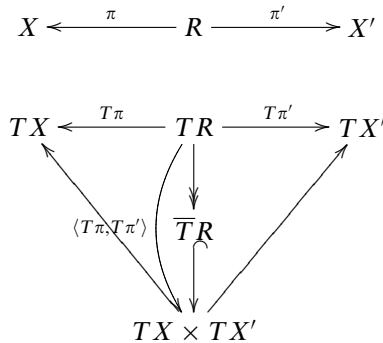
In other words, we apply the functor T to the relation R, seen as a *span*

$$X \xleftarrow{\;\pi\;} R \xrightarrow{\;\pi'\;} X' \,,$$

and define $\overline{T}R$ as the image of TR under the product map $\langle T\pi, T\pi' \rangle$ obtained from the lifted projection maps $T\pi$ and $T\pi'$.

We will now give some concrete examples.

Example 2.4. Fix a relation $R \subseteq X \times X'$. For the identity and constant functors, we find, respectively:

$$\overline{Id}\,R = R$$
$$\overline{C_Q}R = \Delta_Q.$$

The relation lifting associated with the power set functor P can be defined concretely as follows:

$$\overline{P}R = \{(A,A') \in PX \times PX' \mid \forall a \in A\, \exists a' \in A'.aRa' \text{ and } \forall a' \in A'\, \exists a \in A.aRa'\}.$$

This relation is known under many names, of which we will just mention that of the *Egli–Milner* lifting of R. For any standard, weak pullback preserving functor T, it can be shown (Kupke *et al.* 2010) that the lifting of T_ω agrees with that of T in the sense that

$$\overline{T_\omega}R = \overline{T}R \cap (T_\omega X \times T_\omega X').$$

From this it follows that

$$\text{for all } A \in T_\omega X, A' \in T_\omega X' : A\,\overline{P_\omega}R\,A' \text{ iff } A\,\overline{P}R\,A',$$

and for this reason, we shall write $\overline{P}R$ rather than $\overline{P_\omega}R$.

Relation lifting for the finitary multiset functor is slightly more involved: given two maps $\mu \in M_\omega X, \mu' \in M_\omega X'$, we put

$$\mu\,\overline{M_\omega}R\,\mu' \text{ iff there is some map } \rho : R \to \mathbb{N} \text{ such that}$$
$$\forall x \in X. \sum\{\rho(x,x') \mid x' \in X'\} = 1, \text{ and}$$
$$\forall x' \in X'. \sum\{\rho(x,x') \mid x \in X\} = 1.$$

The definition of $\overline{D_\omega}$ is similar.

Finally, relation lifting interacts well with various operations on functors (Hermida and Jacobs 1998). In particular, we have

$$\overline{T_0 \circ T_1}R = \overline{T_0}(\overline{T_1}R)$$
$$\overline{T_0 + T_1}R = \overline{T_0}R \cup \overline{T_1}R$$
$$\overline{T_0 \times T_1}R = \left\{ \big((\xi_0,\xi_1),(\xi'_0,\xi'_1)\big) \mid (\xi_i,\xi'_i) \in \overline{T_i}, \text{ for } i \in \{0,1\} \right\}$$
$$\overline{T^D}R = \{(\varphi,\varphi') \mid (\varphi(d),\varphi'(d)) \in \overline{T}R \text{ for all } d \in D\}.$$

Remark 2.5. Strictly speaking, the definition of the relation lifting of a given relation R depends on the type of the relation, that is, given sets X, X', Y, Y' such that $R \subseteq X \times X'$ and $R \subseteq Y \times Y'$, it matters whether we look at R as a relation from X to X' or as a relation from Y to Y'. We have avoided this potential source of ambiguity by requiring the functor T to be *standard* – see Fact 2.6(6).

Relation lifting has a number of properties that we will use throughout the paper. It can be shown that relation lifting interacts well with the operation of taking the graph of a function $f : X \to X'$, and with most operations on binary relations. Most of the properties below are easy to establish – see Kupke *et al.* (2010) for proofs.

Fact 2.6. Let T be a set functor. Then the relation lifting $\overline{T}$ satisfies the following properties for all functions $f : X \to X'$, all relations $R, S \subseteq X \times X'$ and all subsets $Y \subseteq X$, $Y' \subseteq X'$:

(1) $\overline{T}$ extends T: that is,
$\overline{T}(Gr\, f) = Gr\,(T f)$.
(2) $\overline{T}$ preserves the diagonal: that is,
$\overline{T}(\Delta_X) = \Delta_{TX}$.
(3) $\overline{T}$ commutes with relation converse: that is,
$\overline{T}(R\breve{\ }) = (\overline{T}R)\breve{\ }$.
(4) $\overline{T}$ is monotone: that is,
if $R \subseteq S$ then $\overline{T}(R) \subseteq \overline{T}(S)$.
(5) $\overline{T}$ distributes over composition: that is,
$\overline{T}(R \,;\, S) = \overline{T}(R) \,;\, \overline{T}(S)$, if T preserves weak pullbacks.
(6) $\overline{T}$ commutes with restriction: that is,
$\overline{T}(R{\restriction}_{Y \times Y'}) = \overline{T}R{\restriction}_{TY \times TY'}$ if T is standard and preserves weak pullbacks.

Fact 2.6(5) plays a key role in our work. In fact, distributivity of $\overline{T}$ over relation composition is *equivalent* to T preserving weak-pullbacks; the proof of this equivalence goes back to Trnková (1977).

Many proofs in this paper will be based on Fact 2.6, and we will not always provide all technical details. In the lemma below, we have isolated some facts that will be used a number of times, so that the proof may serve as an example of an argument using properties of relation lifting.

Lemma 2.7. Let $T : \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor. Let X, Y be sets, $f, g : X \to Y$ be two functions and $R \subseteq X \times X$ and $S \subseteq Y \times Y$ be relations. Then:

(1) If (X, R) is a pre-order, then so is $(TX, \overline{T}R)$.
(2) If $f(x)\, S\, g(x)$ for all $x \in X$, then $T f(\alpha)\, \overline{T}S\, Tg(\alpha)$ for all $\alpha \in TX$.
(3) If $x\, R\, y$ implies $f(x)\, S\, g(y)$ for all $x, y \in X$, then $\alpha\, \overline{T}R\, \beta$ implies $(Tf)\alpha\, \overline{T}S\, (Tg)\beta$ for all $\alpha, \beta \in TX$.

Proof.

(1) Observe that (X, R) is a pre-order if and only if $\Delta_X \subseteq R$ and $R \,;\, R \subseteq R$. Hence, if (X, R) is a pre-order, it follows from Fact 2.6(2 and 4) that $\Delta_{TX} = \overline{T}\Delta_X \subseteq \overline{T}R$, and from Fact 2.6(5 and 4) that $\overline{T}R \,;\, \overline{T}R = \overline{T}(R \,;\, R) \subseteq \overline{T}R$, implying that $(TX, \overline{T}R)$ is a pre-order as well.
(2) Observe that the antecedent can be succinctly expressed as

$$(Gr\, f)\breve{\ } \,;\, Gr\, g \subseteq S.$$

Then it follows by the properties of relation lifting that

$$(Gr\ Tf)\check{}\ ;\ Gr\ Tg = (\overline{T}(Gr\ f))\check{}\ ;\ \overline{T}(Gr\ g) \qquad \text{(Fact 2.6(1))}$$
$$= \overline{T}((Gr\ f)\check{})\ ;\ \overline{T}(Gr\ g) \qquad \text{(Fact 2.6(3))}$$
$$= \overline{T}((Gr\ f)\check{}\ ;\ Gr\ g) \qquad \text{(Fact 2.6(5))}$$
$$\subseteq \overline{T}S \qquad \text{(Fact 2.6(4))}.$$

But the inclusion $(Gr\ Tf)\check{}\ ;\ Gr\ Tg \subseteq \overline{T}S$ is just another way of stating the conclusion of part 2.

(3) We reformulate the statement of its antecedent as

$$(Gr\ f)\check{}\ ;\ R\ ;\ Gr\ g \subseteq S.$$

From this we may reason using a completely analogous argument to the one just given that

$$(Gr\ Tf)\check{}\ ;\ \overline{T}R\ ;\ Gr\ Tg \subseteq TS,$$

which is an equivalent way of phrasing the conclusion of part 3. $\qquad\square$

Relation lifting interacts with the map $Base^T$ as follows (see Kupke *et al.* (2010)).

Fact 2.8. Let T be a standard, finitary, weak pullback-preserving functor. Then:

(1) $Base^T$ is a natural transformation $Base^T : T \to P_\omega$. That is, given a map $f : X \to X'$, the following diagram commutes:

$$
\begin{array}{ccc}
TX & \xrightarrow{Base_X^T} & P_\omega X \\
\downarrow{\scriptstyle Tf} & & \downarrow{\scriptstyle Pf} \\
TX' & \xrightarrow{Base_{X'}^T} & P_\omega X'
\end{array}
$$

(2) Given a relation $R \subseteq X \times X'$ and elements $\alpha \in TX$, $\beta \in TY$, it follows from $\alpha\ \overline{T}R\ \beta$ that $Base^T(\alpha)\ \overline{P}R\ Base^T(\beta)$.

An interesting relation to which we shall apply relation lifting is the *membership* relation $\in$. If needed, we will denote the membership relation restricted to a given set X as the relation $\in_X \subseteq X \times PX$. Given a set X and $\Phi \in TPX$, we define

$$\lambda_X^T(\Phi) = \{\alpha \in TX \mid \alpha\ \overline{T}\in_X\ \Phi\}.$$

Elements of $\lambda^T(\Phi)$ will be called *lifted members of* Φ. The properties of λ^T are intimately related to those of $\overline{T}$, as in the following fact (Kupke *et al.* 2010).

Fact 2.9. Let $T : \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor. Then the collection of maps λ_X^T forms a distributive law with respect to both the covariant and the contravariant power set functors. That is, λ^T provides two natural transformations

$$\lambda^T : TP \to PT$$
$$\lambda^T : T\check{P} \to \check{P}T.$$

Remark 2.10. We can strengthen Fact 2.9 as follows. λ^T is actually a distributive law over the *monad* (P, μ, η) in the sense of also being compatible with the unit η and the multiplication μ of P, as given by the following diagrams:

$$
\begin{array}{ccc}
TX \xrightarrow{T\eta_X} TPX \\
\quad {}_{\eta_{TX}}\searrow \quad \downarrow {}^{\lambda^T_X} \\
PTX
\end{array}
\qquad\qquad
\begin{array}{ccccc}
TPPX \xrightarrow{\lambda^T_{PX}} PTPX \xrightarrow{P\lambda^T_X} PPTX \\
\downarrow {}^{T\mu_X} \qquad\qquad\qquad\qquad \downarrow {}^{\mu_X} \\
TPX \xrightarrow[\lambda^T_X]{} PTX
\end{array}
$$

In the terminology of Street (1972), (T, λ^T) is a *monad opfunctor* from the monad P to itself, and there is a one–one correspondence between the monad opfunctors and the functors T equipped with extensions to endofunctors on the *Kleisli category* $\mathsf{Kl}(P)$ associated with P. (The explicit results given in Street (1972), using the 2-functor $\mathbf{Alg}_\mathsf{C}$, are in terms of monad functors and extensions to the category of Eilenberg–Moore algebras. The results for monad opfunctors and the Kleisli category are dual.) Note that the Kleisli category of the power set monad is (isomorphic to) the category Rel with sets as objects, and binary relations as arrows. The correspondence mentioned then links the natural transformation λ^T to the notion of relation lifting $\overline{T}$.

Lemma 2.11. Let T be a standard, finitary, weak pullback-preserving functor. Let X be some set and let $\Phi \in TPX$. Then:

(1) If $\varnothing \in Base^T(\Phi)$, then $\lambda^T(\Phi) = \varnothing$.
(2) If $Base^T(\Phi)$ consists of singletons only, then $\lambda^T(\Phi)$ is a singleton.
(3) If T maps finite sets to finite sets, then for all $\Phi \in TP_\omega X$, $|\lambda^T(\Phi)| < \omega$.

Proof.

(1) We suppose that α is a lifted member of Φ. So we may derive by Fact 2.8 that $Base^T(\alpha) \; \overline{P} \in Base^T(\Phi)$. But from this it would follow that if $\varnothing \in Base^T(\Phi)$, then $Base^T(\alpha)$ contains a member of $\varnothing$, which is clearly impossible. Hence, $\lambda^T(\Phi)$ is empty.

(2) Observe that another way of saying that $Base^T(\Phi)$ consists of singletons only, is that $\Phi \in TS_X$, with $S_X := \{\{x\} \mid x \in X\}$. Let $\theta_X : S_X \to X$ be the inverse of η_X, that is, θ_X is the bijection mapping a singleton $\{x\}$ to its unique member x. Clearly, we then have $(Gr\,\theta_X)^{\smile} = \in\!\restriction_{X \times S_X}$, from which it follows by Fact 2.6 that $(Gr\,T\theta_X)^{\smile} = \overline{T}\in\!\restriction_{TX \times TS_X}$. From this it is immediate that if $\Phi \in TS_X$, then $(T\theta_X)(\Phi)$ is the unique lifted member of Φ.

(3) Since T is finitary, $\Phi \in TP_\omega X$ implies that $\Phi \in TP_\omega Y$ for some finite set Y, and from this it follows that $Base^T(\Phi) \subseteq P_\omega Y$. If α is a lifted member of Φ, then by Fact 2.8, we obtain $Base^T(\alpha) \; \overline{P} \in Base^T(\Phi)$, and thus, in particular, we find $Base^T(\alpha) \subseteq \bigcup Base^T(\Phi) \subseteq Y$. From this it follows that $\lambda^T(\Phi) \subseteq TY$, so $\lambda^T(\Phi)$ must be finite by the assumption on T. $\qquad\square$

2.4. *Frames and their presentations*

A *frame* is a complete lattice in which finite meets distribute over arbitrary joins. The signature of frames consists of arbitrary joins and finite meets, and it will be convenient

for us to include the top and bottom as well. Thus a frame will usually be given as $\mathbb{L} = \langle L, \bigvee, \bigwedge, 0, 1 \rangle$, while we will often consider join and meet as functions $\bigvee_{\mathbb{L}} : PL \to L$ and $\bigwedge_{\mathbb{L}} : P_\omega L \to L$. This enables us, for instance, to define a frame homomorphism $f : \mathbb{L} \to \mathbb{M}$ as a map from L to M satisfying

$$f \circ \bigwedge = \bigwedge \circ (P_\omega f)$$
$$f \circ \bigvee = \bigvee \circ (Pf).$$

We use Fr to denote the category of frames and frame homomorphisms. The initial frame (the lattice of truth values) will be denoted by Ω, and for a given frame $\mathbb{L}$ we will use $!_{\mathbb{L}}$ to denote the unique frame homomorphism from Ω to $\mathbb{L}$, but omitting the subscript when $\mathbb{L}$ is clear from context.

The order relation $\leqslant_{\mathbb{L}}$ of a frame $\mathbb{L}$ is given by $a \leqslant_{\mathbb{L}} b$ if $a \wedge b = a$ (or, equivalently, $a \vee b = b$). We can adjoin an implication operation to a frame $\mathbb{L}$ by defining $a \to b := \bigvee \{c \mid a \wedge c \leqslant b\}$; this operation turns $\mathbb{L}$ into a Heyting algebra. As a special case of implication, we can consider the *negation*:

$$\neg a := \bigvee \{c \mid a \wedge c = 0\}.$$

Generally, neither of these two operations is preserved by frame homomorphisms. A subset S of $\mathbb{L}$ is *directed* if for every $s_0, s_1 \in S$ there is an element $s \in S$ such that $s_0, s_1 \leqslant s$. The join of a directed set S is often denoted by $\bigvee^{\uparrow} S$.

A *frame presentation* is a tuple $\langle G \mid R \rangle$ where G is a set of generators and $R \subseteq PP_\omega G \times PP_\omega G$ is a set of relations. A presentation $\langle G \mid R \rangle$ *presents* a frame $\mathbb{L}$ if there exists a function $f : G \to L$ that is *compatible with R*, that is, such that

$$\text{for all } (t_1, t_2) \in R, \ \bigvee_{A \in t_1} \bigwedge (P_\omega f) A = \bigvee_{B \in t_2} \bigwedge (P_\omega f) B,$$

and for all frames $\mathbb{M}$ and functions $g : G \to M$ compatible with R, there is a unique frame homomorphism $g' : \mathbb{L} \to \mathbb{M}$ such that $g'f = g$. We call f the *insertion of generators* (of G in $\mathbb{L}$).

Fact 2.12. Every frame presentation presents a frame.

The details of the proof of the above fact (which can be found in Vickers (1989, Section 4.4)) tell us how to construct a unique frame given a presentation $\langle G \mid R \rangle$. Omitting these details of the construction, we denote this unique frame by $\text{Fr}\langle G \mid R \rangle$. We will usually write $\bigvee_{i \in I} \bigwedge A_i = \bigvee_{j \in J} \bigwedge B_j$ instead of $(\{A_i \mid i \in I\}, \{B_j \mid j \in J\})$ when specifying relations. In light of the fact that $a \leqslant b$ if and only if $a \vee b = b$, we will also allow ourselves the liberty to specify inequalities of the shape $\bigvee_{i \in I} \bigwedge A_i \leqslant \bigvee_{j \in J} \bigwedge B_j$ as relations. It follows from the proof of Fact 2.12 that if $f : G \to \text{Fr}\langle G \mid R \rangle$ is the insertion of generators, then every element of $\text{Fr}\langle G \mid R \rangle$ can be written as $\bigvee_{i \in I} \bigwedge P_\omega f A$ for some $\{A_i \mid i \in I\} \in PP_\omega G$; in other words, every element of $\text{Fr}\langle G \mid R \rangle$ can be written as an infinite disjunction of finite conjunctions of generators.

We will now introduce flat-site presentations for frames, which have as one of their main advantages that they allow us to assume that an arbitrary element of the frame being presented is an infinite join of generators. A *flat site* is a triple $\langle X, \sqsubseteq, \lhd_0 \rangle$, where

$\langle X, \sqsubseteq \rangle$ is a pre-order and $\lhd_0 \subseteq X \times PX$ is a binary relation such that for all $b \sqsubseteq a \lhd_0 A$, there exists $B \subseteq \downarrow A \cap \downarrow b$ such that $b \lhd_0 B$. A flat-site $\langle X, \sqsubseteq, \lhd_0 \rangle$ presents a frame $\mathbb{L}$ if there exists a function $f : X \to L$ such that

— f is order-preserving,
— $1 \leqslant \bigvee (Pf)X$,
— for all $a, b \in X$, $f(a) \wedge f(b) \leqslant \bigvee (Pf)(\downarrow a \cap \downarrow b)$, and
— for all $a \lhd_0 A$, $f(a) \leqslant \bigvee (Pf)A$

and for all frames $\mathbb{M}$ and all $g : X \to M$ satisfying the above two properties, there exists a unique frame homomorphism $g' : \mathbb{L} \to \mathbb{M}$ such that $g' \circ f = g$. Specifically, for all $a \in \mathbb{L}$,

$$g'(a) = \bigvee \{ g(x) \mid f(x) \leqslant a \}.$$

To put this another way, the frame presented by a flat site is

$$
\begin{aligned}
\mathrm{Fr}\langle X, \sqsubseteq, \lhd_0 \rangle \simeq \mathrm{Fr}\langle X \mid & a \leqslant b \quad (a \sqsubseteq b), \\
& a \leqslant \bigvee A \quad (a \lhd_0 A), \\
& 1 = \bigvee X \\
& a \wedge b = \bigvee \{ c \mid c \sqsubseteq a, c \sqsubseteq b \} \rangle.
\end{aligned}
$$

A *suplattice* is a complete $\bigvee$-semilattice, so a suplattice homomorphism is a map that preserves $\bigvee$. A *suplattice presentation* is a triple $\langle X, \sqsubseteq, \lhd_0 \rangle$, where $\langle X, \sqsubseteq \rangle$ is a pre-order and $\lhd_0 \subseteq X \times PX$. A suplattice presentation $\langle X, \sqsubseteq, \lhd_0 \rangle$ presents a suplattice $\mathbb{L}$ if there exists a function $f : X \to L$ such that:

— f is order-preserving, and
— for all $a \lhd_0 A$, $f(a) \leqslant \bigvee Pf(A)$,

and for all suplattices $\mathbb{M}$ and all functions $g : X \to M$ respecting the above two conditions, there exists a unique suplattice homomorphism $g' : \mathbb{L} \to \mathbb{M}$ such that $g' \circ f = g$. Every suplattice presentation presents a suplattice (Jung *et al.* 2008, Proposition 2.5). Now observe that every flat site can also be viewed as a suplattice presentation with an additional stability condition. Consequently, given a flat site $\langle X, \sqsubseteq, \lhd_0 \rangle$, we can generate two different objects with it: a frame $\mathrm{Fr}\langle X, \sqsubseteq, \lhd_0 \rangle$; and a suplattice $\mathrm{SupLat}\langle X, \sqsubseteq, \lhd_0 \rangle$. The Flat-Site Coverage Theorem (Vickers 2006, Theorem 5) tells us that these two objects are in fact order isomorphic.

Fact 2.13. Let $\langle X, \sqsubseteq, \lhd_0 \rangle$ be a flat site. Then $\mathrm{Fr}\langle X, \sqsubseteq, \lhd_0 \rangle \simeq \mathrm{SupLat}\langle X, \sqsubseteq, \lhd_0 \rangle$.

We will just state the following consequences of the above fact. Suppose $\langle X, \sqsubseteq, \lhd_0 \rangle$ is a flat site that presents a frame $\mathbb{L}$ via $f : X \to L$. Then:

— Every element of $\mathbb{L}$ is of the shape $\bigvee Pf(A)$ for some $A \in PX$.
— We can use $\langle X, \sqsubseteq, \lhd_0 \rangle$ both to define suplattice homomorphisms and frame homomorphisms.

2.5. *Powerlocales through $\square$ and $\diamond$*

We will now introduce the Vietoris powerlocale. In line with our generally algebraic approach, we shall define it directly as a functor on the category of frames rather than

through its opposite, the category of locales. In its full generality, it originates (as the 'Vietoris construction') in Johnstone (1985), with some earlier, more restricted references in Johnstone (1982). For locales it is a localic analogue of hyperspace (with Vietoris topology). The points are (in bijection with) certain sublocales of the original locale. For a full constructive description, see Vickers (1997).

Given a frame $\mathbb{L}$, we first define $L_\square := L$ and $L_\diamond := L$, and then

$$V\mathbb{L} := \mathrm{Fr}\langle L_\square \oplus L_\diamond \mid \square 1 = 1$$
$$\square(a \wedge b) = \square a \wedge \square b$$
$$\square(\textstyle\bigvee^\uparrow A) = \textstyle\bigvee^\uparrow_{a \in A} \square a \quad (A \in PL \text{ directed})$$
$$\diamond(\textstyle\bigvee A) = \textstyle\bigvee_{a \in A} \diamond a \quad (A \in PL)$$
$$\square a \wedge \diamond b \leqslant \diamond(a \wedge b)$$
$$\square(a \vee b) \leqslant \square a \vee \diamond b\rangle.$$

Remark 2.14. We have abused the notation when specifying the relations in the above definition. Strictly speaking, we have two maps, $\square : L_\square \to V\mathbb{L}$ for the left copy of $\mathbb{L}$ and $\diamond : L_\diamond \to V\mathbb{L}$ for the right copy of $\mathbb{L}$, so the insertion of generators is the map $\square \oplus \diamond : L_\square \oplus L_\diamond \to V\mathbb{L}$.

Johnstone (1985) shows that V gives a monad on the category of locales, that is, a comonad on the category of frames. We shall not need the full strength of this here, but some of the ingredients of the comonad structure are easy to check:

— V is functorial:
  If $f : \mathbb{L} \to \mathbb{M}$ is a frame homomorphism, then the function $(\square f) \oplus (\diamond f) : L_\square \oplus L_\diamond \to VM$ is compatible with the relations in the presentation of $V\mathbb{L}$, so there is a frame homomorphism $Vf : V\mathbb{L} \to V\mathbb{M}$ extending this map. It is also easy to show functoriality.
— The counit $i_\mathbb{L} : V\mathbb{L} \to \mathbb{L}$ is given by $\square a \mapsto a$ and $\diamond a \mapsto a$:
  The comultiplication $\mu_\mathbb{L} : V\mathbb{L} \to VV\mathbb{L}$ is given by $\square a \mapsto \square\square a$ and $\diamond a \mapsto \diamond\diamond a$.

3. The T-powerlocale construction

In this section we arrive at the main conceptual contribution of this paper. Given a weak pullback-preserving, standard, finitary functor $T : \mathsf{Set} \to \mathsf{Set}$, we define its associated T-powerlocale functor $V_T : \mathsf{Fr} \to \mathsf{Fr}$ on the category of frames using the Carioca axioms for coalgebraic modal logic. This construction truly generalises the Vietoris powerlocale construction because, as we will see, the P_ω-powerlocale is isomorphic to the Vietoris powerlocale. The other two major results in this section are the fact that we can lift a natural transformation between transition functors $\rho : T' \to T$ to a natural transformation $\widehat{\rho} : V_T \to V_{T'}$ going in the other direction, and the fact that T-powerlocales are join-generated by their generators of the shape $\nabla\alpha$. We will establish the latter fact using the stronger result by showing that $V_T\mathbb{L}$ admits a flat-site presentation. The fact that $V_T\mathbb{L}$ is join-generated by its generators is not entirely surprising since the Carioca axioms were designed with the desirability of conjunction-free disjunctive normal forms in mind

(Bílková *et al.* 2008). However, the precise mathematical formulation of this property, using flat sites and suplattices, is an improvement over what was previously known.

This section is organised as follows:

— In Section 3.1, we introduce the T-powerlocale construction on frames.
— In Section 3.2, we make technical observations about T-powerlocales.
— In Section 3.3, we consider two instantiations of the T-powerlocale construction, the most notable of which is the P_ω-powerlocale, which is isomorphic to the classical Vietoris powerlocale.
— In Section 3.4 we extend the T-powerlocale construction to a functor V_T on the category of frames, and we show how to lift natural transformations between set functors T, T' to natural transformations between powerlocale functors V_T, $V_{T'}$.
— Finally, in Section 3.5, we show that the T-powerlocale construction admits a flat-site presentation, a corollary of which is that each element of $V_T\mathbb{L}$ has a disjunctive normal form.

3.1. *Introducing the T-powerlocale*

In this section, we will use the Carioca axioms for coalgebraic modal logic (Bílková *et al.* 2008) to define the T-powerlocale $V_T\mathbb{L}$ of a given frame $\mathbb{L}$ using a frame presentation, that is, using generators and relations. The generators of $V_T\mathbb{L}$ will be given by the set TL; in order to specify the relations, we will use *relation lifting* (Section 2.3) and *slim redistributions*, which we will introduce below. In addition, we will provide an alternative presentation of $V_T\mathbb{L}$, which does not use slim redistributions. From a conceptual viewpoint, it is not immediately obvious which presentation of $V_T\mathbb{L}$ should be taken as the primary definition. Our choice to use slim redistributions in the primary definition is motivated by the precedent set by the existing literature (Bílková *et al.* 2008; Kupke *et al.* 2008; Kupke *et al.* 2010).

Definition 3.1. Let $T : \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor, X be a set and $\Gamma \in P_\omega TX$. The set of all *slim redistributions* of Γ is defined as follows:

$$SRD(\Gamma) = \left\{ \Psi \in TP_\omega \left(\bigcup_{\gamma \in \Gamma} Base^T(\gamma) \right) \mid \forall \gamma \in \Gamma, \gamma \, \overline{T} \in \Psi \right\}.$$

Intuitively, $\Psi \in TP_\omega X$ is a slim redistribution of $\Gamma \in P_\omega TX$ if:

(i) Ψ is 'obtained from the material of Γ', that is:

$$\Psi \in TP_\omega \left(\bigcup_{\gamma \in \Gamma} Base^T(\gamma) \right).$$

(ii) Every element of Γ is a lifted member of Ψ, or equivalently, $\Gamma \subseteq \lambda^T(\Psi)$.

We will illustrate this definition through the motivating example of slim redistributions, namely slim redistribution for the finite powerset functor.

Example 3.2. Recall from Example 2.4 that if $R \subseteq X \times Y$ is a relation, then $\overline{P_\omega}R \subseteq P_\omega X \times P_\omega Y$ can be characterised as follows:

$$\alpha \overline{P_\omega} R \beta \text{ iff } \forall x \in \alpha, \exists y \in \beta, xRy \text{ and } \forall y \in \beta, \exists x \in \alpha, xRy.$$

In particular, for $\in\; \subseteq X \times PX$, we get $\alpha\, \overline{P_\omega} \in \Gamma$ if and only if $\alpha \subseteq \bigcup\Gamma$ and $\forall \gamma \in \Gamma$, $\gamma \between \alpha$. (Recall that $\gamma \between \alpha$ means that $\gamma \cap \alpha$ is inhabited.) For an order $\leqslant$, we define the *upper, lower* and *convex* pre-orders on finite sets as follows:

$$\alpha \leqslant_L \beta \text{ if } \alpha \subseteq\; \downarrow\beta, \text{ that is, } \forall x \in \alpha,\, \exists y \in \beta,\, x \leqslant y$$

$$\alpha \leqslant_U \beta \text{ if } \uparrow\alpha \supseteq \beta, \text{ that is, } \forall y \in \beta,\, \exists x \in \alpha,\, x \leqslant y$$

$$\alpha \leqslant_C \beta \text{ if } \alpha \leqslant_L \beta \text{ and } \alpha \leqslant_U \beta.$$

Thus $\overline{P_\omega} \leqslant$ is $\leqslant_C$.

Next, if $\alpha \in P_\omega S$, then

$$Base(\alpha) = \bigcap\{S' \in P_\omega(S) \mid \alpha \subseteq S'\} = \alpha.$$

From this, we have if $\Gamma \in P_\omega P_\omega X$, then

$$SRD(\Gamma) = \{\Psi \in P_\omega P_\omega \left(\bigcup\Gamma\right) \mid \forall\gamma \in \Gamma,\, \left(\gamma \subseteq \bigcup\Psi \text{ and } \forall\alpha \in \Psi,\, \alpha \between \gamma\right)$$
$$= \{\Psi \in P_\omega P_\omega (X) \mid \bigcup\Psi = \bigcup\Gamma \text{ and } \forall\gamma \in \Gamma,\, \forall\alpha \in \Psi,\, \alpha \between \gamma\}.$$

Definition 3.3. Let T be a standard, finitary, weak pullback-preserving functor and $\mathbb{L}$ be a frame. We define the *T-powerlocale of* $\mathbb{L}$ by

$$V_T\mathbb{L} := \mathrm{Fr}\langle TL \mid (\nabla 1), (\nabla 2), (\nabla 3)\rangle,$$

where the relations are the *Carioca axioms* (Bílková *et al.* 2008):

$(\nabla 1)$	$\nabla\alpha \leqslant \nabla\beta$	$(\alpha\, \overline{T}{\leqslant}\, \beta)$
$(\nabla 2)$	$\bigwedge_{\alpha\in\Gamma}\nabla\alpha \leqslant \bigvee\{\nabla(T\bigwedge)\Psi \mid \Psi \in SRD(\Gamma)\}$	$(\Gamma \in P_\omega TL)$
$(\nabla 3)$	$\nabla(T\bigvee)\Phi \leqslant \bigvee\{\nabla\beta \mid \beta\, \overline{T}{\in}\, \Phi\}$	$(\Phi \in TPL)$.

Remark 3.4. To be precise, we assume that $\nabla: TL \to V_T L$ is the insertion of generators, so when specifying the relations we should write, for example, $\alpha \leqslant \beta$ rather than $\nabla\alpha \leqslant \nabla\beta$. The way we have specified the relations above is more consistent with Bílková *et al.* (2008).

We will discuss the instantiation of these axioms for $T = P_\omega$ in more detail in Section 3.3.

We will now present a very useful equivalent definition of $V_T\mathbb{L}$. The crucial observation behind the alternative definition of $V_T\mathbb{L}$ is the following technical lemma, which characterises the slim redistributions of a given finite subset Γ of $\langle TL, \overline{T}{\leqslant}\rangle$ as the maximal lower bounds of Γ. Observe that the lemma also holds when $\Gamma = \varnothing$.

Lemma 3.5. Let $T: \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor, $\mathbb{L}$ be a meet-semilattice (for example, a frame) and $\Gamma \in P_\omega TL$. Then for any $\alpha \in TL$, the following are equivalent:

(a) $\alpha \in TL$ is a lower bound of Γ, that is, $\alpha\, \overline{T}{\leqslant}\, \gamma$ for all $\gamma \in \Gamma$.
(b) $\alpha\, \overline{T}{\leqslant}\, (T\bigwedge)\Phi$ for some $\Phi \in SRD(\Gamma)$.

In particular, if $\Phi \in SRD(\Gamma)$, then $(T\bigwedge)\Phi\, \overline{T}{\leqslant}\, \gamma$ for all $\gamma \in \Gamma$.

Proof. Recall that

$$SRD(\Gamma) := \left\{ \Psi \in TP\left(\bigcup_{\gamma \in \Gamma} Base^T(\gamma)\right) \mid \Gamma \subseteq \lambda^T(\Psi) \right\}.$$

For the implication from (b) to (a), observe that for any $a \in L$ and $A \in P_\omega L$, we have that $a \in A$ implies that $\bigwedge A \le a$. By Fact 2.6, it follows that for all $\gamma \in TL$ and $\Psi \in TP_\omega L$, if $\gamma \ \overline{T} \in \Psi$, then $T\bigwedge(\Psi) \ \overline{T}\le \gamma$. We now suppose Ψ is a slim redistribution of Γ. Then $\Gamma \subseteq \lambda^T(\Psi)$, so $(T\bigwedge)\Psi$ is a $\overline{T}\le$-lower bound of Γ. From this, the implication (b) $\Rightarrow$ (a) is immediate.

For the opposite implication, take $\alpha \in TL$ such that $\forall \gamma \in \Gamma, \alpha \ \overline{T}\le \gamma$. Then by Fact 2.8, we obtain $Base^T(\alpha) \ \overline{P}\le Base^T(\gamma)$ for all $\gamma \in \Gamma$. Using the abbreviation $C := \bigcup_{\gamma \in \Gamma} Base^T(\gamma)$ and defining $f : Base^T(\alpha) \to PC$ by

$$f : a \mapsto \uparrow_L a \cap C,$$

that is, $f(a) = \{c \in C \mid a \le c\}$, we get that Tf is a function

$$Tf : T\, Base^T(\alpha) \to TPC.$$

We claim that $\Psi := Tf(\alpha)$ is an element of $SRD(\Gamma)$ and that $\alpha \ \overline{T}\le T\bigwedge(\Psi)$. For the first claim, since $\Psi \in TPC$, all we need to show is that $\Gamma \subseteq \lambda^T(\Psi)$, that is, that for all $\gamma \in \Gamma$, $\gamma \ \overline{T} \in \Psi$. So suppose that $\gamma \in \Gamma$. Then, by assumption, $\alpha \ \overline{T}\le \gamma$, so $Base^T(\alpha) \ \overline{P}\le Base^T(\gamma)$. It follows from the definition of f that for all $b \in Base^T(\gamma)$ and all $a \in Base^T(\alpha)$, if $a \le b$, then $b \in f(a)$. It then follows by Fact 2.6 that

$$\forall \delta \in T\, Base^T(\alpha), \forall \beta \in T\, Base^T(\gamma), \delta \ \overline{T}\le \beta \Rightarrow \beta \ \overline{T} \in Tf(\delta).$$

So, in particular, since $\alpha \in T\, Base^T(\alpha)$, $\gamma \in T\, Base^T(\gamma)$ and $\alpha \ \overline{T}\le \gamma$, we have $\gamma \ \overline{T} \in Tf(\alpha) = \Psi$. Since $\gamma \in \Gamma$ was arbitrary, it follows that $\Gamma \subseteq \lambda^T(\Psi)$. Consequently, $\Psi \in SRD(\Gamma)$, as we wanted to show.

For the second claim, that is, that $\alpha \ \overline{T}\le T\bigwedge(\Psi)$, it suffices to observe that $a \le \bigwedge f(a)$ for all $a \in Base^T(\alpha)$, so by Fact 2.6,

$$\forall \delta \in T\, Base^T(\alpha), \delta \ \overline{T}\le T\bigwedge \circ Tf(\delta).$$

Since $\alpha \in T\, Base^T(\alpha)$ and $\Psi = Tf(\alpha)$, we get that $\alpha \ \overline{T}\le T\bigwedge \circ Tf(\alpha) = T\bigwedge(\Psi)$. $\square$

Corollary 3.6. Let $T : \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor and $\mathbb{L}$ be a frame. Then

$$V_T \mathbb{L} \simeq \mathrm{Fr}\langle TL \mid (\nabla 1), (\nabla 2'), (\nabla 3) \rangle,$$

where the relations are as follows:

$(\nabla 1)$	$\nabla\alpha \le \nabla\beta$	$(\alpha \ \overline{T}\le \beta)$
$(\nabla 2')$	$\bigwedge_{\gamma \in \Gamma} \nabla\gamma \le \bigvee\{\nabla\alpha \mid \forall \gamma \in \Gamma, \alpha \ \overline{T}\le \gamma\}$	$(\Gamma \in P_\omega TL)$
$(\nabla 3)$	$\nabla(T\bigvee)\Phi \le \bigvee\{\nabla\beta \mid \beta \ \overline{T} \in \Phi\}$	$(\Phi \in TPL).$

Proof. Observe that the only difference between $\mathrm{Fr}\langle TL \mid (\nabla 1), (\nabla 2'), (\nabla 3)\rangle$ and the original definition of $V_T \mathbb{L}$ is that we have replaced $(\nabla 2)$,

$(\nabla 2)$ $\qquad\qquad \bigwedge_{\alpha \in \Gamma} \nabla \alpha \leqslant \bigvee \{\nabla (T \bigwedge) \Psi \mid \Psi \in SRD(\Gamma)\} \qquad\qquad (\Gamma \in P_\omega TL),$

by $(\nabla 2')$. To see that the equivalence of these two relations is an immediate corollary of Lemma 3.5, take any $\Gamma \in TP_\omega L$, then

$$\bigvee \{\nabla T \bigwedge (\Psi) \mid \Psi \in SRD(\Gamma)\}$$
$$= \bigvee \{\nabla \alpha \mid \exists \Psi \in SRD(\Gamma), \alpha \, \overline{T} \leqslant \nabla T \bigwedge (\Psi)\} \quad \text{(by order theory and } (\nabla 1))$$
$$= \bigvee \{\nabla \alpha \mid \forall \gamma \in \Gamma, \alpha \, \overline{T} \leqslant \gamma\} \qquad\qquad\qquad \text{(by Lemma 3.5)}.$$

It follows that $V_T \mathbb{L} \simeq \mathrm{Fr}\langle TL \mid (\nabla 1), (\nabla 2'), (\nabla 3)\rangle$. $\qquad\qquad\qquad\qquad\qquad \square$

Remark 3.7. We will see later that both of the axioms $(\nabla 2)$ and $(\nabla 2')$ are equally useful. It seems that $(\nabla 2')$ has not been studied before in the literature on coalgebraic modal logic using the ∇-modality (Palmigiano and Venema 2007; Bílková *et al.* 2008; Kissig and Venema 2009; Kupke *et al.* 2010).

3.2. *Basic properties of the T-powerlocale*

In this section we make some technical observations about slim redistributions and about the structure of the T-powerlocale. We start with two facts about slim redistributions.

Lemma 3.8. Let $T : \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor. Then $SRD(\varnothing) = T\{\varnothing\}$.

Proof. If Φ is a slim redistribution of the empty set, then, by definition, $\Phi \in TP_\omega(\varnothing) = T\{\varnothing\}$. Conversely, any $\Phi \in T\{\varnothing\}$ satisfies the condition that $\varnothing \subseteq \lambda^T(\Phi)$, so $\Phi \in SRD(\varnothing)$. $\qquad\qquad\qquad\qquad \square$

The following Lemma plays an essential role when defining V_T on frame homomorphisms, rather than just on frames. It is crucial when showing that if $f : \mathbb{L} \to \mathbb{M}$ is a frame homomorphism, then $V_T f : V_T \mathbb{L} \to V_T \mathbb{M}$ preserves conjunctions, as we will see in Section 3.4.

Lemma 3.9. Let $T : \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor, X, Y be sets, $f : X \to Y$ be a function and $\Gamma \in P_\omega TX$. Then the restriction of $TP_\omega f : TP_\omega X \to TP_\omega Y$ to $SRD(\Gamma)$ is a surjection onto $SRD(P_\omega Tf\Gamma)$.

Proof. Let X, Y, f and Γ be as in the statement of the Lemma, and abbreviate

$$\Gamma' := (P_\omega Tf)\Gamma$$
$$C := \bigcup_{\gamma \in \Gamma} Base^T(\gamma)$$
$$C' := \bigcup_{\gamma' \in \Gamma'} Base^T(\gamma').$$

Then an easy calculation shows that

$$
\begin{aligned}
C' &= \bigcup_{\gamma \in \Gamma} Base^T (Tf)(\gamma) && \text{(definition of } \Gamma') \\
&= \bigcup_{\gamma \in \Gamma} (Pf) \, Base^T(\gamma) && (Base^T \text{ is natural transformation}) \\
&= (Pf)(C) && \text{(elementary set theory).}
\end{aligned}
$$

We will first show that $TP_\omega f$ maps slim redistributions of Γ to slim redistributions of Γ'. To do this, we take an arbitrary element $\Phi \in SRD(\Gamma)$, and write $\Phi' := (TP_\omega f)\Phi$. We claim that $\Phi' \in SRD(\Gamma')$, and first show that

$$
\Phi' \in TP_\omega C', \tag{1}
$$

or, equivalently, that $Base^T \Phi' \subseteq P_\omega C'$. To prove this inclusion, we take an arbitrary set $A' \in Base^T(\Phi')$. Since by Fact 2.8, $Base^T(\Phi') = (P_\omega P_\omega f)(Base^T(\Phi))$, this means that A' must be of the form $(P_\omega f)(A)$ for some $A \in Base^T(\Phi)$. In particular, A' must be a subset of $(P_\omega f)(\bigcup Base^T(\Phi))$. Also, because Φ is a slim redistribution of Γ, by definition, we have $Base^T(\Phi) \subseteq P_\omega C$, so $\bigcup Base^T(\Phi) \subseteq \bigcup C$. From this it follows that

$$
A' \subseteq (Pf)\left(\bigcup Base^T(\Phi)\right) \subseteq (Pf)\left(\bigcup C\right) = C',
$$

as required.

Our second claim is that

$$
\Gamma' \subseteq \lambda^T(\Phi'). \tag{2}
$$

To prove this, we take an arbitrary element of Γ', say, $(Tf)\gamma$ for some $\gamma \in \Gamma$. We have $\gamma \; \overline{T} \in \Phi$ by the assumption that $\Phi \in SRD(\Gamma)$. But then, since $a \in A$ implies $fa \in (P_\omega f)A$ for any $a \in C$ and $A \subseteq C$, it follows by Lemma 2.7 that

$$
\gamma' = (Tf)\gamma \; \overline{T} \in (TP_\omega f)(\Phi) = \Phi'.
$$

This means that γ' is a lifted member of Φ', as required.

Clearly, the claims (1) and (2) above suffice to prove that $\Phi' \in SRD(\Gamma')$, which means that, indeed, $TP_\omega f$ maps slim redistributions of Γ to slim redistributions of Γ'.

Thus it is just left to prove that every slim redistribution of Γ' is of the form $(TP_\omega f)\Phi$ for some slim redistribution Φ of Γ. Take an arbitrary $\Phi' \in SRD(\Gamma')$, and recall that $\check{P}$ denotes the contravariant power set functor. We restrict f to the map $f^- : C \to C'$, which means that $\check{P}f^- : P_\omega C' \to P_\omega C$. It follows that $T\check{P}f^- : TP_\omega C' \to TP_\omega C$, so we may define $\Phi := (T\check{P}f^-)\Phi'$, and obtain $\Phi \in TP_\omega C$. Hence, in order to prove that

$$
\Phi \in SRD(\Gamma), \tag{3}
$$

it suffices to show that $\Gamma \subseteq \lambda^T(\Phi)$. But this is an immediate consequence of the fact that λ^T is a distributive law of T over $\check{P}$ (Fact 2.9), since for an arbitrary $\gamma \in \Gamma$, we may reason as follows. From $\gamma \in \Gamma$, it follows by the definition of Γ' that $(Tf^-)(\gamma) = (Tf)(\gamma)$ belongs to Γ'. Since $\Gamma' \subseteq \lambda_Y^T(\Phi)$ by assumption, by the definition of $\check{P}$, we find that

$\gamma \in (\check{P} T f) \lambda_Y^T(\Psi)$. But by $\lambda^T : T\check{P} \to \check{P}T$, we know that

$$(\check{P} T f) \lambda_Y^T(\Psi) = \lambda_X^T (T \check{P} f)(\Psi) = \lambda_X^T(\Phi).$$

Thus we find $\gamma \in \lambda^T(\Phi)$, as required.

Finally, note that $f^- : C \to C'$ is surjective, so it follows by properties of the covariant and contravariant power set functors that $P_\omega f^- \circ \check{P} f^- = id_{P_\omega C'}$. From this it is immediate by the functoriality of T that

$$\Phi' = (T P_\omega f^- \circ T \check{P} f^-)\Phi' = (T P_\omega f^-)\Phi = (T P_\omega f)\Phi.$$

This concludes the proof of the Lemma. $\qquad\square$

In the following lemma we gather together some basic observations on the frame structure of the T-powerlocale. These facts generalise results from Kupke *et al.* (2010) to our geometrical setting.

Lemma 3.10. Let T be a standard, finitary, weak pullback-preserving functor and let $\mathbb{L}$ be a frame. Then:

(1) If $\alpha \in T L$ is such that $0_\mathbb{L} \in Base^T(\alpha)$, then $\nabla\alpha = 0_{V_T\mathbb{L}}$.
(2) If $A \subseteq L$ is such that $a \wedge b = 0_\mathbb{L}$ for all $a \neq b$ in A, then $\nabla\alpha \wedge \nabla\beta = 0_{V_T\mathbb{L}}$ for all $\alpha \neq \beta$ in $T A$.
(3) If there is no relation R such that $\alpha \, \overline{T R} \, \beta$, then $\nabla\alpha \wedge \nabla\beta = 0_{V_T\mathbb{L}}$.
(4) $1_{V_T\mathbb{L}} = \bigvee\{\nabla\gamma \mid \gamma \in T\{1_\mathbb{L}\}\}$.
(5) For any $A \subseteq L$ such that $1_\mathbb{L} = \bigvee A$, we have $1_{V_T\mathbb{L}} = \bigvee\{\nabla\alpha \mid \alpha \in T A\}$.

Proof.

(1) Let $\alpha \in T L$ be such that $0_\mathbb{L} \in Base^T(\alpha)$. Consider the map $f : L \to P L$ given by

$$f(a) := \begin{cases} \varnothing & \text{if } a = 0_\mathbb{L} \\ \{a\} & \text{if } a > 0_\mathbb{L}. \end{cases}$$

Then $id_L = \bigvee \circ f$, so $id_{TL} = (T \bigvee) \circ (T f)$ by the functoriality of T. In particular, we get that $\alpha = (T \bigvee)(T f)(\alpha)$, so we may calculate

$$\nabla\alpha = \bigvee \{\nabla\beta \mid \beta \, \overline{T} \in (T f)(\alpha)\} \qquad \text{(axiom } \nabla 2)$$

$$\leqslant \bigvee \{\nabla\beta \mid Base^T(\beta) \, \overline{P} \in Base^T((T f)(\alpha))\} \qquad \text{(Fact 2.8(2))}$$

$$= \bigvee \varnothing \qquad \text{(†– see below)}$$

$$= 0_{V_T\mathbb{L}}.$$

In order to justify the remaining step (†) in this calculation, note that from the naturality of $Base^T$ (Fact 2.8(1)) we have

$$Base^T((T f)(\alpha)) = (P f)(Base^T(\alpha)),$$

so, by the assumption that $0_\mathbb{L} \in Base^T(\alpha)$, we obtain $\varnothing \in Base^T((T f)(\alpha))$. Now suppose in order to show a contradiction that there is some $B \subseteq L$ such that $B \, \overline{P} \in Base^T((T f)(\alpha))$. Then by the definition of $\overline{P}$, there is a $b \in B$ such that $b \in \varnothing$, which provides the desired contradiction. This proves (†), and concludes the proof of part 1.

(2) Let $A \subseteq L$ be such that $a \wedge b = 0_{\mathbb{L}}$ for all $a \neq b$ in A, and take two distinct elements $\alpha, \beta \in TA$. In order to prove that $\nabla\alpha \wedge \nabla\beta = 0_{V_T\mathbb{L}}$, it will be enough by axiom ($\nabla 2$) to show that

$$\nabla(T\textstyle\bigwedge)(\Phi) = 0_{V_T\mathbb{L}}, \text{ for all } \Phi \in SRD\{\alpha, \beta\}. \tag{4}$$

Take an arbitrary slim redistribution Φ of $\{\alpha, \beta\}$. Then by Fact 2.11, $Base^T(\Phi)$ contains a set $A_0 \subseteq_\omega A$ of size > 1. Define the map $d : Base^T(\Phi) \to P_\omega(A) \cup \{\{1_{\mathbb{L}}\}\}$ by putting

$$d(B) := \begin{cases} \varnothing & \text{if } |B| > 1 \\ B & \text{if } |B| = 1 \\ \{1_{\mathbb{L}}\} & \text{if } |B| = 0. \end{cases}$$

It is straightforward to verify from the assumptions on A and the definition of d, that $\bigwedge B \leqslant \bigvee d(B)$ for each $B \in Base^T(\Phi)$. Hence, it follows by Fact 2.6 that

$$\left(T\textstyle\bigwedge\right)(\Phi) \, \overline{T}{\leqslant} \left(T\textstyle\bigvee\right)(Td)(\Phi),$$

so by axiom ($\nabla 1$), we may conclude that

$$\nabla(T\textstyle\bigwedge)(\Phi) \leqslant \nabla(T\textstyle\bigvee)(Td)(\Phi). \tag{5}$$

Finally, it follows from the naturality of $Base^T$ (Fact 2.8(1)) that

$$Base^T(Td)(\Phi) = (Pd)(Base^T(\Phi)).$$

Consequently, for the set $A_0 \in Base^T(\Phi)$ satisfying $|A_0| > 1$, we find

$$\varnothing = d(A_0) \in Base^T(Td)(\Phi),$$

and then

$$0_{\mathbb{L}} = \textstyle\bigvee \varnothing \in \left(P\textstyle\bigvee\right) Base^T(Td)(\Phi) = Base^T\left(T\textstyle\bigvee\right)(Td)(\Phi).$$

Thus, by part (1) of this lemma, we have

$$\nabla(T\textstyle\bigvee)(Td)(\Phi) = 0_{V_T\mathbb{L}}. \tag{6}$$

This completes the proof of part 2, since (4) follows immediately from (5) and (6).

(3) Suppose $\alpha, \beta \in TL$ are not linked by any lifted relation. Consider the (unique) map

$$f : L \to \{1\},$$

and define $\alpha' := (Tf)\alpha$, $\beta' := (Tf)(\beta)$. We now suppose, in order to show a contradiction, that $\alpha' = \beta'$. Then we would find $\alpha \, \overline{T}((Gr\, f)\breve{\ } ; Gr\, f) \, \beta$, contradicting the assumption on α and β. It follows that α' and β are distinct, so by part (2) of this lemma (with $A = \{1_{\mathbb{L}}\}$), we may infer that $\nabla\alpha' \wedge \nabla\beta' = 0_{V_T\mathbb{L}}$. This means that we are done, since it follows from $Gr\, f \subseteq \leqslant$ and the definitions of α', β' that $\alpha \, \overline{T}{\leqslant} \alpha'$ and $\beta \, \overline{T}{\leqslant} \beta'$, and from this we obtain by ($\nabla 1$) that

$$\nabla\alpha \wedge \nabla\beta \leqslant \nabla\alpha' \wedge \nabla\beta' \leqslant 0_{V_T\mathbb{L}}.$$

(4) We reason as follows:

$$1_{V_T \mathbb{L}} = \bigvee \{ \nabla(T\textstyle\bigwedge)(\Phi) \mid \Phi \in SRD(\varnothing) \} \qquad \text{(axiom (∇2) with } A = \varnothing \text{)}$$

$$= \bigvee \{ \nabla(T\textstyle\bigwedge)(\Phi) \mid \Phi \in T\{\varnothing\} \} \qquad \text{(Fact 3.8)}$$

$$= \bigvee \{ \nabla \gamma \mid \gamma \in T\{1_{\mathbb{L}}\} \} \qquad \text{($\ddagger$– see below)}$$

where the last step ($\ddagger$) is justified by the observation that since the map $\bigwedge : P_\omega L \to L$ restricts to a bijection $\bigwedge : \{\varnothing\} \to \{1_{\mathbb{L}}\}$, its lifting restricts to a bijection $T\bigwedge : T\{\varnothing\} \to T\{1_{\mathbb{L}}\}$.

(5) Let $A \subseteq L$ be such that $1_{\mathbb{L}} = \bigvee A$, and consider an arbitrary element $\Phi \in T\{A\}$. We claim that

$$\lambda^T(\Phi) \subseteq TA. \qquad (7)$$

To see this, take an arbitrary lifted element α of Φ. It follows from $\alpha \; \overline{T}{\in}\; \Phi$ that $Base^T(\alpha) \; \overline{P}{\in}\; Base^T(\Phi)$. In particular, each $a \in Base^T(\alpha)$ must belong to some $B \in Base^T(\Phi) \subseteq \{A\}$. In other words, $Base^T(\alpha) \subseteq A$, which is equivalent to saying that $\alpha \in TA$. This proves (7).

By (7) and axiom (∇3) we obtain

$$\nabla(T\textstyle\bigvee)(\Phi) \leqslant \bigvee \{ \nabla\alpha \mid \alpha \in TA \}. \qquad (8)$$

Now we reason as follows:

$$1_{V_T \mathbb{L}} = \bigvee \{ \nabla\alpha \mid \alpha \in T\{1_{\mathbb{L}}\} \} \qquad \text{(part 4)}$$

$$= \bigvee \{ \nabla(T\textstyle\bigvee)(\Phi) \mid \Phi \in T\{A\} \} \qquad \text{($*$ – see below)}$$

$$\leqslant \bigvee \{ \nabla\alpha \mid \alpha \in TA \} \qquad \text{(8)}.$$

To justify the step ($*$), we just note that if we restrict the map $\bigvee : PL \to L$ to the bijection $\bigvee : \{A\} \to \{1_{\mathbb{L}}\}$, as its lifting we obtain a bijection $T\bigvee : T\{A\} \to T\{1_{\mathbb{L}}\}$. $\square$

3.3. *Two examples of the T-powerlocale construction*

In this subsection we will discuss two examples of T-powerlocales. First, we discuss the somewhat trivial example of the Id-powerlocale, and then the defining example of T-powerlocales, namely the P_ω-powerlocale, which is isomorphic to the classical Vietoris powerlocale.

Example 3.11. Let $Id : \mathsf{Set} \to \mathsf{Set}$ be the identity functor on the category of sets. Then for all frames $\mathbb{L}$, we have $V_{Id}\mathbb{L} \simeq \mathbb{L}$.

First recall from Example 2.4 that for any relation $R \subseteq X \times Y$, we have $\overline{Id}R = R$. Moreover, if $A \in Id\, P_\omega L = P_\omega L$, it is straightforward to verify that

$$SRD(A) = \{ \Psi \in P_\omega(\textstyle\bigcup_{c \in A} \{c\}) \mid \forall c \in A, c \in \Psi \}$$

$$= \{A\}.$$

Consequently, the ∇-relations reduce to the following when $T = \mathrm{Id}$:

(∇1) $\qquad\qquad\qquad \nabla a \leqslant \nabla b, \qquad\qquad\qquad\qquad (a \leqslant b)$

(∇2) $\qquad\qquad\qquad \bigwedge_{a \in A} \nabla a \leqslant \nabla \bigwedge A, \qquad\qquad\qquad (A \in P_\omega L)$

(∇3) $\qquad\qquad\qquad \nabla \bigvee A \leqslant \bigvee \{\nabla b \mid b \in A\} \qquad\qquad (A \in P L).$

The identity $id_L : L \to L$ obviously satisfies (∇1), (∇2) and (∇3). Moreover, if we have a frame $\mathbb{M}$ and a function $f : L \to M$ that is compatible with (∇1), (∇2) and (∇3), it is easy to see that f is in fact a frame homomorphism $\mathbb{L} \to \mathbb{M}$. By the universal property of frame presentations, it follows that $V_{\mathrm{Id}} \mathbb{L} \simeq \mathbb{L}$.

We now turn to the P_ω-powerlocale. Recall from Example 2.2 that $P_\omega : \mathsf{Set} \to \mathsf{Set}$, the covariant finite power set functor, is indeed standard, weak pullback-preserving and finitary. We will now show that the P_ω-powerlocale is the Vietoris powerlocale. The equivalence of the ∇ axioms and the $\square, \diamond$ axioms on distributive lattices is already known from Palmigiano and Venema (2007); what is different here is that we consider infinite joins rather than just finite joins.

We will use the presentation using (∇1), (∇2$'$) and (∇3) as our point of departure. Recall that for all $\alpha, \beta \in P_\omega L$,

$$\alpha \leqslant_L \beta \text{ if } \alpha \subseteq \downarrow\beta$$
$$\alpha \leqslant_U \beta \text{ if } \uparrow\alpha \supseteq \beta$$
$$\alpha \leqslant_C \beta \text{ if } \alpha \leqslant_L \beta \text{ and } \alpha \leqslant_U \beta.$$

By Example 3.2, two of the relations presenting $V_{P_\omega} \mathbb{L}$ thus become

(∇2$'$) $\qquad\qquad \bigwedge_{\gamma \in \Gamma} \nabla\gamma \leqslant \bigvee \{\nabla\alpha \mid \forall \gamma \in \Gamma, \alpha \leqslant_C \gamma\}$

(∇3) $\qquad \nabla \left\{ \bigvee \alpha \mid \alpha \in \Phi \right\} \leqslant \bigvee \left\{ \nabla\beta \mid \beta \in P_\omega \left(\bigcup \Phi \right) \text{ and } \forall \alpha \in \Phi, \alpha \between \beta \right\}.$

Lemma 3.12. We consider the presentation of $V_{P_\omega} \mathbb{L}$.

(1) In the presence of (∇1), the relation (∇2$'$) can be replaced by

(∇2.0) $\qquad\qquad\qquad 1 \leqslant \bigvee \{\nabla\beta \mid \beta \in P_\omega L\}$

(∇2.2) $\qquad\qquad \nabla\gamma_1 \wedge \nabla\gamma_2 \leqslant \bigvee \{\nabla\beta \mid \beta \leqslant_C \gamma_1, \beta \leqslant_C \gamma_2\}.$

(2) In the presence of (∇1) and (∇2) (or its equivalent formulations), the relation (∇3) can be replaced by

(∇3.$\uparrow$) $\qquad \nabla \left(\gamma \cup \left\{ \bigvee^{\uparrow} S \right\} \right) \leqslant \bigvee^{\uparrow} \{\nabla(\gamma \cup \{a\}) \mid a \in S\} \quad (S \text{ directed})$

(∇3.0) $\qquad\qquad \nabla(\gamma \cup \{0\}) \leqslant 0$

(∇3.2) $\qquad\qquad \nabla(\gamma \cup \{a_1 \vee a_2\}) \leqslant \nabla(\gamma \cup \{a_1\}) \vee \nabla(\gamma \cup \{a_2\}) \vee \nabla(\gamma \cup \{a_1, a_2\}).$

Proof.

(1) ($\nabla 2.0$) and ($\nabla 2.2$) are the special cases of ($\nabla 2'$) when Γ is empty or a doubleton. To show that they imply ($\nabla 2'$) is an induction on the number of elements needed to enumerate the finite set Γ.

(2) Each of the replacement relations is a special case of ($\nabla 3$) in which all except one of the elements of Φ are singletons. We now show that they are sufficient to imply ($\nabla 3$). First, we will show for any *finite* S that

$$\nabla \left(\gamma \cup \left\{ \bigvee S \right\} \right) \leqslant \bigvee \left\{ \nabla (\gamma \cup \alpha) \mid \varnothing \neq \alpha \in P_\omega S \right\}.$$

We use induction on the length of a finite enumeration of S. The base case, S empty, is ($\nabla 3.0$). Now suppose $S = \{a\} \cup S'$. Then

$$\begin{aligned}
\nabla &\left(\gamma \cup \left\{ \bigvee S \right\} \right) \\
&= \nabla \left(\gamma \cup \left\{ a \vee \bigvee S' \right\} \right) \\
&\leqslant \nabla(\gamma \cup \{a\}) \vee \nabla \left(\gamma \cup \left\{ \bigvee S' \right\} \right) \vee \nabla \left((\gamma \cup \{a\}) \cup \left\{ \bigvee S' \right\} \right) \qquad \text{(by ($\nabla 3.2$))} \\
&\leqslant \nabla(\gamma \cup \{a\}) \vee \bigvee \left\{ \nabla(\gamma \cup \alpha') \mid \varnothing \neq \alpha' \in P_\omega S' \right\} \\
&\qquad \vee \bigvee \left\{ \nabla(\gamma \cup \{a\} \cup \alpha') \mid \varnothing \neq \alpha' \in P_\omega S' \right\} \qquad \text{(by induction)} \\
&= \bigvee \left\{ \nabla(\gamma \cup \alpha) \mid \varnothing \neq \alpha \in P_\omega S \right\}.
\end{aligned}$$

Now we can use ($\nabla 3.\uparrow$) to relax the finiteness condition on S, since for an arbitrary S we have

$$\begin{aligned}
\nabla \left(\gamma \cup \left\{ \bigvee S \right\} \right) &= \nabla \left(\gamma \cup \left\{ \bigvee^\uparrow \left\{ \bigvee S_0 \mid S_0 \in P_\omega S \right\} \right\} \right) \\
&\leqslant \bigvee^\uparrow \left\{ \nabla \left(\gamma \cup \left\{ \bigvee S_0 \right\} \right) \mid S_0 \in P_\omega S \right\}.
\end{aligned}$$

Finally, we can use induction on the length of a finite enumeration of Φ to deduce ($\nabla 3$). More precisely, we can show by induction on n that

$$\nabla \left(\gamma \cup \left\{ \bigvee S_1, \ldots, \bigvee S_n \right\} \right) \leqslant \bigvee \left\{ \nabla(\gamma \cup \alpha) \mid \varnothing \neq \alpha \in P_\omega \left(\bigcup_{i=1}^n S_i \right) \text{ and } \forall i, \, \alpha \between S_i \right\}.$$

This concludes the proof. $\qquad \square$

Remark 3.13. Relation ($\nabla 2.0$) can be weakened even further, to

$$1 \leqslant \nabla \varnothing \vee \nabla \{1\},$$

since if β is non-empty, $\beta \leqslant_C \{1\}$. We can also deduce from ($\nabla 2.2$) that $\nabla \varnothing \wedge \nabla \{1\} = 0$, showing that $\nabla \varnothing$ and $\nabla \{1\}$ are clopen complements.

Lemma 3.14. In $V\mathbb{L}$ we have, for any $S \subseteq \mathbb{L}$,

$$\Box \left(\bigvee S \right) = \bigvee \left\{ \Box \left(\bigvee \alpha \right) \wedge \bigwedge_{a \in \alpha} \Diamond a \mid \alpha \in P_\omega S \right\}.$$

Proof.

— ($\geqslant$):

This direction is immediate.

— ($\leqslant$):

First note that since $\bigvee S$ is a directed join $\bigvee^{\uparrow}_{\alpha \in P_\omega S} \bigvee \alpha$, we have

$$\square \left(\bigvee S \right) \leqslant \bigvee^{\uparrow}_{\alpha \in P_\omega S} \square \left(\bigvee \alpha \right),$$

so we just need to consider the case where S is finite. We will show that for every $\alpha, \beta \in P_\omega S$ we have

$$\square \left(\bigvee \alpha \vee \bigvee \beta \right) \wedge \bigwedge_{a \in \alpha} \diamond a \leqslant \text{RHS in statement},$$

after which the result follows by taking $\beta = S$ and $\alpha = \varnothing$. We use P_ω-induction on β, which is effectively an induction on the length of an enumeration of its elements. The base case, $\beta = \varnothing$, is trivial. For the induction step, suppose $\beta = \beta' \cup \{b\}$. Then

$$\square \left(\bigvee \alpha \vee \bigvee \beta \right) \wedge \bigwedge_{a \in \alpha} \diamond a$$

$$= \square \left(\bigvee \alpha \vee b \vee \bigvee \beta' \right) \wedge \bigwedge_{a \in \alpha} \diamond a$$

$$= \square \left(\bigvee \alpha \vee b \vee \bigvee \beta' \right) \wedge \bigwedge_{a \in \alpha} \diamond a \wedge \left(\square \left(\bigvee \alpha \vee \bigvee \beta' \right) \vee \diamond b \right)$$

$$= \left(\square \left(\bigvee \alpha \vee \bigvee \beta' \right) \wedge \bigwedge_{a \in \alpha} \diamond a \right) \vee \left(\square \left(\bigvee (\alpha \cup \{b\}) \vee \bigvee \beta' \right) \wedge \bigwedge_{a \in \alpha \cup \{b\}} \diamond a \right)$$

$$\leqslant \text{RHS, by induction.}$$

This concludes the proof. $\square$

Theorem 3.15. *Let* $\mathbb{L}$ *be a frame. Then* $V\mathbb{L} \cong V_{P_\omega}\mathbb{L}$.

Proof. First we define a frame homomorphism $\varphi : V_{P_\omega}\mathbb{L} \to V\mathbb{L}$ by

$$\varphi(\nabla\alpha) = \square \left(\bigvee \alpha \right) \wedge \bigwedge_{a \in \alpha} \diamond a.$$

We must check that this respects the relations:

— ($\nabla 1$):

Suppose $\alpha \leqslant_C \beta$. From $\alpha \leqslant_U \beta$ and $\alpha \leqslant_L \beta$ we get

$$\bigwedge_{a \in \alpha} \diamond a \leqslant \bigwedge_{b \in \beta} \diamond b$$
$$\bigvee \alpha \leqslant \bigvee \beta,$$

giving

$$\varphi(\nabla\alpha) \leqslant \varphi(\nabla\beta).$$

— (∇2.0):

We have

$$1 = \Box(0 \vee 1) = \Box 0 \vee (\Box 1 \wedge \Diamond 1) = \varphi(\nabla\varnothing) \vee \varphi(\nabla\{1\}).$$

— (∇2.2):

We have $\varphi(\nabla\gamma_1) \wedge \varphi(\nabla\gamma_2)$ is

$$
\begin{aligned}
\Box\left(\bigvee\gamma_1\right) \wedge \bigwedge_{c\in\gamma_1}\Diamond c \wedge \Box\left(\bigvee\gamma_2\right) &\wedge \bigwedge_{c'\in\gamma_2}\Diamond c' \\
&= \Box\left(\bigvee\gamma_1 \wedge \bigvee\gamma_2\right) \wedge \bigwedge_{c\in\gamma_1}\Diamond c \wedge \bigwedge_{c'\in\gamma_2}\Diamond c' \\
&= \Box\left(\bigvee\gamma_1 \wedge \bigvee\gamma_2\right) \wedge \bigwedge_{c\in\gamma_1}\Diamond\left(c \wedge \bigvee\gamma_1 \wedge \bigvee\gamma_2\right) \wedge \bigwedge_{c'\in\gamma_2}\Diamond\left(c' \wedge \bigvee\gamma_1 \wedge \bigvee\gamma_2\right) \\
&= \Box\left(\bigvee\gamma_1 \wedge \bigvee\gamma_2\right) \wedge \bigwedge_{c\in\gamma_1}\Diamond\left(c \wedge \bigvee\gamma_2\right) \wedge \bigwedge_{c'\in\gamma_2}\Diamond\left(c' \wedge \bigvee\gamma_1\right) \\
&= \Box\left(\bigvee_{c\in\gamma_1}\bigvee_{c'\in\gamma_2}c \wedge c'\right) \wedge \bigwedge_{c\in\gamma_1}\bigvee_{c'\in\gamma_2}\Diamond(c \wedge c') \wedge \bigwedge_{c'\in\gamma_2}\bigvee_{c\in\gamma_1}\Diamond(c \wedge c').
\end{aligned}
$$

Redistributing the disjunctions of the $\Diamond$s, we find that each resulting disjunct is of the form

$$\Box\left(\bigvee_{c\in\gamma_1}\bigvee_{c'\in\gamma_2}c \wedge c'\right) \wedge \bigwedge_{cRc'}\Diamond(c \wedge c')$$

for some $R \in P_\omega(\gamma_1 \times \gamma_2)$ such that $\gamma_1 \overline{P_\omega} R \gamma_2$. Note that for any such R, if we define $\beta_R = \{c \wedge c' \mid cRc'\}$, we have $\beta_R \leqslant_C \gamma_i$ ($i = 1, 2$). Now by Lemma 3.14, we have

$$
\begin{aligned}
\Box\left(\bigvee_{c\in\gamma_1}\bigvee_{c'\in\gamma_2}c \wedge c'\right) &\wedge \bigwedge_{cRc'}\Diamond(c \wedge c') \\
&\leqslant \bigvee\left\{\Box\left(\bigvee_{cR'c'}c \wedge c'\right) \wedge \bigwedge_{c(R\cup R')c'}\Diamond\left(c \wedge c'\right) \mid R' \in P_\omega(\gamma_1 \times \gamma_2)\right\} \\
&\leqslant \bigvee\left\{\Box\left(\bigvee_{cR'c'}c \wedge c'\right) \wedge \bigwedge_{cR'c'}\Diamond\left(c \wedge c'\right) \mid R \subseteq R' \in P_\omega(\gamma_1 \times \gamma_2)\right\} \\
&= \bigvee\left\{\varphi(\nabla\beta_{R'}) \mid R \subseteq R' \in P_\omega(\gamma_1 \times \gamma_2)\right\},
\end{aligned}
$$

and the result follows.

— (∇3.$\uparrow$):

The left-hand side is

$$
\begin{aligned}
\Box\left(\bigvee\gamma \vee \bigvee^{\uparrow}S\right) \wedge \bigwedge_{c\in\gamma}\Diamond c \wedge \Diamond\left(\bigvee^{\uparrow}S\right) \\
&= \bigvee^{\uparrow}\left\{\Box\left(\bigvee\gamma \vee a\right) \mid a \in S\right\} \wedge \bigvee^{\uparrow}\left\{\bigwedge_{c\in\gamma}\Diamond c \wedge \Diamond a \mid a \in S\right\} \\
&= \bigvee^{\uparrow}\left\{\Box\left(\bigvee\gamma \vee a\right) \wedge \bigwedge_{c\in\gamma}\Diamond c \wedge \Diamond a \mid a \in S\right\},
\end{aligned}
$$

which is the right-hand side.

— (∇3.0):

The left-hand side is

$$\Box\left(\bigvee\gamma \vee 0\right) \wedge \bigwedge_{c\in\gamma}\Diamond c \wedge \Diamond 0 = 0.$$

— (∇3.2):
The left-hand side is

$$\square \left(\bigvee \gamma \vee a_1 \vee a_2\right) \wedge \bigwedge_{c \in \gamma} \Diamond c \wedge \Diamond (a_1 \vee a_2)$$

$$= \bigvee_{i=1}^{2} \bigvee \left\{\square \left(\bigvee \beta\right) \wedge \bigwedge_{c \in \beta \cup \gamma \cup \{a_i\}} \Diamond c \mid \beta \in P_\omega \left(\gamma \cup \{a_1, a_2\}\right)\right\}$$

$$\leqslant \bigvee_{i=1}^{2} \bigvee \left\{\varphi \left(\nabla \left(\beta \cup \gamma \cup \{a_i\}\right)\right) \mid \beta \in P_\omega \left(\gamma \cup \{a_1, a_2\}\right)\right\}$$

$$= \bigvee_{i=1}^{2} \bigvee \left\{\varphi \left(\nabla \beta\right) \mid \gamma \cup \{a_i\} \subseteq \beta \in P_\omega \left(\gamma \cup \{a_1, a_2\}\right)\right\}$$

$$= \varphi \left(\nabla \left(\gamma \cup \{a_1\}\right)\right) \vee \varphi \left(\nabla \left(\gamma \cup \{a_2\}\right)\right) \vee \varphi \left(\nabla \left(\gamma \cup \{a_1, a_2\}\right)\right).$$

Next, we define the frame homomorphism $\psi : V\mathbb{L} \to V_{P_\omega}\mathbb{L}$ by

$$\psi \left(\square a\right) = \bigvee \left\{\nabla \alpha \mid \alpha \leqslant_L \{a\}\right\} = \nabla \varnothing \vee \nabla \{a\}$$

$$\psi \left(\Diamond a\right) = \bigvee \left\{\nabla \alpha \mid \alpha \leqslant_U \{a\}\right\} = \bigvee \left\{\nabla \left(\beta \cup \{a\}\right) \mid \beta \in P_\omega L\right\}.$$

(Observe that the expression for $\psi \left(\Diamond a\right)$ could be simplified even further to $\nabla\{1, a\}$.) We
check the relations. First, it is clear that ψ respects the monotonicity of $\square$ and $\Diamond$.

— $\square$ preserves directed joins:

$$\psi \left(\square \left(\bigvee_i^\uparrow a_i\right)\right) = \nabla \varnothing \vee \nabla\{\textstyle\bigvee_i^\uparrow a_i\} = \bigvee_i^\uparrow \psi \left(\square a_i\right).$$

— $\square$ preserves top:
This follows immediately from (∇2.0).

— $\square$ preserves binary meets:

$$\psi \left(\square a_1\right) \wedge \psi \left(\square a_2\right) = \nabla \varnothing \vee \left(\nabla\{a_1\} \wedge \nabla\{a_2\}\right)$$

$$= \nabla \varnothing \vee \bigvee \left\{\nabla \beta \mid \beta \leqslant_C \{a_1\}, \beta \leqslant_C \{a_2\}\right\}$$

$$= \nabla \varnothing \vee \nabla\{a_1 \wedge a_2\} = \psi \left(\square \left(a_1 \wedge a_2\right)\right).$$

— $\Diamond$ preserves joins:

$$\psi \left(\Diamond \left(\bigvee A\right)\right) = \bigvee \left\{\nabla \left(\beta \cup \{\textstyle\bigvee A\}\right) \mid \beta \in P_\omega L\right\}$$

$$= \bigvee \left\{\nabla \left(\beta \cup \alpha\right) \mid \beta \in P_\omega L, \varnothing \neq \alpha \in P_\omega A\right\}$$

$$= \bigvee_{a \in A} \bigvee \left\{\nabla \left(\beta \cup \{a\}\right) \mid \beta \in P_\omega L\right\} = \bigvee_{a \in A} \psi(\Diamond a).$$

— The first mixed relation:
Noting that $\nabla \varnothing \wedge \nabla \left(\beta \cup \{b\}\right) \leqslant \nabla \varnothing \wedge \nabla\{1\} = 0$, we have

$$\psi \left(\square a\right) \wedge \psi \left(\Diamond b\right) = \bigvee_{\beta \in P_\omega L} \left(\nabla \varnothing \vee \nabla\{a\}\right) \wedge \nabla \left(\beta \cup \{b\}\right)$$

$$= \bigvee_{\beta \in P_\omega L} \nabla\{a\} \wedge \nabla \left(\beta \cup \{b\}\right)$$

$$= \bigvee \left\{\nabla \gamma \mid \exists \beta, \gamma \leqslant_C \{a\}, \gamma \leqslant_C \beta \cup \{b\}\right\}$$

$$\leqslant \bigvee_{\beta \in P_\omega L} \nabla \left(\beta \cup \{a \wedge b\}\right) = \psi \left(\Diamond \left(a \wedge b\right)\right).$$

— The second mixed relation:

$$\psi\left(\square\left(a\vee b\right)\right)=\nabla\varnothing\vee\nabla\{a\vee b\}$$
$$=\nabla\varnothing\vee\nabla\{a\}\vee\nabla\{b\}\vee\nabla\{a,b\}$$
$$\leqslant\psi\left(\square a\right)\vee\psi\left(\diamond b\right)$$

since $\nabla\varnothing\vee\nabla\{a\}=\psi\left(\square a\right)$ and $\nabla\{b\}\vee\nabla\{a,b\}\leqslant\psi\left(\diamond b\right)$.

We still need to show that φ and ψ are mutually inverse. First, we have

$$\varphi\left(\psi\left(\square a\right)\right)=\varphi\left(\nabla\varnothing\vee\nabla\{a\}\right)=\square 0\vee\left(\square a\wedge\diamond a\right)=\square a$$

since $\square 0\wedge\diamond a\leqslant\diamond\left(0\wedge a\right)=0$.

Next, to show $\varphi\left(\psi\left(\diamond a\right)\right)=\diamond a$, we have

$$\varphi\left(\psi\left(\diamond a\right)\right)=\bigvee\nolimits_{\beta\in P_{\omega}L}\left(\square\left(\bigvee\beta\vee a\right)\wedge\bigwedge\nolimits_{b\in\beta}\diamond b\wedge\diamond a\right)$$
$$\leqslant\diamond a$$
$$=\square\left(1\vee a\right)\wedge\diamond 1\wedge\diamond a=\varphi\left(\nabla\{1,a\}\right)\leqslant\varphi\left(\psi\left(\diamond a\right)\right).$$

Finally, to show $\psi\left(\varphi\left(\nabla\alpha\right)\right)=\nabla\alpha$, we have

$$\psi\left(\varphi\left(\nabla\alpha\right)\right)=\psi\left(\square\left(\bigvee\alpha\right)\wedge\bigwedge\nolimits_{a\in\alpha}\diamond a\right)$$
$$=\left(\nabla\varnothing\vee\nabla\{\bigvee\alpha\}\right)\wedge\bigwedge\nolimits_{a\in\alpha}\bigvee\nolimits_{\beta_{a}\in P_{\omega}L}\nabla\left(\beta\cup\{a\}\right).$$

Now,

$$\bigwedge\nolimits_{a\in\alpha}\bigvee\nolimits_{\beta_{a}\in P_{\omega}L}\nabla\left(\beta\cup\{a\}\right)=\bigvee\left\{\nabla\gamma\mid\forall a\in\alpha,\exists\beta_{a}\in P_{\omega}L,\gamma\leqslant_{C}\beta_{a}\cup\{a\}\right\}$$
$$=\bigvee\left\{\nabla\gamma\mid\gamma\leqslant_{U}\alpha\right\}.$$

Also

$$\nabla\varnothing\wedge\bigvee\left\{\nabla\gamma\mid\gamma\leqslant_{U}\alpha\right\}=\bigvee\left\{\nabla\delta\mid\delta\leqslant_{C}\varnothing,\delta\leqslant_{U}\alpha\right\}$$
$$=\begin{cases}\nabla\alpha & \text{if }\alpha=\varnothing\\ 0 & \text{if }\alpha\neq\varnothing\end{cases}$$
$$\nabla\{\bigvee\alpha\}\wedge\bigvee\left\{\nabla\gamma\mid\gamma\leqslant_{U}\alpha\right\}=\bigvee\left\{\nabla\delta\mid\delta\leqslant_{C}\{\bigvee\alpha\},\delta\leqslant_{U}\alpha\right\}$$
$$=\nabla\left(\alpha\cup\{\bigvee\alpha\}\right)$$
$$=\bigvee\left\{\nabla\left(\alpha\cup\alpha'\right)\mid\varnothing\neq\alpha'\in P_{\omega}\alpha\right\}$$
$$=\begin{cases}0 & \text{if }\alpha=\varnothing\\ \nabla\alpha & \text{if }\alpha\neq\varnothing\end{cases}$$

It then follows that, independent of whether α is empty or not, $\psi\left(\varphi\left(\nabla\alpha\right)\right)=\nabla\alpha$. $\square$

3.4. Categorical properties of the T-powerlocale

In this section we discuss two categorical properties of the T-powerlocale construction. First we show how to extend the frame construction V_{T} to an endofunctor on the category Fr of frames. We will then show how the natural transformation $i\colon V_{P_{\omega}}\to V_{Id}$ (discussed

in Section 2.5 as $i \colon V \to Id$) can be generalised to a natural transformation

$$\widehat{\rho} \colon V_T \to V_{T'},$$

for any natural transformation $\rho \colon T' \to T$ satisfying some mild conditions (where T and T' are two finitary, weak pullback preserving set functors).

3.4.1. *V_T is a functor.* We start by introducing a natural way to transform a frame homomorphism $f \colon \mathbb{L} \to \mathbb{M}$ into a frame homomorphism from $V_T \mathbb{L}$ to $V_T \mathbb{M}$. To do this, we first prove the following technical lemma.

Lemma 3.16. Let $T \colon \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor, $\mathbb{L}, \mathbb{M}$ be frames and $f \colon \mathbb{L} \to \mathbb{M}$ be a frame homomorphism. Then the map $\nabla \circ Tf \colon TL \to V_T M$, that is, $\alpha \mapsto \nabla(Tf)(\alpha)$, is compatible with the relations $(\nabla 1)$, $(\nabla 2)$ and $(\nabla 3)$.

Proof. We use the abbreviation $\heartsuit := \nabla \circ Tf$, that is, for $\alpha \in TL$, we define $\heartsuit \alpha := \nabla(Tf)(\alpha)$.

In order to prove that $\heartsuit$ is compatible with $(\nabla 1)$, we need to show that

$$\text{for all } \alpha, \beta \in TL \colon \ \alpha \ \overline{T} {\leqslant}_{\mathbb{L}} \ \beta \text{ implies } \heartsuit \alpha \leqslant_{V_T M} \heartsuit \beta. \tag{9}$$

To see this, we assume that $\alpha, \beta \in TL$ are such that $\alpha \ \overline{T} {\leqslant}_{\mathbb{L}} \ \beta$. From this it follows, by Lemma 2.7 and the assumption that f is a frame homomorphism, that $(Tf)(\alpha) \ \overline{T} {\leqslant}_{\mathbb{M}} (Tf)(\beta)$. Then by $(\nabla 1)_{\mathbb{M}}$, we get that $\heartsuit \alpha \leqslant_{V_T M} \heartsuit \beta$, as required.

Proving compatibility with $(\nabla 2)$ boils down to showing

$$\text{for all } \Gamma \in P_\omega TL \colon \ \bigwedge_{\alpha \in \Gamma} \heartsuit \alpha \leqslant \bigvee \{\heartsuit(T\textstyle\bigwedge)(\Psi) \mid \Psi \in SRD(\Gamma)\}. \tag{10}$$

To do this, given $\Gamma \in P_\omega TL$, we use $\Gamma' \in P_\omega TM$ to denote the set

$$\Gamma' := (P_\omega Tf)(\Gamma) = \{(Tf)(\alpha) \mid \alpha \in \Gamma\}.$$

Then we may observe

$$\bigwedge_{\alpha \in \Gamma} \heartsuit \alpha = \bigvee \{\nabla(T\textstyle\bigwedge)(\Psi) \mid \Psi \in SRD(\Gamma')\} \tag{$\nabla 1$}$$

$$\leqslant \bigvee \{\nabla(T\textstyle\bigwedge)(TP_\omega f)(\Phi) \mid \Phi \in SRD(\Gamma)\} \tag{Lemma 3.9}$$

$$= \bigvee \{\nabla(Tf)(T\textstyle\bigwedge)(\Phi) \mid \Phi \in SRD(\Gamma)\} \tag{$\dagger$– see below}$$

$$= \bigvee \{\heartsuit(T\textstyle\bigwedge)(\Phi) \mid \Phi \in SRD(\Gamma)\} \tag{definition of $\heartsuit$}.$$

Here the identity marked $(\dagger)$ is easily justified from the fact that f is a homomorphism: it follows from $f \circ \bigwedge = \bigwedge \circ (P_\omega f)$ and the functoriality of T that

$$(Tf) \circ (T\textstyle\bigwedge) = (T\textstyle\bigwedge) \circ (TP_\omega f).$$

Finally, for compatibility with $(\nabla 3)$, we need to verify that

$$\text{for all } \Phi \in TPL \colon \ \heartsuit(T\textstyle\bigvee)(\Phi) \leqslant \bigvee \{\heartsuit \beta \mid \beta \ \overline{T} \in \Phi\}. \tag{11}$$

To prove this, we calculate for a given $\Phi \in TPL$:

$$
\begin{aligned}
\heartsuit(T\bigvee)(\Phi) &= \nabla(Tf)(T\bigvee)(\Phi) && \text{(definition of } \heartsuit) \\
&= \nabla(T\bigvee)(TPf)(\Phi) && \text{(} f \text{ a frame homomorphism)} \\
&\leqslant \bigvee\{\nabla\beta \mid \beta\,\overline{T}\in(TPf)(\Phi)\} && ((\nabla 3)_\mathbb{M}) \\
&= \bigvee\{\nabla(Tf)(\gamma) \mid \gamma\,\overline{T}\in\Phi\} && (\ddagger\text{-- see below}) \\
&= \bigvee\{\heartsuit\gamma \mid \gamma\,\overline{T}\in\Phi\} && \text{(definition of } \heartsuit).
\end{aligned}
$$

Here the identity ($\ddagger$) follows from the observation that for all $\beta \in TM$ and $\Phi \in TPL$, we have $\beta\,\overline{T}\in(TPf)(\Phi)$ if and only if β is of the form $\beta = (Tf)(\gamma)$ for some $\gamma \in TL$. Using Fact 2.6, this is easily derived from the observation that for $b \in M$ and $A \in PL$, we have $b \in (Pf)A$ if and only if $b = f(c)$ for some $c \in A$. □

Lemma 3.16 justifies the following definition.

Definition 3.17. Let $T : \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor and $f : \mathbb{L} \to \mathbb{M}$ be a frame homomorphism. We define $V_T f : V_T\mathbb{L} \to V_T\mathbb{M}$ to be the unique frame homomorphism extending

$$
\nabla \circ Tf : TL \to V_T M.
$$

Theorem 3.18. Let T be a standard, finitary, weak pullback-preserving functor. Then the operation V_T defined above is an endofunctor on the category Fr.

Proof. Since for an arbitrary $f : \mathbb{L} \to \mathbb{M}$ we have ensured by definition that $V_T F$ is a frame homomorphism from $V_T\mathbb{L}$ to $V_T\mathbb{M}$, it is just left to show that V_T maps the identity map of a frame to the identity map of its T-powerlocale, and distributes over function composition. We will confine our attention to the second property.

Let $f : \mathbb{K} \to \mathbb{L}$ and $g : \mathbb{L} \to \mathbb{M}$ be two frame homomorphisms. In order to show that $V_T(g \circ f) = V_T g \circ V_T f$, we first recall that $V_T(g \circ f)$ is by definition the unique frame homomorphism extending the map $\nabla_\mathbb{M} \circ T(g \circ f) : TK \to V_T\mathbb{M}$. Hence, it suffices to prove that the map $V_T g \circ V_T f$, which is obviously a frame homomorphism, extends $\nabla_\mathbb{M} \circ T(g \circ f)$. But it is easy to see why this is the case, since, given an arbitrary element $\alpha \in TK$, a straightforward unravelling of definitions shows that

$$
(V_T g \circ V_T f)(\alpha) = V_T g(\nabla_\mathbb{L}(Tf)(\alpha)) = \nabla_\mathbb{M}(Tg)(Tf)(\alpha) = \nabla_\mathbb{M} T(g \circ f)(\alpha),
$$

as required. □

3.4.2. *Natural transformations between V_T and $V_{T'}$.* Now that we have seen how each finitary, weak pullback preserving set functor T induces a functor V_T on the category of frames, we can investigate the relation *between* two such functors $V_T, V_{T'}$. In fact, we have already seen an example of this: recall that in Section 2.5 we mentioned Johnstone's result (Johnstone 1985) that the standard Vietoris functor V is in fact a *comonad* on the category of frames. In our nabla-based presentation of this functor as $V = V_{P_\omega}$, thinking of the identity functor on the category Fr as the Vietoris functor V_{Id}, we can view the

counit of this comonad as a natural transformation

$$i \colon V_{P_\omega} \to V_{Id},$$

given by $i_{\mathbb{L}} \colon \nabla A \mapsto \bigwedge A$. More precisely, we can show that the map $\heartsuit \colon P_\omega L \to L$ given by $\heartsuit A := \bigwedge A$ is compatible with the ∇-axioms, and hence can be uniquely extended to the homomorphism $i_{\mathbb{L}}$; subsequently, we can show that this i is natural in $\mathbb{L}$. Recall that in the case of a concrete topological space (X, τ), this counit corresponds on the dual side to the singleton map $\sigma_X \colon s \mapsto \{s\}$, which provides an embedding of a compact Hausdorff topology into its Vietoris space.

We will now see how to generalise this picture of the natural transformation $i \colon V_{P_\omega} \to V_{Id}$ being induced by the singleton natural transformation $\sigma \colon Id \to P_\omega$ to a more general setting. First consider the following definition.

Definition 3.19. Let T and T' be standard, finitary, weak pullback-preserving functors. A natural transformation $\rho \colon T' \to T$ is said to *respect relation lifting* if for any relation $R \subseteq X \times Y$ we have, for all $\alpha' \in T'X$ and $\beta' \in T'Y$,

$$\text{if } \alpha' \; \overline{T'}R \; \beta' \text{ then } \rho_X(\alpha') \; \overline{T}R \; \rho_Y(\beta'). \tag{12}$$

We say ρ is *base-invariant* if it commutes with *Base*, that is,

$$Base^{T'} = Base^T \circ \rho \tag{13}$$

for any set X.

Example 3.20. We can give three simple examples of base-invariant natural transformations that respect relation lifting:

(1) The base transformation $Base^T \colon T \to P_\omega$.
(2) The singleton natural transformation $\sigma \colon \mathrm{Id} \to P_\omega$, which is in fact a special case of (1).
(3) The diagonal map δ (given by $\delta_X \colon x \mapsto (x, x)$); it is straightforward to check that as a natural transformation, $\delta \colon \mathrm{Id} \to \mathrm{Id} \times \mathrm{Id}$ also satisfies both properties of Definition 3.19.

As we will see next, every base-invariant natural transformation $\rho \colon T' \to T$ that respects relation lifting induces a natural transformation $\widehat{\rho} \colon V_T \to V_{T'}$. In particular, the natural transformation $i \colon V \to \mathrm{Id}$ can be seen as $i = \widehat{\sigma}$, where $\sigma \colon \mathrm{Id} \to P_\omega$ is the singleton transformation discussed above.

Theorem 3.21. Let T and T' be standard, finitary, weak pullback-preserving functors. We assume that $\rho \colon T' \to T$ is a base-invariant natural transformation that respects relation lifting, and let $\mathbb{L}$ be a frame. Then the map from TL to $V_{T'}L$ given by

$$\alpha \mapsto \bigvee \{\nabla \alpha' \mid \alpha' \in T'L, \rho(\alpha') \; \overline{T} \leqslant \alpha\}$$

specifies a frame homomorphism

$$\widehat{\rho}_{\mathbb{L}} \colon V_T \mathbb{L} \to V_{T'} \mathbb{L}$$

that is natural in $\mathbb{L}$.

Proof. We use $\heartsuit : TL \to L$ to denote the map given in the statement of the Theorem, that is,

$$\heartsuit\alpha := \bigvee\{\nabla\alpha' \mid \alpha' \in T'L, \rho(\alpha') \, \overline{T}{\leqslant}\, \alpha\}.$$

We will first prove that this map is compatible with (∇1), (∇2) and (∇3), and then show the naturality of the induced frame homomorphism.

Claim 1. The map $\heartsuit$ is compatible with (∇1).

Proof of claim. To show that $\heartsuit$ is compatible with (∇1), we take two elements $\alpha, \beta \in TL$ such that $\alpha \, \overline{T}{\leqslant}\, \beta$. Then for any $\alpha' \in T'L$ such that $\rho(\alpha') \, \overline{T}{\leqslant}\, \alpha$, by the transitivity of $\overline{T}{\leqslant}$ (Fact 2.6(5)), we get $\rho(\alpha') \, \overline{T}{\leqslant}\, \beta$. From this it is immediate that $\heartsuit\alpha \leqslant \heartsuit\beta$, as required. $\qquad\square$

Claim 2. The map $\heartsuit$ is compatible with (∇2).

Proof of claim. For compatibility with (∇2), it suffices to show compatibility with (∇2′). That is, for $\Gamma \in P_\omega TL$, we will verify that

$$\bigwedge\{\heartsuit\gamma \mid \gamma \in \Gamma\} \leqslant \bigvee\{\heartsuit\beta \mid \beta \, \overline{T}{\leqslant}\, \gamma, \text{ for all } \gamma \in \Gamma\}. \tag{14}$$

We start by rewriting the left-hand side of (14) into

$$\bigwedge\{\heartsuit\gamma \mid \gamma \in \Gamma\} = \bigwedge\left\{\bigvee\left\{\nabla\gamma' \mid \rho(\gamma') \, \overline{T}{\leqslant}\, \gamma\right\} \mid \gamma \in \Gamma\right\} \qquad \text{(definition of } \heartsuit\text{)}$$

$$= \bigvee\left\{\bigwedge\left\{\varphi_\gamma \mid \gamma \in \Gamma\right\} \mid \varphi \in \mathscr{C}_\Gamma\right\} \qquad \text{(frame distributivity)}$$

where we define

$$\mathscr{C}_\Gamma := \{\varphi : \Gamma \to T'L \mid \rho(\varphi_\gamma) \, \overline{T}{\leqslant}\, \gamma, \text{ for all } \gamma \in \Gamma\}.$$

For any map $\varphi \in \mathscr{C}_\Gamma$, we may calculate

$$\bigwedge\{\varphi_\gamma \mid \gamma \in \Gamma\}$$

$$= \bigvee\{\nabla\gamma' \mid \gamma' \, \overline{T'}{\leqslant}\, \varphi_\gamma, \forall\gamma \in \Gamma\} \qquad (\nabla 2')$$

$$\leqslant \bigvee\{\nabla\gamma' \mid \rho(\gamma') \, \overline{T}{\leqslant}\, \rho(\varphi_\gamma), \forall\gamma \in \Gamma\} \qquad (\rho \text{ respects relation lifting})$$

$$\leqslant \bigvee\{\nabla\gamma' \mid \rho(\gamma') \, \overline{T}{\leqslant}\, \gamma, \forall\gamma \in \Gamma\} \qquad (\varphi \in \mathscr{C}_\Gamma, \text{ transitivity of } \overline{T}{\leqslant})$$

$$= \bigvee\left\{\bigvee\{\nabla\gamma' \mid \rho(\gamma') \, \overline{T}{\leqslant}\, \beta\} \mid \beta \, \overline{T}{\leqslant}\, \gamma, \forall\gamma \in \Gamma\right\} \qquad (\text{associativity of } \bigvee)$$

$$= \bigvee\{\heartsuit\beta \mid \beta \, \overline{T}{\leqslant}\, \gamma, \forall\gamma \in \Gamma\} \qquad (\text{definition of } \heartsuit).$$

(14) then follows immediately from the above calculations. $\qquad\square$

Claim 3. The map $\heartsuit$ is compatible with (∇3).

Proof of claim. We need to show, for an arbitrary but fixed set $\Phi \in TPL$, that

$$\heartsuit(T\textstyle\bigvee)(\Phi) = \bigvee\{\heartsuit\alpha \mid \alpha \, \overline{T}{\in}\, \Phi\}. \tag{15}$$

By definition, on the left-hand side of (15) we have

$$\heartsuit(T\textstyle\bigvee)(\Phi) = \bigvee\{\nabla\beta' \mid \rho(\beta') \, \overline{T}{\leqslant}\, (T\textstyle\bigvee)(\Phi)\},$$

while on the right-hand side we obtain, by the definition of $\heartsuit$,

$$\bigvee\{\heartsuit\alpha \mid \alpha\ \overline{T}\in\Phi\} = \bigvee\{\bigvee\{\nabla\alpha' \mid \rho(\alpha')\ \overline{T}\leqslant\alpha\} \mid \alpha\ \overline{T}\in\Phi\}$$
$$= \bigvee\{\nabla\alpha' \mid \rho(\alpha')\ \overline{T}(\leqslant\ ;\in)\ \Phi\}$$

where the second equality is by the associativity of $\bigvee$ and the compositionality of relation lifting (Fact 2.6(5)).

As a consequence, in order to establish the compatibility of $\heartsuit$ with (∇3), it suffices to show that

$$\nabla\beta' \leqslant \bigvee\{\nabla\alpha' \mid \rho(\alpha')\ \overline{T}(\leqslant\ ;\in)\ \Phi\}, \text{ for any } \beta' \text{ with } \rho(\beta')\ \overline{T}\leqslant (T\bigvee)(\Phi). \qquad (16)$$

Let β' be an arbitrary element of TL such that $\rho(\beta')\ \overline{T}\leqslant (T\bigvee)(\Phi)$. Our goal will be to find a set $\Phi' \in T'PL$ satisfying (20), (21) and (22) below, which is clearly enough to prove (16).

By Fact 2.8, we have

$$Base^T(\rho\beta')\ \overline{P}\leqslant Base^T((T\bigvee)(\Phi)) = (P\bigvee)\ Base^T(\Phi),$$

and since ρ is base-invariant, we have $Base^{T'}(\beta') = Base^T(\rho\beta')$. Combining these facts, we can see that $Base^{T'}(\beta')\ \overline{P}\leqslant (P\bigvee)\ Base^T(\Phi)$. This motivates the definition of the following map $\mathscr{H} : Base^{T'}(\beta') \to P_\omega PL$:

$$\mathscr{H}(b) := \{B \in Base^T(\Phi) \mid b \leqslant \bigvee B\}.$$

From the definitions, it is immediate that

$$\text{for all } b \in Base^{T'}(\beta') : \ b \leqslant \bigwedge\{\bigvee B \mid B \in \mathscr{H}(b)\}. \qquad (17)$$

Also, given a set $\mathscr{B} \in P_\omega PL$, let $\mathscr{C}_B$ be the collection of choice functions on $\mathscr{B}$, that is,

$$\mathscr{C}_{\mathscr{B}} := \{f : \mathscr{B} \to L \mid f(B) \in B \text{ for all } B \in \mathscr{B}\}.$$

Then it follows by frame distributivity that

$$\bigwedge\{\bigvee B \mid B \in \mathscr{B}\} = \bigvee\{\bigwedge(Pf)(\mathscr{B}) \mid f \in \mathscr{C}_{\mathscr{B}}\}. \qquad (18)$$

We define the map $K : P_\omega PL \to PL$ by

$$K(\mathscr{B}) := \{\bigwedge(Pf)(\mathscr{B}) \mid f \in \mathscr{C}_{\mathscr{B}}\}.$$

It then follows from (17) and (18) and the definitions that

$$\text{for all } b \in Base^{T'}(\beta') : \ b \leqslant \bigvee K(\mathscr{H}(b)). \qquad (19)$$

As a corollary, if we define

$$\Phi' := (T'K)(T'\mathscr{H})(\beta'),$$

it follows from (19), by the properties of relation lifting, that $\beta'\ \overline{T'}\leqslant (T'\bigvee)(\Phi')$, so an application of ($\nabla$1) yields

$$\nabla\beta' \leqslant \nabla(T'\bigvee)(\Phi'). \qquad (20)$$

Also, on the basis of an application of (∇3), we may conclude that

$$\nabla(T'\bigvee)(\Phi') \leqslant \bigvee\{\nabla\gamma' \mid \gamma'\ \overline{T'}\in\Phi'\}. \qquad (21)$$

This means that we will have proved (16) if we can show that

$$\text{for any } \gamma' \in T'L, \text{ if } \gamma' \ \overline{T} \in \Phi' \text{ then } \rho(\gamma') \ \overline{T}(\leqslant ; \in) \ \Phi. \tag{22}$$

For a proof of (22), let γ' be an arbitrary T'-lifted member of Φ' and recall that $\Phi' = (TK)(T\mathscr{H})(\beta')$. Then it follows by the assumption that ρ respects relation lifting that

$$\rho(\gamma') \ \overline{T} \in \rho(\Phi') = (TK)(T\mathscr{H})(\rho(\beta')).$$

Given our assumption on β', this means that the relation between $\rho(\gamma')$ and Φ can be summarised as

$$\rho(\gamma') \ \overline{T} \in (TK)(T\mathscr{H})(\beta) \text{ and } \beta \ \overline{T} \leqslant (T\bigvee)(\Phi) \text{ for some } \beta \in T \ Base^{T'}(\beta'), \tag{23}$$

where we may take $\rho(\beta')$ for β.

Returning to the ground level, observe that for any $c \in L$, $A \in Base^T(\Phi)$, we have

$$\text{if } c \in K\mathscr{H}(b) \text{ and } b \leqslant \bigvee A, \text{ for some } b \in Base^{T'}(\beta'), \text{ then } c \ (\leqslant ; \in) \ A. \tag{24}$$

To see why this is the case, assume that $c \in K\mathscr{H}(b)$ and $b \leqslant \bigvee A$, for some $b \in Base^{T'}(\beta')$. Then by the definition of $\mathscr{H}$, we find $A \in \mathscr{H}(b)$, while $c \in K\mathscr{H}(b)$ simply means that $c = \bigwedge\{f(B) \mid B \in \mathscr{H}(b)\}$, for some $f \in \mathscr{C}_{\mathscr{H}(b)}$. But then it is immediate that $c \leqslant f(A)$, while $f(A) \in A$ by the definition of $\mathscr{C}_{\mathscr{H}(b)}$. Thus $f(A)$ is the required element witnessing the fact that $c \ (\leqslant ; \in) \ A$.

But by the properties of relation lifting, we may derive from (24) that

$$\begin{aligned} &\text{if } \gamma \ \overline{T} \in (TK)(T\mathscr{H})(\beta) \text{ and } \beta \ \overline{T} \leqslant (T\bigvee)(\Phi) \text{ for some } \beta \in T \ Base^{T'}(\beta'), \\ &\text{then } \gamma \ \overline{T}(\leqslant ; \in) \ \Phi, \end{aligned} \tag{25}$$

so it is immediate by (23) that $\rho(\gamma') \ \overline{T}(\leqslant ; \in) \ \Phi$. This proves (22).

As we have already mentioned, the compatibility of $\heartsuit$ with ($\nabla 3$) is immediate by (20), (21) and (22), so this completes the proof of Claim 3. $\qquad\square$

As a result of Claims 1–3, we may uniquely extend $\heartsuit$ to a homomorphism $\widehat{\rho}_{\mathbb{L}} : V_T\mathbb{L} \to V_{T'}\mathbb{L}$. Hence, it is clear that to prove the theorem it now suffices to prove the following claim.

Claim 4. The family of homomorphisms $\widehat{\rho}_{\mathbb{L}}$ constitutes a natural transformation

$$\widehat{\rho} : V_T \to V_{T'}.$$

Proof of claim. Given two frames $\mathbb{L}$ and $\mathbb{M}$ and a frame homomorphism $f : \mathbb{L} \to \mathbb{M}$, we need to show that the following diagram commutes:

$$\begin{array}{ccc} V_T\mathbb{L} & \xrightarrow{\ \widehat{\rho}_{\mathbb{L}}\ } & V_{T'}\mathbb{L} \\ {\scriptstyle V_Tf}\downarrow & & \downarrow{\scriptstyle V_{T'}f} \\ V_T\mathbb{M} & \xrightarrow{\ \widehat{\rho}_{\mathbb{M}}\ } & V_{T'}\mathbb{M} \end{array}$$

To show this, we take an arbitrary element $\alpha \in TL$, and consider the following calculation:

$$(V_{T'}f)(\widehat{\rho}_{\mathbb{L}}(\nabla\alpha))$$

$$= (V_{T'}f)(\heartsuit\alpha) \qquad\qquad \text{(definition of } \widehat{\rho}_{\mathbb{L}})$$

$$= (V_{T'}f)\left(\bigvee\{\nabla\beta' \mid \rho_L(\beta')\,\overline{T}{\leqslant}\,\alpha\}\right) \qquad\qquad \text{(definition of } \heartsuit)$$

$$= \bigvee\{(V_T'f)(\nabla\beta') \mid \rho_L(\beta')\,\overline{T}{\leqslant}\,\alpha\} \qquad (V_{T'}f \text{ is a frame homomorphism)}$$

$$= \bigvee\{\nabla(T'f)(\beta') \mid \rho_L(\beta')\,\overline{T}{\leqslant}\,\alpha\} \qquad\qquad \text{(definition of } V_{T'}f)$$

$$= \bigvee\{\nabla\delta' \mid \rho_M(\delta')\,\overline{T}{\leqslant}\,(Tf)(\alpha)\} \qquad\qquad (\dagger\text{– see below)}$$

$$= \heartsuit(Tf)(\alpha) \qquad\qquad \text{(definition of } \heartsuit)$$

$$= \widehat{\rho}_{\mathbb{M}}(\nabla(Tf)(\alpha)) \qquad\qquad \text{(definition of } \widehat{\rho}_{\mathbb{M}})$$

$$= \widehat{\rho}_{\mathbb{M}}((V_Tf)(\nabla\alpha)) \qquad\qquad \text{(definition of } V_Tf).$$

Here the crucial step, marked (†), is proved by establishing the $\leqslant$ and $\geqslant$ inequalities as follows:

— ($\leqslant$):

It is straightforward to show that the set of joinands on the left-hand side is included in that on the right-hand side, and this follows from

$$\rho_L(\beta')\,\overline{T}{\leqslant}\,\alpha \text{ implies } \rho_M((T'f)(\beta'))\,\overline{T}{\leqslant}\,(Tf)(\alpha). \qquad (26)$$

To prove (26), suppose $\rho_L(\beta')\,\overline{T}{\leqslant}\,\alpha$. Then it follows by the fact that f is a homomorphism, and hence, monotone, that $(Tf)(\rho_L(\beta'))\,\overline{T}{\leqslant}\,(Tf)(\alpha)$. But since ρ is a natural transformation, we also have $(Tf)(\rho_L(\beta')) = \rho_M(T'f)(\beta')$, and from this (26) is immediate.

— ($\geqslant$):

We need to prove

$$\bigvee\{\nabla\delta' \mid \rho_M(\delta')\,\overline{T}{\leqslant}\,(Tf)(\alpha)\} \leqslant \bigvee\{\nabla(T'f)(\beta') \mid \rho_L(\beta')\,\overline{T}{\leqslant}\,\alpha\}. \qquad (27)$$

So we fix an arbitrary element $\delta' \in TL$ such that $\rho_M(\delta')\,\overline{T}{\leqslant}\,(Tf)(\alpha)$.
We define the map $h\colon Base^{T'}(\delta') \to L$ by putting

$$h(d) := \bigwedge\{a \in Base^T(\alpha) \mid d \leqslant f(a)\}.$$

Then, for all $d \in Base^{T'}(\delta')$ and all $a \in Base^T(\alpha)$, we find that $d \leqslant fa$ implies $hd \leqslant a$, which can be expressed by the relational inclusion

$$Gr\,f\,;\geqslant;\,Gr\,h \;\subseteq\; \geqslant,$$

so, by the properties of relation lifting, we may conclude that

$$Gr\,(Tf)\,;\overline{T}{\geqslant};\,Gr\,(Th) \;\subseteq\; \overline{T}{\geqslant},$$

which is just another way of saying that for all $\delta \in T\,Base^{T'}(\delta')$, we have

$$\delta\,\overline{T}{\leqslant}\,(Tf)(\alpha) \text{ only if } (Th)(\delta)\,\overline{T}{\leqslant}\,\alpha. \qquad (28)$$

We now define

$$\beta' := (T'h)(\delta'),$$

and can conclude from the fact that ρ respects relation lifting that

$$\rho_L(\beta') = (Th)\rho_M(\delta'),$$

and thus, by the assumption that

$$\rho_M(\delta') \, \overline{T}{\leqslant} (Tf)(\alpha),$$

we get by (28) that

$$\rho_L(\beta') \, \overline{T}{\leqslant} \alpha. \tag{29}$$

Similarly, from the fact that $d \leqslant fhd$ for each $d \in Base^{T'}(\delta')$, we can get $\delta' \, \overline{T'}{\leqslant} (T'f)(\beta')$, and thus, by ($\nabla$1), we may conclude that

$$\nabla\delta' \leqslant \nabla(T'f)(\beta'). \tag{30}$$

Finally, since (26) follows immediately from (25) and (30), this concludes the proof of Claim 4. $\qquad\square$

This completes the proof of the theorem. $\qquad\square$

Remark 3.22. The definition of the $\widehat{\rho}_{\mathbb{L}} : V_T\mathbb{L} \to V_{T'}\mathbb{L}$ using the assignment

$$\alpha \mapsto \bigvee \{\nabla\alpha' \mid \alpha' \in T'L, \rho(\alpha') \, \overline{T}{\leqslant} \alpha\}$$

is very similar to that of a right adjoint. If it were the case that $\widehat{\rho}_{\mathbb{L}}$ preserved *all* meets, then the adjoint functor theorem would allow us to define its left adjoint. However, we only have a proof that $\widehat{\rho}_{\mathbb{L}} : V_T\mathbb{L} \to V_{T'}\mathbb{L}$ preserves *finite* conjunctions, so it is not at all obvious at this point that there even is a left adjoint to $\widehat{\rho}_{\mathbb{L}}$. This is an interesting question for future work.

3.5. *T-powerlocales through flat sites*

In this section, we will show that $V_T\mathbb{L}$, the T-powerlocale of a given frame $\mathbb{L}$, has a flat-site presentation as $V_T\mathbb{L} \simeq \mathrm{Fr}\langle TL, \overline{T}{\leqslant}, \lhd_0^{\mathbb{L}} \rangle$. It then follows by the Flat-Site Coverage Theorem that every element of $V_T\mathbb{L}$ has a disjunctive normal form, and that the suplattice reduct of $V_T\mathbb{L}$ has a presentation defined only in terms of the order $\overline{T}{\leqslant}$ and the lifted join function $T\bigvee : TPL \to TL$.

Recall that $\langle X, \sqsubseteq, \lhd_0 \rangle$ is a flat site if $\langle X, \sqsubseteq \rangle$ is a pre-order and $\lhd_0$ is a basic cover relation compatible with $\sqsubseteq$. In that case, we know that $\langle X, \sqsubseteq, \lhd_0 \rangle$ presents a frame $\mathrm{Fr}\langle X, \sqsubseteq, \lhd_0 \rangle$, and that if we denote the insertion of generators by $\heartsuit : X \to \mathrm{Fr}\langle X, \sqsubseteq, \lhd_0 \rangle$, then

$$\begin{aligned}
\mathrm{Fr}\langle X, \sqsubseteq, \lhd_0 \rangle \simeq \mathrm{Fr}\langle X \mid &\heartsuit a \leqslant \heartsuit b \quad (a \sqsubseteq b), \\
&1 = \bigvee \{\heartsuit a \mid a \in X\} \\
&\heartsuit a \wedge \heartsuit b = \bigvee \{\heartsuit c \mid c \sqsubseteq a, c \sqsubseteq b\} \\
&\heartsuit a \leqslant \bigvee \{\heartsuit b \mid b \in A\} \quad (a \lhd_0 A)\rangle.
\end{aligned}$$

Note that this is very similar to our presentation of $V_T\mathbb{L}$ from Corollary 3.6 using ($\nabla 1$), ($\nabla 2'$) and ($\nabla 3$), namely,

$$V_T\mathbb{L} \simeq \mathrm{Fr}\langle TL \mid \nabla\alpha \leqslant \nabla\beta \quad (\alpha\,\overline{T}{\leqslant}\,\beta),$$
$$\textstyle\bigwedge_\Gamma \nabla\gamma = \bigvee\{\nabla\delta \mid \forall\gamma \in \Gamma,\ \delta\,\overline{T}{\leqslant}\,\gamma\} \quad (\Gamma \in TP_\omega L)$$
$$\textstyle\nabla T\bigvee(\Phi) \leqslant \bigvee\{\nabla\beta \mid \beta \in \lambda^T(\Phi)\} \quad (\Phi \in TPL)\rangle.$$

We will see below that if we define a cover relation $\vartriangleleft_0^{\mathbb{L}}$, which is inspired by ($\nabla 3$), we obtain a flat site $\langle TL, \overline{T}{\leqslant}, \vartriangleleft_0^{\mathbb{L}}\rangle$, and this flat site presents $V_T\mathbb{L}$.

So how do we go about defining a basic cover relation $\vartriangleleft_0^{\mathbb{L}} \subseteq TL \times PTL$ so we can give a presentation of $V_T\mathbb{L}$? Intuitively, we would like to take the T-lifting of the relation

$$\textstyle\{(a,A) \in L \times PL \mid a \leqslant \bigvee A\} = {\leqslant}\,;\left(Gr\,\bigvee\right)^{\smallsmile}.$$

However, the T-lifting of this relation is of type $TL \times TPL$, while a basic cover relation on $\langle TL, \overline{T}{\leqslant}\rangle$ should be of type $TL \times PTL$. We solve this by involving the natural transformation $\lambda^T : TP \to PT$, given by

$$\lambda^T(\Phi) := \{\beta \in TL \mid \beta\,\overline{T}{\in}\,\Phi\},$$

and assigning to each $\Phi \in TPL$ the set of its lifted members. That is, we define

$$\vartriangleleft_0^{\mathbb{L}} := \textstyle\{(\alpha, \lambda^T(\Phi)) \in L \times PTL \mid \alpha\,\overline{T}{\leqslant}\,T\bigvee(\Phi)\}.$$

In other words, we put $\alpha\,\vartriangleleft_0^{\mathbb{L}}\,\Gamma$ if and only if Γ is of the form $\lambda^T(\Phi)$ for some $\Phi \in TPL$ such that $\alpha\,\overline{T}{\leqslant}\,(T\bigvee)\Phi$. Two tasks lie ahead of us: first, we must show that $\langle TL, \overline{T}{\leqslant}, \vartriangleleft_0^{\mathbb{L}}\rangle$ is a flat site, meaning that $\vartriangleleft_0$ is compatible with $\overline{T}{\leqslant}$; and then we must show that $\langle TL, \overline{T}{\leqslant}, \vartriangleleft_0^{\mathbb{L}}\rangle$ presents $V_T\mathbb{L}$. The following technical observation about the relation $\alpha\,\overline{T}{\leqslant}\,T\bigvee(\Phi)$ is the main reason $V_T\mathbb{L}$ admits a flat-site presentation. The reason for introducing a $\wedge$-semilattice $\mathbb{M}$ below will become apparent in Section 4.3.

Lemma 3.23. Let $T : \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor, $\mathbb{L}$ be a frame and $\mathbb{M}$ be a $\wedge$-subsemilattice of $\mathbb{L}$. Then for all $\alpha \in TM$ and $\Phi \in TPM$ such that $\alpha\,\overline{T}{\leqslant}\,T\bigvee(\Phi)$, there exists $\Phi' \in TPM$ such that:

(1) $\alpha\,\overline{T}{\leqslant}\,T\bigvee(\Phi')$.
(2) $\Phi'\,\overline{T}{\subseteq}\,T{\downarrow}_L \circ T\eta(\alpha)$.
(3) $\Phi'\,\overline{T}{\subseteq}\,T{\downarrow}_L(\Phi)$.

Proof. We first define the following relation on $M \times PM$:

$$R := \textstyle\{(a,A) \in M \times PM \mid a \leqslant \bigvee A\} = \left({\leqslant}\,;(Gr\,\bigvee)^{\smallsmile}\right) \upharpoonright_{M \times PM}.$$

Consider the span $M \xleftarrow{p_1} R \xrightarrow{p_2} PM$. We define the following function $f : R \to R$:

$$f : (a,A) \mapsto (a, a \wedge A),$$

where $a \wedge A := \{a \wedge b \mid b \in A\}$. To see that this function is well defined, first observe that $a \wedge A \in PM$ because $\mathbb{M}$ is a $\wedge$-subsemilattice of $\mathbb{L}$. Moreover, by frame distributivity, we see that if $(a,A) \in R$, that is, if $a \leqslant \bigvee A$, then $a \leqslant \bigvee(a \wedge A)$ also, so $(a, a \wedge A) \in R$. Now

observe that $f : R \to R$ satisfies an equation and two inequations: for all $(a, A) \in R$,

$$p_1 \circ f(a, A) = a = p_1(a, A) \qquad \text{(by the definition of } f\text{)}$$

$$p_2 \circ f(a, A) = a \wedge A \sqsubseteq_L \downarrow_L\{a\} = \downarrow_L \circ \eta_L \circ p_1(a, A) \qquad \text{(since } \forall b \in A, a \wedge b \leqslant a\text{)}$$

$$p_2 \circ f(a, A) = a \wedge A \sqsubseteq_L \downarrow_L A = \downarrow_L \circ p_2(a, A) \qquad \text{(since } \forall b \in A, a \wedge b \leqslant b \in A\text{)}.$$

Now consider the lifted diagram

$$TM \xleftarrow{\ Tp_1\ } TR \xrightarrow{\ Tp_2\ } TPM.$$

It follows from Lemma 2.7 and the equation/inequations above that for each $\delta \in TR$, we have

$$Tp_1 \circ Tf(\delta) = Tp_1(\delta) \tag{31}$$

$$Tp_2 \circ Tf(\delta)\ \overline{T} \sqsubseteq_L T\downarrow_L \circ T\eta_L \circ Tp_1(\delta) \tag{32}$$

$$Tp_2 \circ Tf(\delta)\ \overline{T} \sqsubseteq_L T\downarrow_L \circ Tp_2(\delta). \tag{33}$$

Now recall that by Fact 2.6,

$$\overline{T}\leqslant\ ;\ Gr(T\bigvee)\check{\ } = \overline{T}(\leqslant\ ;\ (Gr\ \bigvee)\check{\ }) = \overline{T}R,$$

so we see that $\alpha\ \overline{T}\leqslant T\bigvee(\Phi)$ if and only if $\alpha\ \overline{T}R\ \Phi$. So we let $\alpha \in TM$ and $\Phi \in TPM$ such that $\alpha\ \overline{T}\leqslant T\bigvee(\Phi)$, that is, such that $\alpha\ \overline{T}R\ \Phi$. We will now show that there is a $\Phi' \in TPM$ satisfying properties (1)–(3). First, observe that by the definition of relation lifting, there must exist some $\delta \in TR$ such that

$$Tp_1(\delta) = \alpha$$

$$Tp_2(\delta) = \Phi.$$

We claim that $\Phi' := Tp_2 \circ Tf(\delta)$ satisfies properties (1)–(3). We know by the definition of relation lifting that $(Tp_1 \circ Tf(\delta))\ \overline{T}R\ (Tp_2 \circ Tf(\delta))$. But, since

$$Tp_1 \circ Tf(\delta) = Tp_1(\delta) \qquad \text{(by (31))}$$

$$= \alpha \qquad \text{(by assumption)},$$

it follows that $\alpha\ \overline{T}R\ \Phi'$, that is, $\alpha\ \overline{T}\leqslant T\bigvee(\Phi')$, and we can conclude that (1) holds. Moreover, it follows immediately from (32) that (2) holds. Similarly, it follows immediately from (33) that (3) holds. $\qquad\square$

In the above lemma, we have used the lifted inclusion relation $T\subseteq$ and the lifted downset function $T\downarrow$. In the lemma below we record some elementary observations about the interaction between $T\subseteq$, $T\downarrow$ and the natural transformation $\lambda^T : TP \to PT$.

Lemma 3.24. Let $T :$ Set $\to$ Set be a standard, finitary, weak pullback-preserving functor, $\langle X, \sqsubseteq \rangle$ be a pre-order, $\alpha \in TX$ and $\Phi, \Phi' \in TPX$. Then:

(1) $\downarrow_{TX}\lambda^T(\Phi) = \lambda^T(T\downarrow_X(\Phi))$.
(2) $\downarrow_{TX}\{\alpha\} = \lambda^T(T\downarrow_X \circ T\eta_X(\alpha))$.
(3) If $\Phi'\ \overline{T}\sqsubseteq_X \Phi$, then also $\lambda^T(\Phi') \subseteq \lambda^T(\Phi)$.

Proof.

(1) For all $a \in X$ and all $A \in PX$, we have $a \leqslant ; \in A$ if and only if $a \in \downarrow_X A$. Consequently,

$$\forall \alpha \in TL, \forall \Phi \in TPL, \alpha\, \overline{T}{\leqslant}\, ; \overline{T} \in \Phi \text{ iff } \alpha\, \overline{T} \in T{\downarrow_X}(\Phi).$$

Now we can see that

$$\begin{aligned}
\alpha \in \downarrow_{TX} \lambda^T(\Phi) &\Leftrightarrow \alpha\, \overline{T}{\leqslant}\, ; \overline{T} \in \Phi && \text{(by the definition of } \downarrow \text{ and } \lambda^T) \\
&\Leftrightarrow \alpha\, \overline{T} \in T{\downarrow_X}(\Phi) && \text{(by the above)} \\
&\Leftrightarrow \alpha \in \lambda^T \left(T{\downarrow_X}(\Phi)\right) && \text{(by the definition of } \lambda^T).
\end{aligned}$$

(2) For all $a, b \in X$, we have $b \leqslant a$ if and only if $b \in \downarrow_X \{a\}$. It follows by relation lifting that

$$\forall \alpha, \beta \in TX, \beta\, \overline{T}{\leqslant}\, \alpha \text{ iff } \beta\, \overline{T} \in T{\downarrow_X} \circ T\eta_X(\alpha).$$

It now follows by an argument analogous to that used for part (1) that (2) also holds.

(3) Observe that for all $A, A' \in PX$ and all $a \in X$, we have that $a \in A' \subseteq A$ implies that $a \in A$. The statement then follows by relation lifting. $\square$

We are now ready to prove that $\langle TL, \overline{T}{\leqslant}, \lhd_0^{\mathbb{L}} \rangle$ is indeed a flat site.

Lemma 3.25. Let $T : \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor. If $\mathbb{L}$ is a frame, then $\langle TL, \overline{T}{\leqslant}, \lhd_0^{\mathbb{L}} \rangle$ is a flat site.

Proof. We already know from Lemma 2.7 that $\langle TL, \overline{T}{\leqslant} \rangle$ is a pre-order, so what remains to be shown is that the relation $\lhd_0^{\mathbb{L}}$ is compatible with the pre-order. Fix $\alpha \in TL$ and $\Phi \in TPL$ such that $\alpha\, \overline{T}{\leqslant}\, T\bigvee(\Phi)$ so that $\alpha \lhd_0^{\mathbb{L}} \lambda^T(\Phi)$. We need to show that

$$\forall \beta \in TL, \text{ if } \beta\, \overline{T}{\leqslant}\, \alpha \text{ then } \exists \Gamma \in TPL \text{ with } \Gamma \subseteq \downarrow_{TL}\{\beta\} \cap \downarrow_{TL} \lambda^T(\Phi) \text{ and } \beta \lhd_0^{\mathbb{L}} \Gamma. \quad (34)$$

But this is easy to see, since if $\beta\, \overline{T}{\leqslant}\, \alpha$, then, since $\alpha\, \overline{T}{\leqslant}\, T\bigvee(\Phi)$, it follows by transitivity of $\overline{T}{\leqslant}$ that $\beta\, \overline{T}{\leqslant}\, T\bigvee(\Phi)$. Now by Lemma 3.23, there exists $\Phi' \in TPL$ such that

$$\begin{aligned}
\alpha\, \overline{T} &\leqslant T\bigvee(\Phi') \\
\Phi'\, \overline{T} &\subseteq T{\downarrow_L} \circ T\eta(\beta) \\
\Phi'\, \overline{T} &\subseteq T{\downarrow_L}\Phi.
\end{aligned}$$

If we define $\Gamma := \lambda^T(\Phi')$, then we have $\beta \lhd_0 \Gamma$ by the definition of $\lhd_0^{\mathbb{L}}$. Moreover, it now follows from Lemma 3.24 that $\Gamma \subseteq \downarrow_{TL}\{\beta\} \cap \downarrow_{TL} \lambda^T(\Phi)$. We can now conclude that (34) holds. Hence, since $\alpha \in TL$ and $\Phi \in TPL$ were arbitrary, we have shown that $\lhd_0^{\mathbb{L}}$ is compatible with the order $\overline{T}{\leqslant}$, so $\langle TL, \overline{T}{\leqslant}, \lhd_0^{\mathbb{L}} \rangle$ is a flat site. $\square$

Having established that $\langle TL, \overline{T}{\leqslant}, \lhd_0^{\mathbb{L}} \rangle$ is a flat site, we will now prove that it presents $V_T\mathbb{L}$, that is, that $V_T\mathbb{L} \simeq \mathrm{Fr}\langle TL, \overline{T}{\leqslant}, \lhd_0^{\mathbb{L}} \rangle$.

Theorem 3.26. Let $\mathbb{L}$ be a frame and T be a standard, finitary, weak pullback-preserving functor. Then $V_T\mathbb{L}$ admits the following flat-site presentation:

$$V_T\mathbb{L} \simeq \mathrm{Fr}\langle TL, \overline{T}{\leqslant}, \lhd_0^{\mathbb{L}} \rangle,$$

where

$$\lhd_0^{\mathbb{L}} = \{(\alpha, \lambda^T(\Phi)) \in L \times PTL \mid \alpha \, \overline{T} \leqslant T\bigvee(\Phi)\},$$

and in each direction the isomorphism is the unique frame homomorphism extending the identity map id_{TL} on the set of generators of $V_T\mathbb{L}$ and $\mathrm{Fr}\langle TL, \overline{T}\leqslant, \lhd_0^{\mathbb{L}}\rangle$, respectively.

Proof. We denote the insertion of generators from TL to $V_T\mathbb{L}$ by ∇, and from TL to $\mathrm{Fr}\langle TL, \overline{T}\leqslant, \lhd_0^{\mathbb{L}}\rangle$ by $\heartsuit$. We will show that:

(1) The function $\heartsuit \colon TL \to \mathrm{Fr}\langle TL, \overline{T}\leqslant, \lhd_0^{\mathbb{L}}\rangle$ is compatible with the relations (∇1), (∇2′) and (∇3).

(2) The function $\nabla \colon TL \to V_T\mathbb{L}$ has the following properties:

    (a) ∇ is order-preserving.

    (b) $1 = \bigvee\{\nabla\alpha \mid \alpha \in TL\}$.

    (c) For all $\alpha, \beta \in TL$, we have $\nabla\alpha \wedge \nabla\beta = \bigwedge\{\nabla\gamma \mid \delta \, \overline{T}\leqslant \alpha, \beta\}$.

    (d) For all $\alpha \lhd_0^{\mathbb{L}} \Gamma$, we have $\nabla\alpha \leqslant \bigvee\{\nabla\beta \mid \beta \in \Gamma\}$.

(1) We consider compatability with (∇1), (∇2′) and (∇3) in turn:

  — (∇1):

    Suppose $\alpha, \beta \in TL$ such that $\alpha \, \overline{T}\leqslant \beta$. We have to show that $\heartsuit\alpha \leqslant \heartsuit\beta$. This follows immediately from the fact that $\heartsuit \colon TL \to \mathrm{Fr}\langle TL, \overline{T}\leqslant, \lhd_0^{\mathbb{L}}\rangle$ is order-preserving.

  — (∇2′):

    Let $\Gamma \in P_\omega TL$. We have to show that

$$\bigwedge_{\gamma \in \Gamma}\heartsuit\gamma \leqslant \bigvee\{\heartsuit\delta \mid \forall\gamma \in \Gamma, \delta \, \overline{T}\leqslant \gamma\}. \tag{35}$$

    Recall from Section 2.4 that since $\langle TL, \overline{T}\leqslant, \lhd_0^{\mathbb{L}}\rangle$ is a flat site, we know that

$$1 = \bigvee\{\heartsuit\alpha \mid \alpha \in TL\},$$

    and that for all $\alpha, \beta \in TL$,

$$\heartsuit\alpha \wedge \heartsuit\beta = \bigwedge\{\heartsuit\gamma \mid \delta \, \overline{T}\leqslant \alpha, \beta\}.$$

    It now follows by induction on the size of Γ that (35) holds.

  — (∇3):

    We take $\Phi \in TPL$. We have to show that

$$\heartsuit T\bigvee(\Phi) \leqslant \bigvee\{\heartsuit\beta \mid \beta \in \lambda^T(\Phi)\}.$$

    But this follows immediately from the definition of $\lhd_0^{\mathbb{L}}$, since

$$T\bigvee(\Phi) \, \overline{T}\leqslant T\bigvee(\Phi).$$

So we can conclude that

$$\heartsuit \colon TL \to \mathrm{Fr}\langle TL, \overline{T}\leqslant, \lhd_0^{\mathbb{L}}\rangle$$

is compatible with the relations (∇1), (∇2) and (∇3) and thus there must be a unique frame homomorphism

$$f \colon V_T\mathbb{L} \to \mathrm{Fr}\langle TL, \overline{T}\leqslant, \lhd_0^{\mathbb{L}}\rangle$$

that extends $\heartsuit$.

(a) We first have to show that ∇ is order-preserving, that is, that if $\alpha \, \overline{T}{\leqslant} \, \beta$, then $\nabla\alpha \leqslant \nabla\beta$. But this follows immediately from (∇1).

(b) This follows immediately from (∇2').

(c) This follows immediately from (∇2').

(d) We suppose that $\alpha \vartriangleleft_0^{\mathbb{L}} \Gamma$. By the definition of $\vartriangleleft_0^{\mathbb{L}}$, there is some $\Phi \in TPL$ such that $\alpha \, \overline{T}{\leqslant} \, T\bigvee(\Phi)$ and $\lambda^T(\Phi) = \Gamma$. Now we need to show that $\nabla\alpha \leqslant \bigvee\{\nabla\beta \mid \beta \in \lambda^T(\Phi)\}$. But this is easy to see, since

$$\nabla\alpha \leqslant \nabla T \bigvee(\Phi) \qquad\qquad \text{(by (∇1))}$$
$$\leqslant \bigvee\{\nabla\beta \mid \beta \in \lambda^T(\Phi)\} \qquad\qquad \text{(by (∇3)).}$$

This completes the proof that (2)(d) holds.

Consequently, there exists a unique frame homomorphism

$$g : \mathrm{Fr}\langle TL, \overline{T}{\leqslant}, \vartriangleleft_0^{\mathbb{L}}\rangle \to V_T\mathbb{L}$$

extending ∇.

Finally, it is easy to see that

$$gf = id_{V_T\mathbb{L}}$$
$$fg = id_{\langle TL, \overline{T}{\leqslant}, \vartriangleleft_0^{\mathbb{L}}\rangle},$$

so that we do indeed have $V_T\mathbb{L} \simeq \mathrm{Fr}\langle TL, \overline{T}{\leqslant}, \vartriangleleft_0^{\mathbb{L}}\rangle$. $\qquad\square$

In the light of Theorem 3.26, we will denote the insertion of generators by

$$\nabla : TL \to \mathrm{Fr}\langle TL, \overline{T}{\leqslant}, \vartriangleleft_0^{\mathbb{L}}\rangle.$$

We now arrive at the most important corollary of Theorem 3.26, which says that every element of $V_T\mathbb{L}$ has a disjunctive normal form.

Corollary 3.27. Let $T :$ Set $\to$ Set be a standard, finitary, weak pullback-preserving functor and $\mathbb{L}$ be a frame. Then for all $x \in V_T\mathbb{L}$, there is a $\Gamma \in PTL$ such that $x = \bigvee\{\nabla\gamma \mid \gamma \in \Gamma\}$.

Proof. By Theorem 3.26 we know that $V_T\mathbb{L} \simeq \mathrm{SupLat}\langle TL, \overline{T}{\leqslant}, \vartriangleleft_0^{\mathbb{L}}\rangle$. The corollary now follows by Fact 2.13. $\qquad\square$

Remark 3.28. It is not hard to show that

$$\mathrm{SupLat}\langle TL, \overline{T}{\leqslant}, \vartriangleleft_0^{\mathbb{L}}\rangle \simeq \mathrm{SupLat}\langle TL \mid (\nabla 1), (\nabla 3)\rangle.$$

Consequently, by Theorem 3.26 and Fact 2.13, the *order* on $V_T\mathbb{L}$ is uniquely determined by the relations (∇1) and (∇3).

4. Preservation results

Now that we have established the T-powerlocale construction, we can set about proving that it is well behaved. One particular kind of good behaviour is to ask that it

preserves algebraic properties. In this section, we present several initial results in this area. We begin in Section 4.1 by briefly reviewing some of the preservation properties of V, the usual Vietoris powerlocale, and then prove that V preserves compactness. In Section 4.2, we show that V_T, the T-powerlocale construction, preserves regularity and zero-dimensionality. Finally, in Section 4.3, we show that if we assume that T maps finite sets to finite sets, then V_T preserves the combination of compactness and zero-dimensionality.

4.1. *Preservation properties of V*

There are various relations between the properties of $\mathbb{L}$ and $V\mathbb{L}$. For instance, Johnstone (1985) shows that $\mathbb{L}$ is regular, completely regular, zero-dimensional or compact regular if and only if $V\mathbb{L}$ is, and also that if $\mathbb{L}$ is locally compact, then so is $V\mathbb{L}$. The same paper also mentions without proof that if $\mathbb{L}$ is compact, then so is $V\mathbb{L}$, referring to a proof by transfinite induction similar to that used for the localic Tychonoff theorem in Johnstone (1982). The paper leaves open the converse question of whether $V\mathbb{L}$ compact implies that $\mathbb{L}$ is too. We shall give here a constructive (topos-valid) proof using preframe techniques that $\mathbb{L}$ is compact if and only if $V\mathbb{L}$ is.

Definition 4.1. A frame $\mathbb{L}$ (or, more properly, its locale) is *compact* if whenever $1 \leqslant \bigvee_i^{\uparrow} a_i$, then $1 \leqslant a_i$ for some i.

The following constructive proof is a routine application of the techniques in Johnstone and Vickers (1991).

Theorem 4.2. $\mathbb{L}$ is compact if and only if $V\mathbb{L}$ is.

Proof.

— ($\Rightarrow$):

$\mathbb{L}$ is compact if and only if the function $\mathbb{L} \to \Omega$ that maps $a \in \mathbb{L}$ to the truth value of $a = 1$ is a preframe homomorphism, that is, it preserves finite meets and directed joins. This function is characterised by being right adjoint to the unique frame homomorphism $! : \Omega \to \mathbb{L}$, so to prove compactness, it suffices to define a preframe homomorphism $\mathbb{L} \to \Omega$ and show that it is right adjoint to !. If $\mathbb{L}$ is presented – as a *frame* – by generators and relations, then the 'preframe coverage theorem' of Johnstone and Vickers (1991) shows how to derive a presentation as a preframe, which can then be used for defining preframe homomorphisms from $\mathbb{L}$. The strategy is to generate a $\vee$-semilattice from the generators and then add relations to ensure a $\vee$-*stability* condition analogous to the $\wedge$-stability used in Johnstone's coverage theorem (Johnstone 1982).

Our first step is to apply the preframe coverage theorem to derive a preframe presentation of $V\mathbb{L}$. We show

$$V\mathbb{L} \cong \mathrm{Fr}\langle P_\omega\mathbb{L} \times \mathbb{L} \ (\text{qua} \ \vee\text{-semilattice}) \mid$$
$$1 \leqslant (\gamma \cup \{1\}, d)$$
$$(\gamma \cup \{a\}, d) \wedge (\gamma \cup \{b\}, d) \leqslant (\gamma \cup \{a \wedge b\}, d)$$
$$(\gamma \cup \{\textstyle\bigvee^\uparrow A\}, d) \leqslant \textstyle\bigvee^\uparrow_{a \in A}(\gamma \cup \{a\}, d) \qquad (A \ \text{directed})$$
$$(\gamma, \textstyle\bigvee^\uparrow A \vee d) \leqslant \textstyle\bigvee^\uparrow_{a \in A}(\gamma, a \vee d) \qquad (A \ \text{directed})$$
$$(\gamma \cup \{a\}, d) \wedge (\gamma, b \vee d) \leqslant (\gamma, (a \wedge b) \vee d)$$
$$(\gamma \cup \{a \vee b\}, d) \leqslant (\gamma \cup \{a\}, b \vee d)\rangle.$$

The $\vee$-semilattice structure on $P_\omega\mathbb{L} \times \mathbb{L}$ is the product structure from $\cup$ on $P_\omega\mathbb{L}$ and $\vee$ on $\mathbb{L}$. The homomorphisms between the frame presented above and $V\mathbb{L}$ are given by

$$\Box a \mapsto (\{a\}, 0), \quad \Diamond a \mapsto (\varnothing, a)$$
$$(\gamma, d) \mapsto \bigvee_{c \in \gamma} \Box c \vee \Diamond d.$$

The $\vee$-stable, so the preframe coverage shows that

$$V\mathbb{L} \cong \mathrm{PreFr}\langle P_\omega\mathbb{L} \times \mathbb{L} \ (\text{qua poset}) \mid \text{same relations as above} \rangle.$$

We can now define a preframe homomorphism $\varphi \colon V\mathbb{L} \to \Omega$ by

$$\varphi(\gamma, d) = \exists c \in \gamma. \, c \vee d = 1.$$

To motivate this, we want criteria for $\bigvee_{c \in \gamma} \Box c \vee \Diamond d = 1$, and, intuitively, this means that for every sublocale K corresponding to a point of $V\mathbb{L}$, either K is included in some $c \in \gamma$ or K meets d. Taking K to be the closed complement of d, we get the given condition. This is not a rigorous argument, since that closed complement is not necessarily a point of $V\mathbb{L}$. However, the rest of our argument validates the choice. The relations in the preframe presentation of $V\mathbb{L}$ are generally easy to check, and we shall just mention the penultimate one. Suppose $(\gamma \cup \{a\}, d)$ and $(\gamma, b \vee d)$ are both mapped to 1. We have either some $c \in \gamma$ with $c \vee d = 1$, in which case $c \vee (a \wedge b) \vee d = 1$, or we have $a \vee d = 1$ and in addition some $c' \in \gamma$ with $c' \vee b \vee d = 1$. In this latter case, $c' \vee (a \wedge b) \vee d = 1$.

Next we show that φ is right adjoint to $! \colon \Omega \to V\mathbb{L}$, the unique frame homomorphism defined by

$$!(p) = \bigvee \{1 \mid p\} = \textstyle\bigvee^\uparrow \left(\{0\} \cup \{1 \mid p\} \right).$$

We must show $\varphi(!(p)) \geqslant p$ and $!(\varphi(\gamma, d)) \leqslant (\gamma, d)$. For the former, we have

$$\varphi(!(p)) = \varphi \left(\textstyle\bigvee^\uparrow \left(\{0\} \cup \{1 \mid p\} \right) \right)$$
$$= \varphi(\varnothing, 0) \vee \bigvee \{\varphi(\{1\}, 0) \mid p\} \geqslant p$$

since if p holds, the disjuncts include $\varphi(\{1\}, 0) = 1$. For the second inequality, we must show that

$$\bigvee \{1 \mid \varphi(\gamma, d)\} \leqslant (\gamma, d).$$

If $\varphi(\gamma, d)$ holds, then $c \vee d = 1$ for some $c \in \gamma$, so

$$1 = (\{1\}, 0) = (\{c \vee d\}, 0) \leqslant (\{c\}, d) \leqslant (\gamma, d).$$

— ($\Leftarrow$):

Suppose we have $1 = \bigvee_i^\uparrow a_i$ in $\mathbb{L}$. Then in $V\mathbb{L}$ we have $1 = \square 1 = \bigvee_i^\uparrow \square a_i$, so $1 = \square a_i$ for some i. Applying i to both sides then gives $1 = a_i$. $\square$

4.2. *Regularity and zero-dimensionality*

This section is devoted to proving that the operation V_T preserves the regularity and zero-dimensionality of frames. Both of these notions are defined in terms of the well-inside relation $\leqslant$. Accordingly, the main technical result of this section states that if $\alpha \, \overline{T} \leqslant \beta$, then $\nabla \alpha \leqslant_{V_T\mathbb{L}} \nabla \beta$ also. We first recall some notions leading up to the definition of regularity.

Definition 4.3.   Given two elements a, b of a distributive lattice $\mathbb{L}$, we say that a is *well inside* b (notation, $a \leqslant b$) if there is some c in $\mathbb{L}$ such that $a \wedge c = 0$ and $b \vee c = 1$. If $a \leqslant a$, we say a is *clopen*. We denote the clopen elements of $\mathbb{L}$ by $C_\mathbb{L}$.

If $\mathbb{L}$ is a frame, we may always take the Heyting complementation $\neg a$ of a for the element c witnessing $a \leqslant b$ in the definition of $\leqslant$. In other words, $a \leqslant b$ if and only if $b \vee \neg a = 1$. Consequently, if a is clopen, $a \vee \neg a = 1$. In the following we will use both this fact and the following properties of $\leqslant$ without further reference – see Johnstone (1982, Section III-1.1) for proofs.

Fact 4.4.   Let $\mathbb{L}$ be a frame. Then:

(1) $\leqslant \, \subseteq \, \leqslant$.
(2) $\leqslant \, ; \, \leqslant \, ; \, \leqslant \, \subseteq \, \leqslant$.
(3) for $X \in P_\omega L$, if $\forall x \in X . x \leqslant y$, then $\bigvee X \leqslant y$.
(4) for $X \in P_\omega L$, if $\forall x \in X . y \leqslant x$, then $y \leqslant \bigwedge X$.
(5) $a \leqslant a$ if and only if a has a complement.

Definition 4.5.   A frame $\mathbb{L}$ is *regular* if every $a \in \mathbb{L}$ satisfies

$$a = \bigvee \{b \in L \mid b \leqslant a\}.$$

We say $\mathbb{L}$ is *zero-dimensional* if for all $a \in \mathbb{L}$,

$$a = \bigvee \{b \in C_\mathbb{L} \mid b \leqslant a\}.$$

We will just state the following useful property of $C_\mathbb{L}$ (Johnstone 1982, Section III-1.1).

Fact 4.6.   Let $\mathbb{L}$ be a frame. Then $\langle C_\mathbb{L}, \wedge, \vee, 0, 1 \rangle$ is a sublattice of $\mathbb{L}$.

We define a function $\Downarrow : PL \to PC_\mathbb{L}$ that maps $A \in PL$ to $\downarrow A \cap C_\mathbb{L}$.

Lemma 4.7. Let $T : \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor. If $\mathbb{L}$ is a zero-dimensional frame, then:

(1) $\forall \alpha \in TL, \nabla\alpha = \bigvee\{\nabla\beta \mid \beta \in TC_{\mathbb{L}}, \beta \ \overline{T}{\leqslant}\ \alpha\}$.
(2) $\forall \Phi \in TPL, T\bigvee(\Phi) = T\bigvee \circ T \Downarrow(\Phi)$.
(3) $\forall \Phi \in TPL, \forall \alpha \in TL, [\alpha \in TC_{\mathbb{L}} \text{ and } \alpha \ \overline{T}{\leqslant}\ ; \overline{T}{\in} \Phi]$ if and only if $\alpha \in \lambda^T (T \Downarrow(\Phi))$.

Similarly to (1), if $\mathbb{L}$ is regular, then $\forall \alpha \in TL, \nabla\alpha = \bigvee\{\nabla\beta \mid \beta \in TL, \beta \ \overline{T}{\lessdot}\ \alpha\}$.

  Proof.

(1) First, observe that for all $a \in L$, we have

$$a = \bigvee\{b \in C_{\mathbb{L}} \mid b \leqslant a\} \qquad\qquad \text{(by zero-dimensionality)}$$
$$= \bigvee \Downarrow\{a\} \qquad\qquad\qquad \text{(by the definition of } \Downarrow)$$
$$= \bigvee \Downarrow \circ \eta(a) \qquad\qquad \text{(by the definition of } \eta : \mathrm{Id}_{\mathsf{Set}} \to P).$$

By relation lifting, it follows that

$$\forall \alpha \in TL, \alpha = T\bigvee \circ T \Downarrow \circ T\eta(\alpha). \qquad (36)$$

Now observe that for all $a, b \in L$, we have $b \in \Downarrow\eta(a)$ if and only if $b \in C_{\mathbb{L}}$ and $b \leqslant a$. By relation lifting, it follows that

$$\forall \alpha, \beta \in TL, \left[\beta \ \overline{T}{\in}\ T \Downarrow \circ T\eta(\alpha) \text{ iff } \beta \in TC_{\mathbb{L}} \text{ and } \beta \ \overline{T}{\leqslant}\ \alpha \right]. \qquad (37)$$

Combining these two observations, we see that

$$\nabla\alpha = \nabla \left(T\bigvee \circ T \Downarrow \circ T\eta(\alpha) \right) \qquad\qquad \text{(by (36))}$$
$$= \bigvee\{\nabla\beta \mid \beta \ \overline{T}{\in}\ T \Downarrow \circ T\eta(\alpha)\} \qquad \text{(by } (\nabla 3))$$
$$= \bigvee\{\nabla\beta \mid \beta \in TC_{\mathbb{L}}, \beta \ \overline{T}{\leqslant}\ \alpha\} \qquad \text{(by (37))}.$$

(2) It follows by the zero-dimensionality of $\mathbb{L}$ that for all $A \in PL$, we have $\bigvee A = \bigvee \Downarrow A$. Consequently, by relation lifting, (2) holds.

(3) Take $a \in L$ and $A \in PL$. Then

$$a \in \Downarrow A \Leftrightarrow a \in C_{\mathbb{L}} \text{ and } \exists b \in A, a \leqslant b \qquad\qquad \text{(by the definition of } \Downarrow)$$
$$\Leftrightarrow a \in C_{\mathbb{L}} \text{ and } a \leqslant ; \in A \qquad \text{(by the definition of relation composition)}.$$

It follows by relation lifting that

$$\forall \Phi \in TPL, \forall \alpha \in TL, \alpha \ \overline{T}{\in}\ T \Downarrow(\Phi) \text{ iff } \alpha \in TC_{\mathbb{L}} \text{ and } \alpha \ T{\leqslant}\ ; T{\in} \Phi.$$

It now follows by the definition of $\lambda^T (\Phi)$ that (3) holds.

  For the final part of the lemma, first observe that if $\mathbb{L}$ is regular, then for all $a \in L$, $a = \bigvee w(a)$, where we temporarily define $w : L \to PL$ as

$$w : a \mapsto \{b \in L \mid b \lessdot a\}.$$

By relation lifting, it follows that

$$T\bigvee \circ Tw = id_L. \qquad (38)$$

Moreover, it follows by the definition of $w : L \to PL$ that for all $a, b \in L$, $b \in w(a)$ if and only if $b \ll a$. Consequently,

$$\forall \alpha, \beta \in TL, \ \beta \ \overline{T}{\in} \ Tw(\alpha) \text{ iff } \beta \ \overline{T}{\ll} \ \alpha. \tag{39}$$

Now we see that for any $\alpha \in TL$,

$$
\begin{aligned}
\nabla \alpha &= \nabla \left(T\bigvee \circ Tw(\alpha) \right) && \text{(by (38))} \\
&= \bigvee \{ \nabla \beta \mid \beta \ \overline{T}{\in} \ Tw(\alpha) \} && \text{(by (∇3))} \\
&= \bigvee \{ \nabla \beta \mid \beta \ \overline{T}{\ll} \ \alpha \} && \text{(by (39)).}
\end{aligned}
$$

This completes the proof of the lemma. $\qquad\qquad\square$

The key technical lemma of this subsection states that relation lifting preserves the $\ll$-relation.

Lemma 4.8. Let T be a standard, finitary, weak pullback-preserving functor and $\mathbb{L}$ be a frame. Then

$$\text{for all } \alpha, \beta \in TL : \ \alpha \ \overline{T}{\ll} \ \beta \text{ implies } \nabla \alpha \ll_{V_T \mathbb{L}} \nabla \beta. \tag{40}$$

Proof. Let $\alpha, \beta \in TL$ be such that $\alpha \ \overline{T}{\ll} \ \beta$. Our aim will be to show that $\nabla \alpha \ll_{V_T \mathbb{L}} \nabla \beta$. We may assume without loss of generality that

$$\beta = (Tf)\alpha \text{ for some } f : Base^T(\alpha) \to Base^T(\beta)$$
$$\text{such that } a \ll fa \text{ for all } a \in Base^T(\alpha). \tag{41}$$

To justify this assumption, we assume that we have a proof of (40) for all β satisfying (41). To derive (40) in the general case, consider arbitrary elements $\alpha, \beta' \in TL$ such that $\alpha \ \overline{T}{\ll} \ \beta'$. In order to show that $\nabla \alpha \ \overline{T}{\ll} \ \nabla \beta'$, consider the map $f : Base^T(\alpha) \to L$ given by

$$f(a) := \bigwedge \{ b \in Base^T(\beta') \mid a \ll b \}.$$

It is not difficult to see from Fact 4.4 that $Gr(f) \subseteq \ll$, so by the properties of relation lifting, we obtain $Gr(Tf) \subseteq \overline{T}{\ll}$. In particular, we find that $\alpha \ \overline{T}{\ll} \ (Tf)\alpha$. Thus, by our assumption, we may conclude that $\nabla \alpha \ll \nabla(Tf)\alpha$. Also, note that $a \ll b$ implies $fa \leqslant b$, for all $a \in Base^T(\alpha)$ and $b \in Base^T(\beta')$. Hence, by Lemma 2.7, we may conclude from $\alpha \ \overline{T}{\ll} \ \beta'$ that $(Tf)\alpha \ \overline{T}{\leqslant} \ \beta'$, which gives $\nabla(Tf)\alpha \leqslant \nabla \beta'$. Combining our observations thus far, by Fact 4.4 it follows from $\nabla \alpha \ll \nabla(Tf)\alpha$ and $\nabla(Tf)\alpha \leqslant \nabla \beta'$ that $\nabla \alpha \ll \nabla \beta'$ does indeed hold, and thus that our assumption (41) is justified.

Turning to the proof itself, consider the map $h : P \, Base^T(\alpha) \to L$ given by

$$h(A) := \bigwedge \left(\{ \neg a \mid a \in A \} \cup \{ fa \mid a \notin A \} \right).$$

Our first observation is that since $\neg a \vee fa = 1_{\mathbb{L}}$ for each $a \in Base^T(\alpha)$ by assumption, we may infer that

$$1_{\mathbb{L}} = \bigwedge \{ \neg a \vee fa \mid a \in Base^T(\alpha) \}.$$

A straightforward application of the (finitary) distributive law then yields

$$1_{\mathbb{L}} = \bigvee \{ h(A) \mid A \in P \, Base^T(\alpha) \}. \tag{42}$$

The crucial observation now is that

$$\geqslant' \,;\, \leqslant' \,;\, (Gr\ h)^{\smile} \,\subseteq\, \notin. \tag{47}$$

To prove this, we take a pair $(a, A) \in L \times PL$ in the left-hand side of (47), and suppose in order to show a contradiction that $a \in A$. Then, by the definition of h, we obtain $h(A) \leqslant \neg a$, so $a \wedge h(A) = 0_{\mathbb{L}}$. But if $a \geqslant' \,;\, \leqslant' \,;\, (Gr\ h)^{\smile} A$, there must be some b such that $b \leqslant' a, h(A)$, and, by the definition of $\leqslant'$, this can only be the case if $b > 0_{\mathbb{L}}$. This gives the desired contradiction.

Finally, by the monotonicity of relation lifting, it is an immediate consequence of (46) and (47) that $\alpha\ \overline{T}\!\notin \Phi$.

This completes the proof of Claim 1. $\qquad\square$

Given Claim 1, the rest of the proof is straightforward. We define

$$c := \bigvee \left\{ Th(\Phi) \mid \Phi \in TP\,Base^T(\alpha) \text{ such that } (\alpha, \Phi) \notin \overline{T}\!\notin \right\}.$$

Then we may calculate that

$$c \vee \nabla\beta$$
$$\geqslant c \vee \bigvee \left\{ Th(\Phi) \mid \Phi \in TP\,Base^T(\alpha) \text{ such that } (\alpha, \Phi) \in \overline{T}\!\notin \right\} \qquad \text{(Claim 1(a))}$$
$$= \bigvee \left\{ Th(\Phi) \mid \Phi \in TP\,Base^T(\alpha) \right\} \qquad\qquad\qquad \text{(definition of } c\text{)}$$
$$= 1_{V_T\mathbb{L}} \qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad \text{(equation (44))}$$

and

$$\nabla\alpha \wedge c$$
$$= \bigvee \left\{ \nabla\alpha \wedge Th(\Phi) \mid \Phi \in TP\,Base^T(\alpha) \text{ such that } (\alpha, \Phi) \notin \overline{T}\!\notin \right\} \quad \text{(distributivity)}$$
$$= \bigvee \left\{ 0_{V_T\mathbb{L}} \mid \Phi \in TP\,Base^T(\alpha) \text{ such that } (\alpha, \Phi) \notin \overline{T}\!\notin \right\} \qquad \text{(Claim 1(b))}$$
$$= 0_{V_T\mathbb{L}}.$$

In other words, c witnesses the fact that $\nabla\alpha \leqslant_{V_T\mathbb{L}} \nabla\beta$. $\qquad\square$

We now arrive at the main result of this section, namely, that the T-powerlocale construction preserves regularity and zero-dimensionality.

Theorem 4.9. Let $\mathbb{L}$ be a frame and T be a standard, finitary, weak pullback-preserving functor. Then:

(1) If $\mathbb{L}$ is regular, then so is $V_T\mathbb{L}$.
(2) If $\mathbb{L}$ is zero-dimensional, then so is $V_T\mathbb{L}$.

Proof.

(1) By Corollary 3.27, it suffices to show that for all $\alpha \in TL$,

$$\nabla\alpha = \bigvee\{\nabla\beta \in V_T\mathbb{L} \mid \nabla\beta \lessdot \nabla\alpha\}. \tag{48}$$

Take $\alpha \in TL$. We see that

$$
\begin{aligned}
\nabla\alpha &= \bigvee\{\nabla\beta \mid \beta \; \overline{T}{\lessdot}\; \alpha\} && \text{(by Lemma 4.7)} \\
&\leqslant \bigvee\{\nabla\beta \mid \nabla\beta \lessdot_{V_T\mathbb{L}} \nabla\alpha\} && \text{(by Lemma 4.8)} \\
&\leqslant \nabla\alpha && \text{(since } \lessdot \,\subseteq\, \leqslant).
\end{aligned}
$$

It follows that (48) holds, which concludes the proof of part (1).

(2) Again by Corollary 3.27, it suffices to show that for all $\alpha \in TL$,

$$
\nabla\alpha = \bigvee\{\nabla\beta \mid \nabla\beta \in C_{V_T\mathbb{L}}, \nabla\beta \leqslant \nabla\alpha\}. \tag{49}
$$

The main observation here is that

$$
\begin{aligned}
&\forall \beta \in TC_{\mathbb{L}} \\
&\nabla\beta \in C_{V_T\mathbb{L}}.
\end{aligned} \tag{50}
$$

To see why this is the case, recall that $C_{\mathbb{L}} := \{b \in L \mid b \lessdot b\}$, so for all $b \in C_{\mathbb{L}}$, we have $b = b$ implies $b \lessdot b$. Consequently, by relation lifting,

$$
\forall \beta \in TC_{\mathbb{L}}, \; \beta \; \overline{T}{\lessdot}\; \beta.
$$

It then follows by Lemma 4.8 that (50) holds. Now,

$$
\begin{aligned}
\nabla\alpha &= \bigvee\{\nabla\beta \mid \beta \in TC_{\mathbb{L}}, \beta \; \overline{T}{\leqslant}\; \alpha\} && \text{(by Lemma 4.7(1))} \\
&\leqslant \bigvee\{\nabla\beta \in C_{V_T\mathbb{L}} \mid \beta \; \overline{T}{\leqslant}\; \alpha\} && \text{(by (50))} \\
&\leqslant \bigvee\{\nabla\beta \in C_{V_T\mathbb{L}} \mid \nabla\beta \leqslant \nabla\alpha\} && \text{(by (∇1))} \\
&= \nabla\alpha && \text{(by order theory).}
\end{aligned}
$$

It now follows that (49) holds, so (2) also holds. $\qquad\square$

4.3. *Compactness + zero-dimensionality*

In this section we will show that if $\mathbb{L}$ is compact and zero-dimensional, then so is $V_T\mathbb{L}$. Our proof strategy is as follows. Given a compact zero-dimensional frame $\mathbb{L}$, we will define a new construction $V_T^C\mathbb{L}$ that is guaranteed to be compact, and then show that $V_T\mathbb{L} \simeq V_T^C\mathbb{L}$.

We define a flat-site presentation $\langle TC_{\mathbb{L}}, \overline{T}{\leqslant}, \lessdot_0^C \rangle$, where

$$
\lessdot_0^C := \{(\alpha, \lambda^T(\Phi)) \in TC_{\mathbb{L}} \times PTL \mid \alpha \; \overline{T}{\leqslant}\; T\bigvee(\Phi), \Phi \in TP_\omega C_{\mathbb{L}}\}.
$$

Note that we view $TC_{\mathbb{L}}$ as a substructure of TL, which is justified by the fact that $C_{\mathbb{L}}$ is a sublattice of $\mathbb{L}$ (Fact 4.6): this fact tells us that $\bigvee : PL \to L$ restricts to a function from $P_\omega C_{\mathbb{L}}$ to $C_{\mathbb{L}}$, so $T\bigvee$ maps $TP_\omega C_{\mathbb{L}}$ to $TC_{\mathbb{L}}$ by the standardness of T. We will need the following property of relation lifting with respect to ordered sets.

Lemma 4.10. Let $T : \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor and $\mathbb{P}$ be a poset with a top element 1. Then for every $\beta \in T\mathbb{P}$, there is some $\alpha \in T\{1\}$ such that $\beta \; \overline{T}{\leqslant}\; \alpha$.

Proof. Consider the following function at the ground level: $f : P \to \{1\}$, where f is the constant function $f : b \mapsto 1$. Then for all $b \in P$, we have $b \leqslant f(b)$ and $f(b) \in \{1\}$. By relation lifting, we have for all $\beta \in TP$, that $\beta \, \overline{T}{\leqslant}\, Tf(\beta)$ and $Tf(\beta) \in T\{1\}$. The statement then follows. $\qquad\square$

Lemma 4.11. Let $T : \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor and let $\mathbb{L}$ be a frame. Then $\langle TC_{\mathbb{L}}, \overline{T}{\leqslant}, \lhd_0^C \rangle$ is a flat site. Moreover, if T maps finite sets to finite sets, then $\mathrm{Fr}\langle TC_{\mathbb{L}}, \overline{T}{\leqslant}, \lhd_0^C \rangle$ is a compact frame.

Proof. Because $C_{\mathbb{L}}$ is a meet-subsemilattice of $\mathbb{L}$, we can apply Lemma 3.23 to $TC_{\mathbb{L}}$. Now the proof that $\langle TC_{\mathbb{L}}, \overline{T}{\leqslant}, \lhd_0^C \rangle$ is a flat site is analogous to that of Lemma 3.25.

Now suppose that T maps finite sets to finite sets. Then for all $\Phi \in TP_\omega C_{\mathbb{L}}$, it follows by Fact 2.11(3) that $\lambda^T(\Phi)$ is finite. Consequently,

$$\forall \alpha \lhd_0^C \lambda^T(\Phi), \ \lambda^T(\Phi) \text{ is finite.}$$

Moreover, by Lemma 4.10,

$$TC_{\mathbb{L}} = \downarrow_{TC_{\mathbb{L}}} T\{1_{\mathbb{L}}\}$$

since $1_{\mathbb{L}} \in C_{\mathbb{L}}$ as $C_{\mathbb{L}}$ is a sublattice of $\mathbb{L}$. Since we assumed that T maps finite sets to finite sets, the set $T\{1_{\mathbb{L}}\}$ must be finite. It now follows from a straightforward generalisation of Vickers (2006, Proposition 11) that $\mathrm{Fr}\langle TC_{\mathbb{L}}, \overline{T}{\leqslant}, \lhd_0^C \rangle$ is a compact frame. (The only change we need to make to Vickers (2006, Proposition 11) is to generalise from using single finite trees to using disjoint unions of $|T\{1_{\mathbb{L}}\}|$-many trees so that we can cover each element of $T\{1_{\mathbb{L}}\}$.) $\qquad\square$

We define $V_T^C\mathbb{L} := \mathrm{Fr}\langle TC_{\mathbb{L}}, \overline{T}{\leqslant}, \lhd_0^C \rangle$, and for the time being we denote the insertion of generators by $\heartsuit : TC_{\mathbb{L}} \to V_T^C\mathbb{L}$. Our goal now is to show that $V_T\mathbb{L} \simeq V_T^C\mathbb{L}$. We will use a shortcut, exploiting the fact that both $V_T\mathbb{L}$ and $V_T^C\mathbb{L}$ have flat-site presentations: we will define *suplattice* homomorphisms $f' : V_T\mathbb{L} \to V_T^C\mathbb{L}$ and $g' : V_T^C\mathbb{L} \to V_T\mathbb{L}$. We will then show that $g' \circ f' = id$ and $f' \circ g' = id$, so $V_T\mathbb{L}$ and $V_T^C\mathbb{L}$ are isomorphic as suplattices. It then follows from order theory that they are also isomorphic as frames. We begin by defining a function $g : TC_{\mathbb{L}} \to V_T\mathbb{L}$, defined as

$$g : \alpha \mapsto \nabla\alpha.$$

Lemma 4.12. Let $T : \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor and $\mathbb{L}$ be a frame. Then the function g defined above extends to a *suplattice* homomorphism $g' : V_T^C\mathbb{L} \to V_T\mathbb{L}$ such that $g' \circ \heartsuit = g$.

Proof. We need to show that $g : TC_{\mathbb{L}} \to V_T\mathbb{L}$ preserves the order on $TC_{\mathbb{L}}$ and preserves covers into joins. If $\alpha \lhd_0^C \lambda^T(\Phi)$, where $\alpha \in TC_{\mathbb{L}}$, $\Phi \in TPC_{\mathbb{L}}$ and $\alpha \, \overline{T}{\leqslant} \bigvee(\Phi)$, then

$g(\alpha) \leqslant \bigvee\{g(\beta) \mid \beta \in \lambda^T(\Phi)\}$. Both of the required properties then follow straightforwardly from the fact that $\langle TC_{\mathbb{L}}, \overline{T}\leqslant, \lessdot_0^C \rangle$ is a substructure of $\langle TL, \overline{T}\leqslant, \lessdot_0^{\mathbb{L}} \rangle$. $\qquad\square$

The next step is to define the suplattice homomorphism $f' : V_T\mathbb{L} \to V_T^C\mathbb{L}$. This requires a little more work than the definition of $g' : V_T^C\mathbb{L} \to V_T\mathbb{L}$, beginning with the following lemma.

Lemma 4.13. Let $T : \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor and $\mathbb{L}$ be a compact frame. If $\alpha \in TC_{\mathbb{L}}$ and $\Phi \in TPC_{\mathbb{L}}$ such that $\alpha \, \overline{T}\leqslant T\bigvee(\Phi)$, then there exists $\Phi_\alpha \in TP_\omega C_{\mathbb{L}}$ such that $\Phi_\alpha \, \overline{T}\subseteq_L \Phi$ and $\alpha \, \overline{T}\leqslant T\bigvee(\Phi_\alpha)$.

Proof. Since $\mathbb{L}$ is compact, we can show that

$$\text{for all } a \in C_{\mathbb{L}}, a \text{ is compact.} \tag{51}$$

After all, if $a \in C_{\mathbb{L}}$ and $A \in PL$ such that $a \leqslant \bigvee A$, then we also have

$$1 \leqslant a \vee \neg a \leqslant \bigvee A \cup \{\neg a\},$$

so by the compactness of $\mathbb{L}$, there exists a finite $A' \subseteq A$ such that

$$a \vee \neg a \leqslant \bigvee A' \cup \{\neg a\}.$$

Consequently, $a \leqslant \bigvee A'$. Since A was arbitrary, it follows that a is compact.

We define

$$S := \left(\leqslant \, ; Gr(\bigvee)^{\smallsmile} \right) \restriction_{C_{\mathbb{L}} \times PC_{\mathbb{L}}};$$

so that $(a, A) \in S$ if and only if $a \in C_{\mathbb{L}}$, $A \in PC_{\mathbb{L}}$ and $a \leqslant \bigvee A$. By (51), we can define a function $h : S \to S$ where $h : (a, A) \mapsto (a', A')$ such that $a = a'$, $A' \subseteq A$, $a' \leqslant \bigvee A'$ (otherwise h would not be well defined) and such that A' is finite, that is, $A' \in P_\omega C_{\mathbb{L}}$. In other words, $h : S \to S$ is a function that assigns a finite subcover A' to a set of zero-dimensional opens A covering a zero-dimensional open element a . If we denote the projection functions of S by

$$C_{\mathbb{L}} \xleftarrow{\ p_1\ } S \xrightarrow{\ p_2\ } PC_{\mathbb{L}},$$

we can encode the above-mentioned properties of h as follows:

$$\forall x \in S, p_1 \circ h(x) = p_1(x)$$
$$\forall x \in S, p_2 \circ h(x) \subseteq p_2(x)$$
$$\forall x \in S, p_2 \circ h(x) \in P_\omega C_{\mathbb{L}}.$$

By relation lifting, it follows that

$$\forall x \in TS, Tp_1 \circ Th(x) = Tp_1(x) \tag{52}$$
$$\forall x \in TS, Tp_2 \circ Th(x) \, \overline{T}\subseteq Tp_2(x) \tag{53}$$
$$\forall x \in TS, Tp_2 \circ Th(x) \in TP_\omega C_{\mathbb{L}}. \tag{54}$$

Finally, observe that it follows by relation lifting that

$$\forall \alpha \in TC_{\mathbb{L}}, \forall \Phi \in TPC_{\mathbb{L}}, \alpha \, \overline{T}\leqslant \bigvee(\Phi) \text{ iff } \alpha \, \overline{T}S \, \Phi.$$

Now take $\alpha \in TC_{\mathbb{L}}$ and $\Phi \in TPC_{\mathbb{L}}$ such that $\alpha \overline{T} \leqslant \bigvee(\Phi)$. Then by the above, we have $\alpha \overline{TS} \Phi$, so by the definition of $\overline{T}$, there must exist some $x \in TS$ such that $Tp_1(x) = \alpha$ and $Tp_2(x) = \Phi$. We define $\Phi_\alpha := Tp_2 \circ Th(x)$; note that $Tp_1 \circ Th(x) = Tp_1(x) = \alpha$ by (52). Since Th is a function from TS to TS, we have $\alpha \overline{TS} \Phi_\alpha$, so $\alpha \overline{T} \leqslant T\bigvee(\Phi_\alpha)$. Moreover, by (53), we have $\Phi_\alpha \overline{T} \subseteq \Phi$ and by (54), we have $\Phi_\alpha \in TP_\omega C_{\mathbb{L}}$. This concludes the proof. $\square$

We now define a map $f: TL \to V_T^C \mathbb{L}$ by sending

$$f: \alpha \mapsto \bigvee\{\heartsuit\beta \mid \beta \in TC_{\mathbb{L}}, \beta \overline{T} \leqslant \alpha\}.$$

This will give us our suplattice homomorphism $f': V_T \mathbb{L} \to V_T^C \mathbb{L}$.

Lemma 4.14. Let $T: \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving functor. If $\mathbb{L}$ is a compact zero-dimensional frame, then $f: TL \to V_T^C \mathbb{L}$ defined above extends to a *suplattice* homomorphism $f': V_T \mathbb{L} \to V_T^C \mathbb{L}$, where $f' \circ \nabla = f$.

Proof. In order to show that $f: TL \to V_T^C \mathbb{L}$ extends to a suplattice homomorphism, we need to show that f preserves the order on TL and f transforms covers into joins, that is, that for all $(\alpha, \lambda^T(\Phi)) \in \vartriangleleft_0$, where $\alpha \overline{T} \leqslant T\bigvee(\Phi)$, we have $f(\alpha) \leqslant \bigvee\{f(\gamma) \mid \gamma \in \lambda^T(\Phi)\}$. To see why f is order-preserving, we suppose that $\alpha_0, \alpha_1 \in TL$ and $\alpha_0 \overline{T} \leqslant \alpha_1$. Then

$$
\begin{aligned}
f(\alpha_0) &= \bigvee\{\heartsuit\beta \mid \beta \in TC_{\mathbb{L}}, \beta \overline{T} \leqslant \alpha_0\} && \text{(by the definition of } f) \\
&\leqslant \bigvee\{\heartsuit\beta \mid \beta \in TC_{\mathbb{L}}, \beta \overline{T} \leqslant \alpha_1\} && \text{(since } \beta \overline{T} \leqslant \alpha_0 \overline{T} \leqslant \alpha_1 \Rightarrow \beta \overline{T} \leqslant \alpha_1) \\
&= f(\alpha_1) && \text{(by the definition of } f).
\end{aligned}
$$

Before we show that f transforms covers $\alpha \vartriangleleft_0 \lambda^T(\Phi)$ into joins, we will first show that the expression $\bigvee\{f(\gamma) \mid \gamma \in \lambda^T(\Phi)\}$ can be simplified:

$$\forall \Phi \in TPL, \bigvee\{f(\gamma) \mid \gamma \in \lambda^T(\Phi)\} = \bigvee\{\heartsuit\beta \mid \beta \in \lambda^T(T\Downarrow(\Phi))\}. \tag{55}$$

To see how we do this, observe that

$$
\begin{aligned}
&\bigvee\{f(\gamma) \mid \gamma \in \lambda^T(\Phi)\} \\
&= \bigvee\left\{\bigvee\{\heartsuit\beta \mid \beta \in TC_{\mathbb{L}}, \beta \leqslant \gamma\} \mid \gamma \in \lambda^T(\Phi)\right\} && \text{(by the definition of } f) \\
&= \bigvee\left\{\bigvee\{\heartsuit\beta \mid \beta \in TC_{\mathbb{L}}, \beta \leqslant \gamma\} \mid \gamma \overline{T} \in \Phi\right\} && \text{(by the definition of } \lambda^T) \\
&= \bigvee\{\heartsuit\beta \mid \beta \in TC_{\mathbb{L}}, \exists \gamma \overline{T} \in \Phi, \beta \leqslant \gamma\} && \text{(by associativity of } \bigvee) \\
&= \bigvee\{\heartsuit\beta \mid \beta \in TC_{\mathbb{L}}, \beta \overline{T} \leqslant ; \overline{T} \in \Phi\} && \text{(by the definition of relation composition)} \\
&= \bigvee\{\heartsuit\beta \mid \beta \in \lambda^T(T\Downarrow(\Phi))\} && \text{(by Lemma 4.7(3)).}
\end{aligned}
$$

Let $\alpha \in TL$ and $\Phi \in TPL$ such that $\alpha \ \overline{T}{\leqslant}\ T\bigvee(\Phi)$; we need to show that $f(\alpha) \leqslant \bigvee\{f(\gamma) \mid \gamma \in \lambda^T(\Phi)\}$. By (55), it suffices to show that

$$f(\alpha) \leqslant \bigvee\{\heartsuit\gamma \mid \gamma \in \lambda^T(T\Downarrow(\Phi))\}. \tag{56}$$

Recall that

$$f(\alpha) = \bigvee\{\heartsuit\beta \mid \beta \in TC_{\mathbb{L}}, \beta \leqslant \alpha\}.$$

We will show that

$$\forall \beta \in TC_{\mathbb{L}}, \beta \ \overline{T}{\leqslant}\ \alpha \Rightarrow \heartsuit\beta \leqslant \bigvee\{\heartsuit\gamma \mid \gamma \in \lambda^T(T\Downarrow(\Phi))\}. \tag{57}$$

Suppose $\beta \in TC_{\mathbb{L}}$ and $\beta \ \overline{T}{\leqslant}\ \alpha$. Then, since we have assumed that $\alpha \ \overline{T}{\leqslant}\ T\bigvee(\Phi)$, it follows that $\beta \ \overline{T}{\leqslant}\ T\bigvee(\Phi)$. By Lemma 4.7(2), we know that $T\bigvee(\Phi) = T\bigvee \circ T\Downarrow(\Phi)$, so

$$\beta \ \overline{T}{\leqslant}\ T\bigvee \circ T\Downarrow(\Phi).$$

Now, since $T\Downarrow(\Phi) \in TPC_{\mathbb{L}}$, we can apply Lemma 4.13 to conclude that there must be some $\Phi' \in TP_\omega C_{\mathbb{L}}$ such that $\Phi' \ \overline{T}{\subseteq}\ T\Downarrow(\Phi)$ and $\beta \ \overline{T}{\leqslant}\ \bigvee\Phi'$. It now follows from the definition of $\lhd_0^C$ that $\beta \lhd_0^C \lambda^T(\Phi')$. Now

$$\begin{aligned} \heartsuit\beta &\leqslant \bigvee\{\heartsuit\gamma \mid \gamma \in \lambda^T(\Phi')\} & \text{(since } \beta \lhd_0^C \lambda^T(\Phi')) \\ &\leqslant \bigvee\{\heartsuit\gamma \mid \gamma \in \lambda^T(T\Downarrow(\Phi))\} & \text{(by Lemma 3.24 since } \Phi' \ \overline{T}{\subseteq}\ T\Downarrow(\Phi)). \end{aligned}$$

Since $\beta \in TC_{\mathbb{L}}$ was arbitrary, it follows that (57) holds. Consequently, (56) holds, so we may indeed conclude that f transforms covers into joins. We conclude that $f : TL \to V_T^C\mathbb{L}$ extends to a suplattice homomorphism $f' : V_T\mathbb{L} \to V_T^C\mathbb{L}$. $\qquad\square$

Now that we have established the existence of suplattice homomorphisms

$$\begin{aligned} f' &: V_T\mathbb{L} \to V_T^C\mathbb{L} \\ g' &: V_T^C\mathbb{L} \to V_T\mathbb{L}, \end{aligned}$$

we are ready to prove the theorem of this section.

Theorem 4.15. Let $T : \mathsf{Set} \to \mathsf{Set}$ be a standard, finitary, weak pullback-preserving set functor that maps finite sets to finite sets and let $\mathbb{L}$ be a frame. If $\mathbb{L}$ is compact and zero-dimensional, then so is $V_T\mathbb{L}$.

Proof. It follows by Theorem 4.9 that $V_T\mathbb{L}$ is zero-dimensional. To show that $V_T\mathbb{L}$ is compact, it suffices to show that $V_T\mathbb{L} \simeq V_T^C\mathbb{L}$ by Lemma 4.11. We will establish that $V_T\mathbb{L} \simeq V_T^C\mathbb{L}$ by showing that $g' : V_T^C\mathbb{L} \to V_T\mathbb{L}$ and $f' : V_T\mathbb{L} \to V_T^C\mathbb{L}$ are *suplattice isomorphisms*, because $g' \circ f' = id_{V_T\mathbb{L}}$ and $f' \circ g' = id_{V_T^C\mathbb{L}}$. This is sufficient, since by order theory, any suplattice isomorphism is also a frame isomorphism. We begin by making the following claim:

$$\forall \alpha \in TL, g' \circ f(\alpha) = \nabla\alpha. \tag{58}$$

After all, if $\alpha \in TL$, then

$$
\begin{aligned}
g' \circ f(\alpha) &= g'\left(\bigvee\{\heartsuit\beta \mid \beta \in TC_{\mathbb{L}}, \beta\ \overline{T}\leqslant \alpha\}\right) && \text{(by the definition of } f) \\
&= \bigvee\{g'(\heartsuit\beta) \mid \beta \in TC_{\mathbb{L}}, \beta\ \overline{T}\leqslant \alpha\} && \text{(since } g' \text{ preserves } \bigvee) \\
&= \bigvee\{g(\beta) \mid \beta \in TC_{\mathbb{L}}, \beta\ \overline{T}\leqslant \alpha\} && \text{(by Lemma 4.12)} \\
&= \bigvee\{\nabla\beta \mid \beta \in TC_{\mathbb{L}}, \beta\ \overline{T}\leqslant \alpha\} && \text{(by the definition of } g) \\
&= \nabla\alpha && \text{(by Lemma 4.7(1)).}
\end{aligned}
$$

It then follows that (58) holds. Conversely, we claim that

$$
\forall \alpha \in TC_{\mathbb{L}},\ f' \circ g(\alpha) = \heartsuit\alpha. \tag{59}
$$

This is also not hard to see since if we take $\alpha \in TC_{\mathbb{L}}$, then

$$
\begin{aligned}
f' \circ g(\alpha) &= f'(\nabla\alpha) && \text{(by the definition of } g) \\
&= f(\alpha) && \text{(by Lemma 4.14)} \\
&= \bigvee\{\heartsuit\beta \mid \beta \in TC_{\mathbb{L}}, \beta\ \overline{T}\leqslant \alpha\} && \text{(by the definition of } f) \\
&= \heartsuit\alpha && \text{(since } \alpha \in TC_{\mathbb{L}} \text{ and } \heartsuit \text{ is order-preserving).}
\end{aligned}
$$

It follows that (59) holds. Now we see that for all $\alpha \in TL$,

$$
\begin{aligned}
g' \circ f'(\nabla\alpha) &= g' \circ f(\alpha) && \text{(since } f' \circ \nabla = f) \\
&= \nabla\alpha && \text{(by (58))} \\
&= id_{V_T\mathbb{L}}(\nabla\alpha).
\end{aligned}
$$

In other words, $g' \circ f'$ and $id_{V_T\mathbb{L}}$ agree on the generators of $V_T\mathbb{L}$. It follows that $g' \circ f' = id_{V_T\mathbb{L}}$. An analogous argument shows that $f' \circ g' = id_{V_T^C\mathbb{L}}$. We conclude that $V_T\mathbb{L}$ and $V_T^C\mathbb{L}$ are isomorphic as suplattices and consequently also as frames. It follows that $V_T\mathbb{L}$ is compact. $\qquad\square$

5. Future work

We will conclude this paper by listing some open problems and directions for future work.

5.1. *Preservation properties*

The main technical problems we would like to solve are concerned with possible further preservation properties of our construction. In particular, we are very eager to find out for which functors T the T-power construction preserves compactness, or the combination of compactness and regularity. Note that any functor satisfying this property must map finite sets to finite sets: if TA were infinite for some finite A subset of $\mathbb{L}$, then we could have $1_{V_T\mathbb{L}} = \bigvee\{\nabla\alpha \mid \alpha \in A\}$, without there being a finite subcover. We conjecture that this condition (that is, of T restricting to finite sets) is in fact not only necessary, but also sufficient to prove the preservation of compactness.

5.2. *Functorial properties*

In Section 3.4 we saw that certain natural transformations $\rho : T' \to T$ induce natural transformations $\widehat{\rho} : V_T \to V_{T'}$, with the unit of the Vietoris comonad V_{P_ω} providing an instance of this phenomenon. There are some natural open questions related to this. In particular, we are interested, for the case where T is actually a *monad*, in whether V_T is a co-monad.

Another question related to the natural transformation $\widehat{\rho}$ is whether $\widehat{\rho}_{\mathbb{L}} : V_T \mathbb{L} \to V_{T'} \mathbb{L}$ always has a right adjoint – see Remark 3.22.

5.3. *Spatiality and compact Hausdorff spaces*

Palmigiano and Venema (2007) introduced a lifting construction on Chu spaces to prove that for Stone spaces, the Vietoris construction can be generalised from the power set case to an arbitrary set functor T (meeting the same constraints as in the current paper). We are led to ask whether we can generalise this to arbitrary topological spaces, or at least to compact Hausdorff spaces.

Assume that, for any functor T mapping finite sets to finite sets, we can prove that our T-powerlocale construction V_T preserves the combination of compactness and regularity. Then, using the well-known duality between compact regular locales and compact Hausdorff spaces, we obtain a Vietoris-like functor on compact Hausdorff spaces for free. The question is then whether we can give a more direct, insightful description of this functor.

5.4. *Locales and constructivity*

In this paper, we have mostly adopted a frame- rather than a locale-oriented perspective. However, Theorem 3.21 suggests that if we want to understand the relationship between coalgebra functors $T : \mathsf{Set} \to \mathsf{Set}$ and the V_T construction, we should think of V_T as a functor on *locales*, since natural transformations $T' \to T$ satisfying the conditions of Theorem 3.21 correspond to frame natural transformations $V_T \to V_{T'}$. It would be interesting to pursue this idea further, especially in conjunction with the use of *constructive mathematics*. We have seen that certain constructive techniques, such as frame, flat-site and preframe presentations, can be brought over to the framework of coalgebraic logic. Making the entire approach of this paper constructive would be a lot of work, but we believe that this would be a promising line for further research.

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